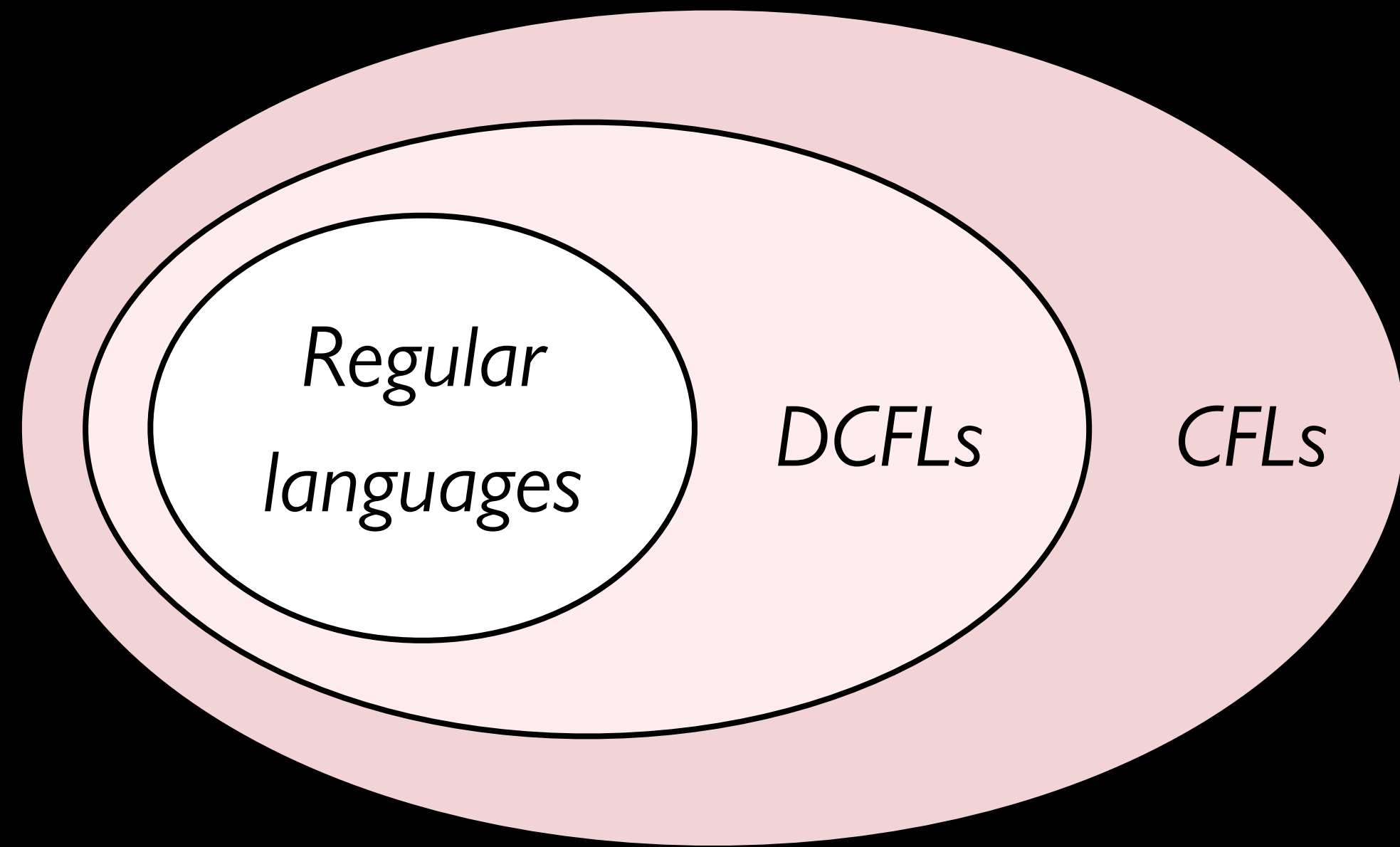




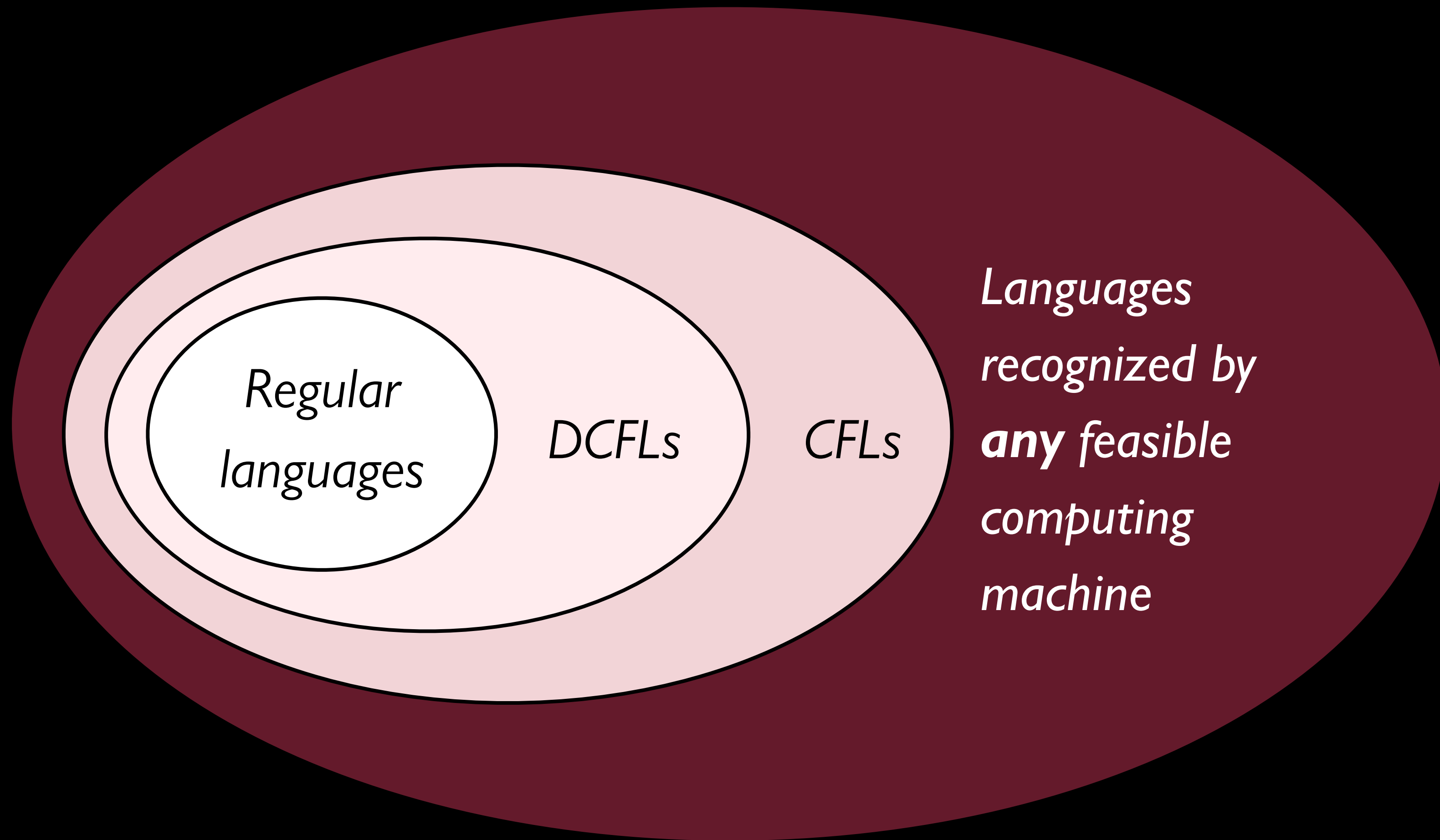
What problems can we solve with a computer?

