

Chapter 11

Trees

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11 A-1

Terminology

- Definition of a general tree
 - A general tree T is a set of one or more nodes such that T is partitioned into disjoint subsets:
 - A single node r , the root
 - Sets that are general trees, called subtrees of r
- Trees have hierarchical structure
 - parent/child/sibling
 - root/leaf
 - ancestor/descendant
 - subtree

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11 A-2

Terminology

- Definition of a binary tree
 - A binary tree is a set T of nodes such that either
 - T is empty, or
 - T is partitioned into three disjoint subsets:
 - A single node r , the root
 - Two possibly empty sets that are binary trees, called left and right subtrees of r

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11 A-3

Terminology

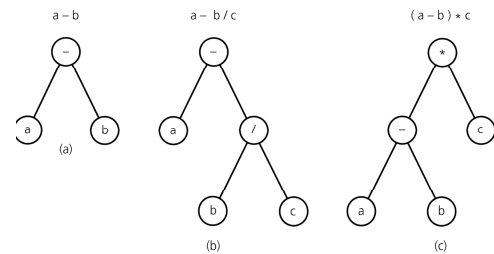


Figure 11-4
Binary trees that represent algebraic expressions

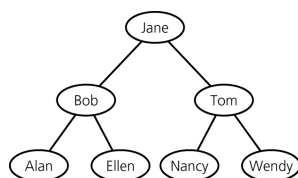
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11 A-4

Terminology

- A binary search tree
 - A binary tree that has the following properties for each node n
 - n 's value is greater than all values in its left subtree T_L
 - n 's value is less than all values in its right subtree T_R
 - Both T_L and T_R are binary search trees

Figure 11-5
A binary search tree of names



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11 A-5

Terminology

- The height of trees
 - Level of a node n in a tree T
 - If n is the root of T , it is at level 1
 - If n is not the root of T , its level is 1 greater than the level of its parent
 - Height of a tree T defined in terms of the levels of its nodes
 - If T is empty, its height is 0
 - If T is not empty, its height is equal to the maximum level of its nodes

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11 A-6

Terminology

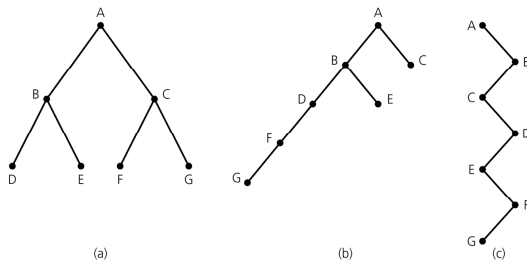


Figure 11-6
Binary trees with the same nodes but different heights

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11 A-7

Terminology

- Full, complete, and balanced binary trees
 - Recursive definition of a full binary tree
 - If T is empty, T is a full binary tree of height 0
 - If T is not empty and has height $h > 0$, T is a full binary tree if its root's subtrees are both full binary trees of height $h - 1$

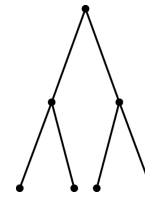


Figure 11-7
A full binary tree of height 3

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11 A-8

Terminology

- Complete binary trees
 - A binary tree T of height h is complete if
 - All nodes at level $h - 2$ and above have two children each, and
 - When a node at level $h - 1$ has children, all nodes to its left at the same level have two children each, and
 - When a node at level $h - 1$ has one child, it is a left child

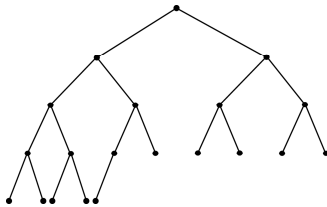


Figure 11-8
A complete binary tree

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Terminology

- Balanced binary trees
 - A binary tree is balanced if the height of any node's right subtree differs from the height of the node's left subtree by no more than 1
- Full binary trees are complete
- Complete binary trees are balanced

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11 A-10

Terminology

- Summary of tree terminology
 - General tree
 - A set of one or more nodes, partitioned into a root node and subsets that are general subtrees of the root
 - Parent of node n
 - The node directly above node n in the tree
 - Child of node n
 - A node directly below node n in the tree
 - Root
 - The only node in the tree with no parent

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11 A-11

Terminology

- Summary of tree terminology (Continued)
 - Leaf
 - A node with no children
 - Siblings
 - Nodes with a common parent
 - Ancestor of node n
 - A node on the path from the root to n
 - Descendant of node n
 - A node on a path from n to a leaf
 - Subtree of node n
 - A tree that consists of a child (if any) of n and the child's descendants