*Finite automata* represent computers with bounded memory. They recognize the *regular languages*.

*Pushdown automata* represent computers with a limited form of unbounded memory. They recognize the *context-free languages*.



*All languages*



Regular  
languages

DCFLs

CFLs

Languages  
recognized by  
**any** feasible  
computing  
machine

All languages

To talk about what languages could be recognized by “any feasible computing machine”, we need our final, most powerful model of computation.



*Entscheidungsproblem* = Decision problem



# ON COMPUTABLE NUMBERS, WITH AN APPLICATION TO THE ENTSCHIEDUNGSPROBLEM

*By* A. M. TURING.

[Received 28 May, 1936.—Read 12 November, 1936.]

The “computable” numbers may be described briefly as the real numbers whose expressions as a decimal are calculable by finite means. Although the subject of this paper is ostensibly the computable *numbers*, it is almost equally easy to define and investigate computable functions of an integral variable or a real or computable variable, computable predicates, and so forth. The fundamental problems involved are, however, the same in each case, and I have chosen the computable numbers for explicit treatment as involving the least cumbrous technique. I hope shortly to give an account of the relations of the computable numbers, functions, and so forth to one another. This will include a development of the theory of functions of a real variable expressed in terms of computable numbers. According to my definition, a number is computable if its decimal can be written down by a machine.

ON COMPUTABLE NUMBERS, WITH AN APPLICATION TO  
THE ENTSCHIEDUNGSPROBLEM. **A CORRECTION**

*By* A. M. TURING.

In a paper entitled "On computable numbers, with an application to the Entscheidungsproblem"\* the author gave a proof of the insolubility of the Entscheidungsproblem of the "engere Funktionenkalkül". This proof contained some formal errors† which will be corrected here: there are also some other statements in the same paper which should be modified, although they are not actually false as they stand.

The expression for  $\text{Inst} \{q_i S_j S_k Lq_l\}$  on p. 260 of the paper quoted should read

$$(x, y, x', y') \left\{ \left( R_{S_j}(x, y) \& I(x, y) \& K_{q_i}(x) \& F(x, x') \& F(y', y) \right) \right. \\ \rightarrow \left( I(x', y') \& R_{S_k}(x', y) \& K_{q_l}(x') \& F(y', z) \vee \left[ \left( R_{S_0}(x, z) \rightarrow R_{S_0}(x', z) \right) \right. \right. \\ \left. \left. \& \left( R_{S_1}(x, z) \rightarrow R_{S_1}(x', z) \right) \& \dots \& \left( R_{S_n}(x, z) \rightarrow R_{S_n}(x', z) \right) \right] \right) \left. \right\},$$

At the same time, *Alonzo Church* published similar ideas and results – the  $\lambda$  calculus.

However, the Turing model has become the standard in theoretical computer science



$$\begin{array}{r} 27182818284590 \\ +31415926535897 \\ \hline \end{array}$$

$$\begin{array}{r} 27182818284590 \\ +31415926535897 \\ \hline 7 \end{array}$$

$$\begin{array}{r} 27182818284590 \\ +31415926535897 \\ \hline 87 \end{array}$$

$$\begin{array}{r} \phantom{2718281828}4\mathbf{5}90 \\ +3141592653\mathbf{5}897 \\ \hline \phantom{2718281828}4\mathbf{8}7 \end{array}$$

$$\begin{array}{r} \phantom{2718281828}111 \\ 27182818284590 \\ +31415926535897 \\ \hline \phantom{2718281828}0487 \end{array}$$

$$\begin{array}{r} \phantom{271828182} \phantom{84590} \phantom{+} \phantom{314159265} \phantom{35897} \\ \phantom{271828182} \phantom{84590} \phantom{+} \phantom{314159265} \phantom{35897} \phantom{1} \phantom{1} \phantom{1} \phantom{1} \\ 271828182\mathbf{8}4590 \\ +314159265\mathbf{3}5897 \\ \hline \phantom{271828182} \phantom{84590} \phantom{+} \phantom{314159265} \phantom{35897} \phantom{1} \phantom{1} \phantom{1} \phantom{1} \\ \phantom{271828182} \phantom{84590} \phantom{+} \phantom{314159265} \phantom{35897} \mathbf{20487} \end{array}$$

$$\begin{array}{r} \phantom{27182818} 1\ 1\ 1\ 1 \\ 27182818\mathbf{2}84590 \\ +31415926\mathbf{5}35897 \\ \hline \phantom{27182818} \phantom{+31415926} \mathbf{8}20487 \end{array}$$

$$\begin{array}{r} \phantom{2718281}1\phantom{284590} \\ \phantom{2718281}1\phantom{284590} \\ \phantom{2718281}1\phantom{284590} \\ \phantom{2718281}1\phantom{284590} \\ \phantom{2718281}1\phantom{284590} \\ 2718281\mathbf{8}284590 \\ +3141592\mathbf{6}535897 \\ \hline 4820487 \end{array}$$

$$\begin{array}{r} \phantom{271828}1\phantom{8284590} \\ \phantom{271828}1\phantom{8284590} \\ \phantom{271828}1\phantom{8284590} \\ \phantom{271828}1\phantom{8284590} \\ \phantom{271828}1\phantom{8284590} \\ 27182818284590 \\ +31415926535897 \\ \hline 44820487 \end{array}$$

$$\begin{array}{r} \phantom{2} \phantom{7} \phantom{1} \phantom{8} \phantom{2} \phantom{8} \phantom{1} \phantom{8} \phantom{2} \phantom{8} \phantom{4} \phantom{5} \phantom{9} \phantom{0} \\ \phantom{2} \phantom{7} \phantom{1} \phantom{8} \phantom{2} \phantom{8} \phantom{1} \phantom{8} \phantom{2} \phantom{8} \phantom{4} \phantom{5} \phantom{9} \phantom{0} \\ + 3 \phantom{1} \phantom{4} \phantom{1} \phantom{5} \phantom{9} \phantom{2} \phantom{6} \phantom{5} \phantom{3} \phantom{5} \phantom{8} \phantom{9} \phantom{7} \\ \hline 7 \phantom{4} \phantom{4} \phantom{8} \phantom{2} \phantom{0} \phantom{4} \phantom{8} \phantom{7} \end{array}$$

$$\begin{array}{r} \phantom{2718}2818284590 \\ +31415926535897 \\ \hline 8744820487 \end{array}$$

1 1 1 1 1 1

$$\begin{array}{r}
 \phantom{+}271\mathbf{8}2818284590 \\
 +314\mathbf{1}5926535897 \\
 \hline
 98744820487
 \end{array}$$

The image shows a vertical addition problem with carry-over indicators. The numbers are aligned to the right, with the first number having eight digits and the second having nine. The result has nine digits. The carry-over indicators are '1' placed above the second, third, fourth, fifth, sixth, and seventh digits of the top number. The bolded digits represent the carry-over points.

$$\begin{array}{r} \phantom{+}27182818284590 \\ +31415926535897 \\ \hline 598744820487 \end{array}$$

1 1 1 1 1 1



$$\begin{array}{r} \phantom{+}27182818284590 \\ +31415926535897 \\ \hline 58598744820487 \end{array}$$

1 1 1 1 1 1

$$\begin{array}{r}
 \phantom{+}27182818284590 \\
 +31415926535897 \\
 \hline
 58598744820487
 \end{array}$$