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11 A-12

Terminology

- Summary of tree terminology (Continued)
 - Height
 - The number of nodes on the longest path from the root to a leaf
 - Binary tree
 - A set of nodes that is either empty or partitioned into a root node and one or two subsets that are binary subtrees of the root
 - Each node has at most two children, the left child and the right child
 - Left (right) child of node n
 - A node directly below and to the left (right) of node n in a binary tree

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11 A-13

Terminology

- Summary of tree terminology (Continued)
 - Left (right) subtree of node n
 - In a binary tree, the left (right) child (if any) of node n plus its descendants
 - Binary search tree
 - A binary tree where the value in any node n is greater than the value in every node in n's left subtree, but less than the value of every node in n's right subtree
 - Empty binary tree
 - A binary tree with no nodes

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11 A-14

Terminology

- Summary of tree terminology (Continued)
 - Full binary tree
 - A binary tree of height h with no missing nodes
 - All leaves are at level h and all other nodes each have two children
 - Complete binary tree
 - A binary tree of height h that is full to level h - 1 and has level h filled in from left to right
 - Balanced binary tree
 - A binary tree in which the left and right subtrees of any node have heights that differ by at most 1

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Possible Representations of a Binary Tree

- An array-based representation of a complete tree
 - If the binary tree is complete and remains complete
 - A memory-efficient array-based implementation can be used

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11 A-16

Possible Representations of a Binary Tree

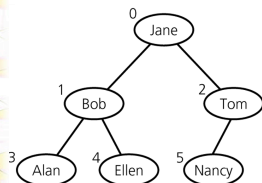


Figure 11-12
Level-by-level numbering of a complete binary tree

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0	Jane
1	Bob
2	Tom
3	Alan
4	Ellen
5	Nancy
6	
7	

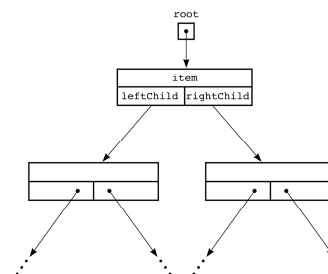
Figure 11-13
An array-based implementation of the complete binary tree in Figure 10-12

11 A-17

Possible Representations of a Binary Tree

- A reference-based representation
 - Java references can be used to link the nodes in the tree

Figure 11-14
A reference-based implementation of a binary tree



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11 A-18

A Comparison of Sorting Algorithms

	<u>Worst case</u>	<u>Average case</u>
Selection sort	n^2	n^2
Bubble sort	n^2	n^2
Insertion sort	n^2	n^2
Mergesort	$n * \log n$	$n * \log n$
Quicksort	n^2	$n * \log n$
Radix sort	n	n
Treesort	n^2	$n * \log n$
Heapsort	$n * \log n$	$n * \log n$

Figure 10-22

Approximate growth rates of time required for eight sorting algorithms

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11 A-19

Traversals of a Binary Tree

- A traversal algorithm for a binary tree visits each node in the tree
- Recursive traversal algorithms
 - Preorder traversal
 - Inorder traversal
 - Postorder traversal
- Traversal is $O(n)$

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Traversal of a Binary Tree

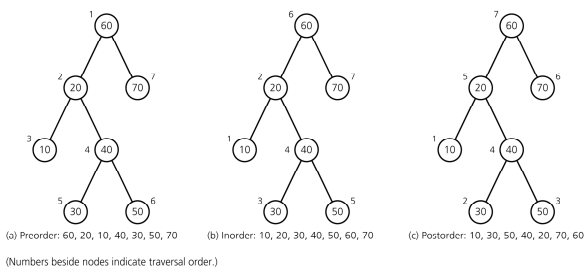


Figure 11-10

Traversals of a binary tree: a) preorder; b) inorder; c) postorder

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11 A-21

Heaps

- A heap is a complete binary tree
 - That is empty
- or
- Whose root contains a search key greater than or equal to the search key in each of its children, and
- Whose root has heaps as its subtrees

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Heaps

- Maxheap
 - A heap in which the root contains the item with the largest search key
- Minheap
 - A heap in which the root contains the item with the smallest search key

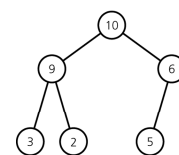
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11 A-23

Heaps: An Array-based Implementation of a Heap

- Data fields
 - items: an array of heap items
 - size: an integer equal to the number of items in the heap

Figure 12-11
A heap with its array representation



0	10
1	9
2	6
3	3
4	2
5	5

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