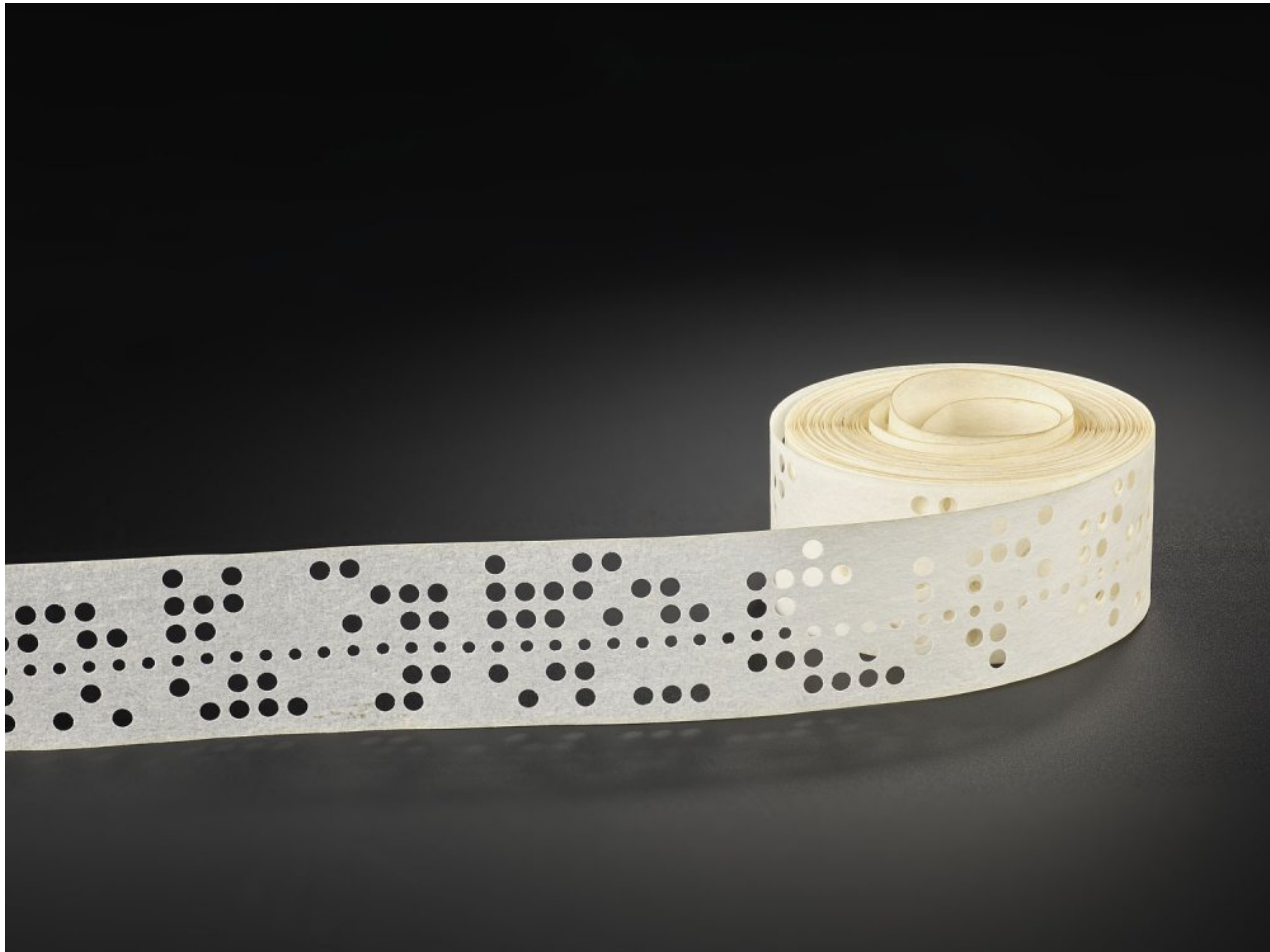
*Key idea:*

You might need huge amounts of scratch space to perform a calculation...

But at each point in the calculation you only need access to a small amount of that scratch space!

To provide his machines extra memory, Turing gave them access to an *infinite tape*, subdivided into *tape cells*.

A Turing machine can only see one tape cell at a time, the one pointed at by the *tape head*.

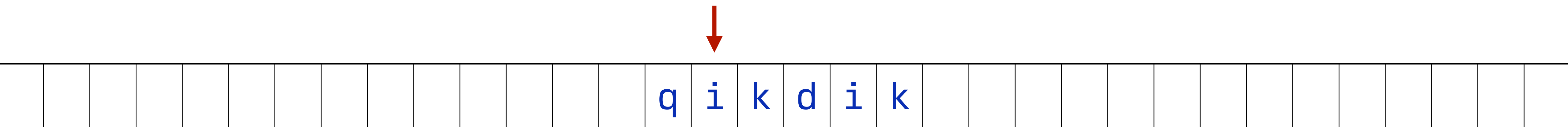


*Punched paper computer tape (1955–65)*



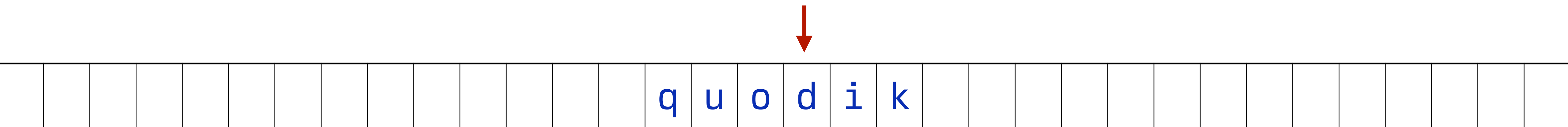


At every step, the Turing machine can read the cell under the tape head, change which symbol was written under the tape head, and move its tape head to the left or to the right.





At every step, the Turing machine can  
read the cell under the tape head,  
change which symbol was written under the tape head, and  
move its tape head to the left or to the right.





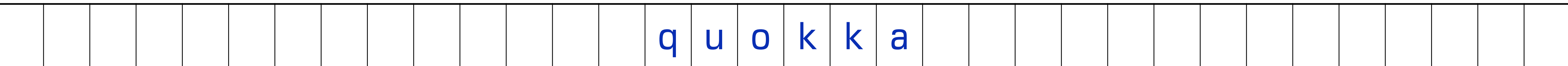




At every step, the Turing machine can read the cell under the tape head, change which symbol was written under the tape head, and move its tape head to the left or to the right.



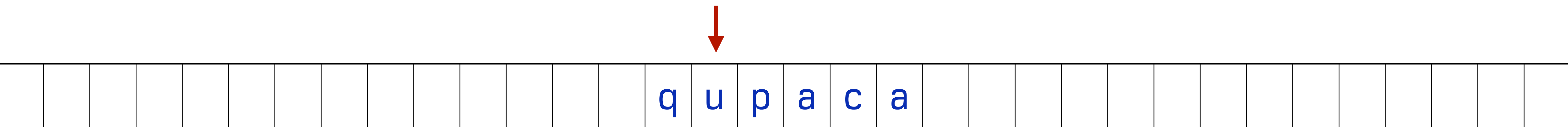
*A quokka*





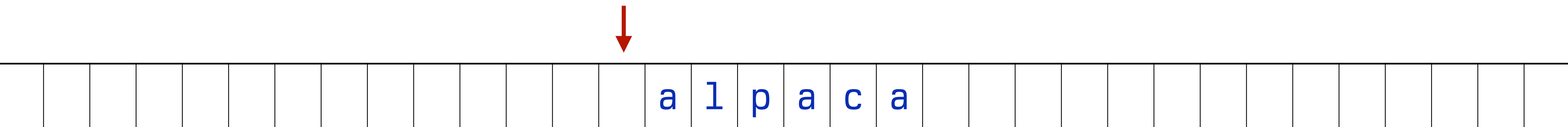


At every step, the Turing machine can read the cell under the tape head, change which symbol was written under the tape head, and move its tape head to the left or to the right.

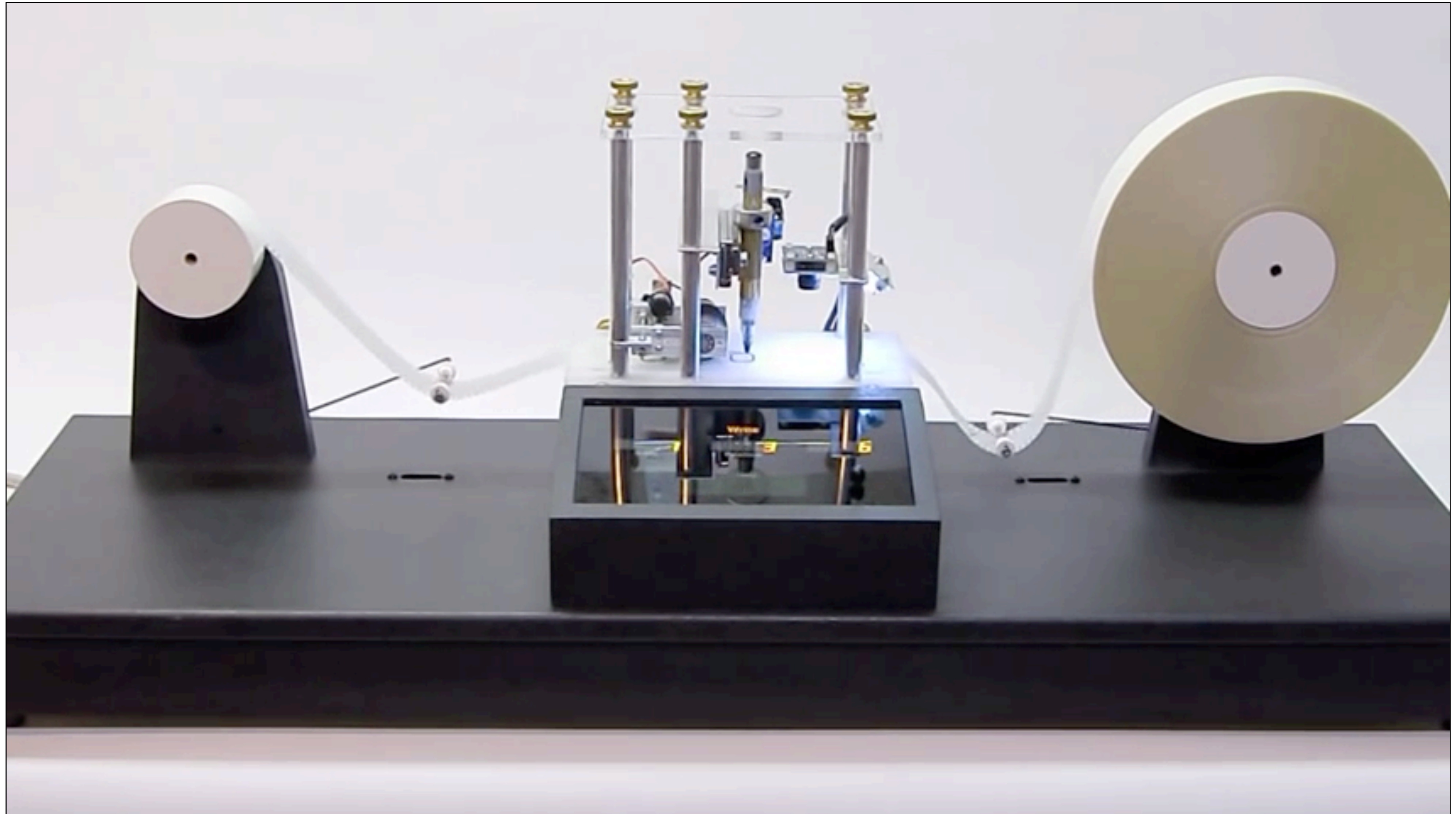




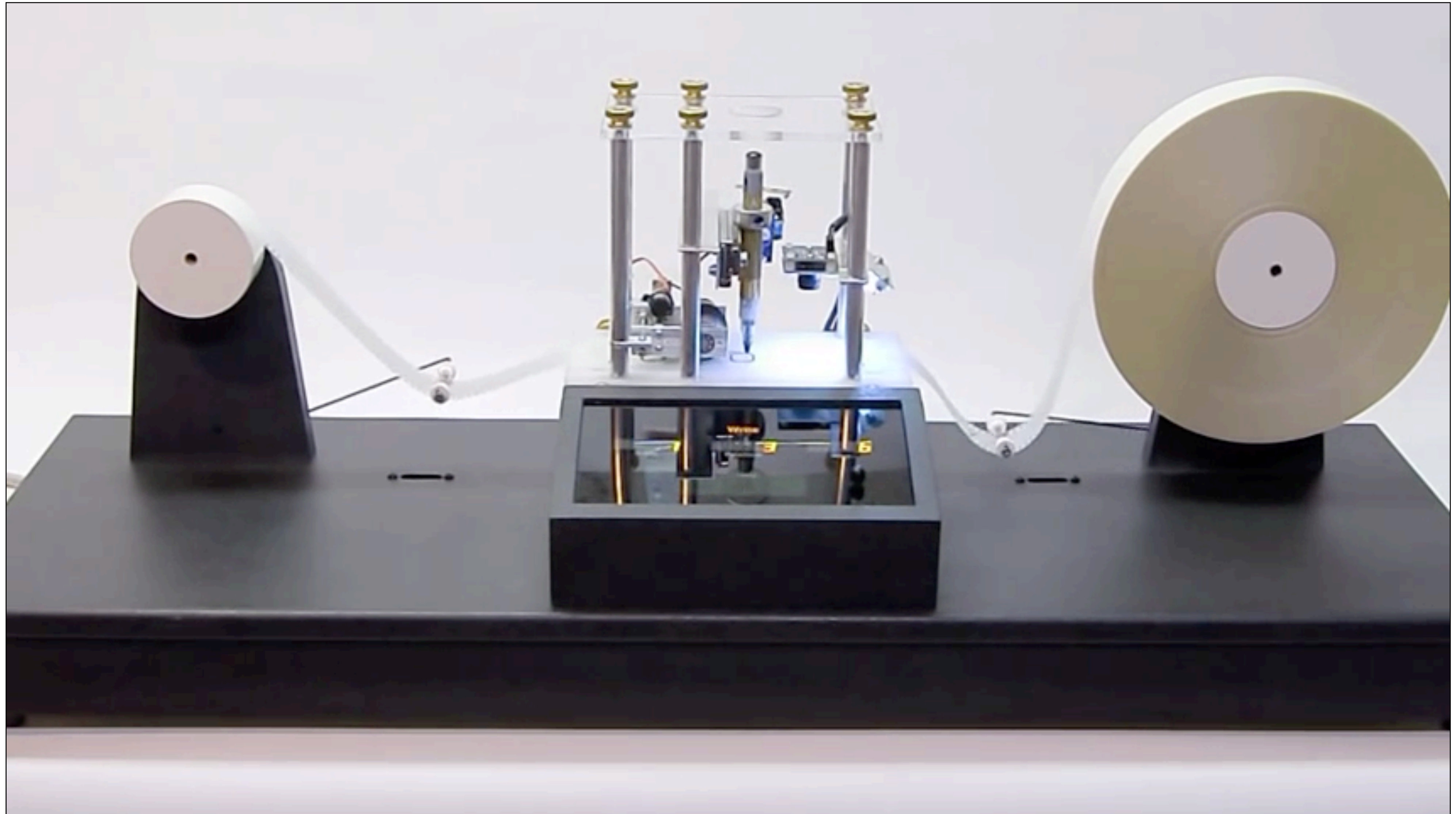
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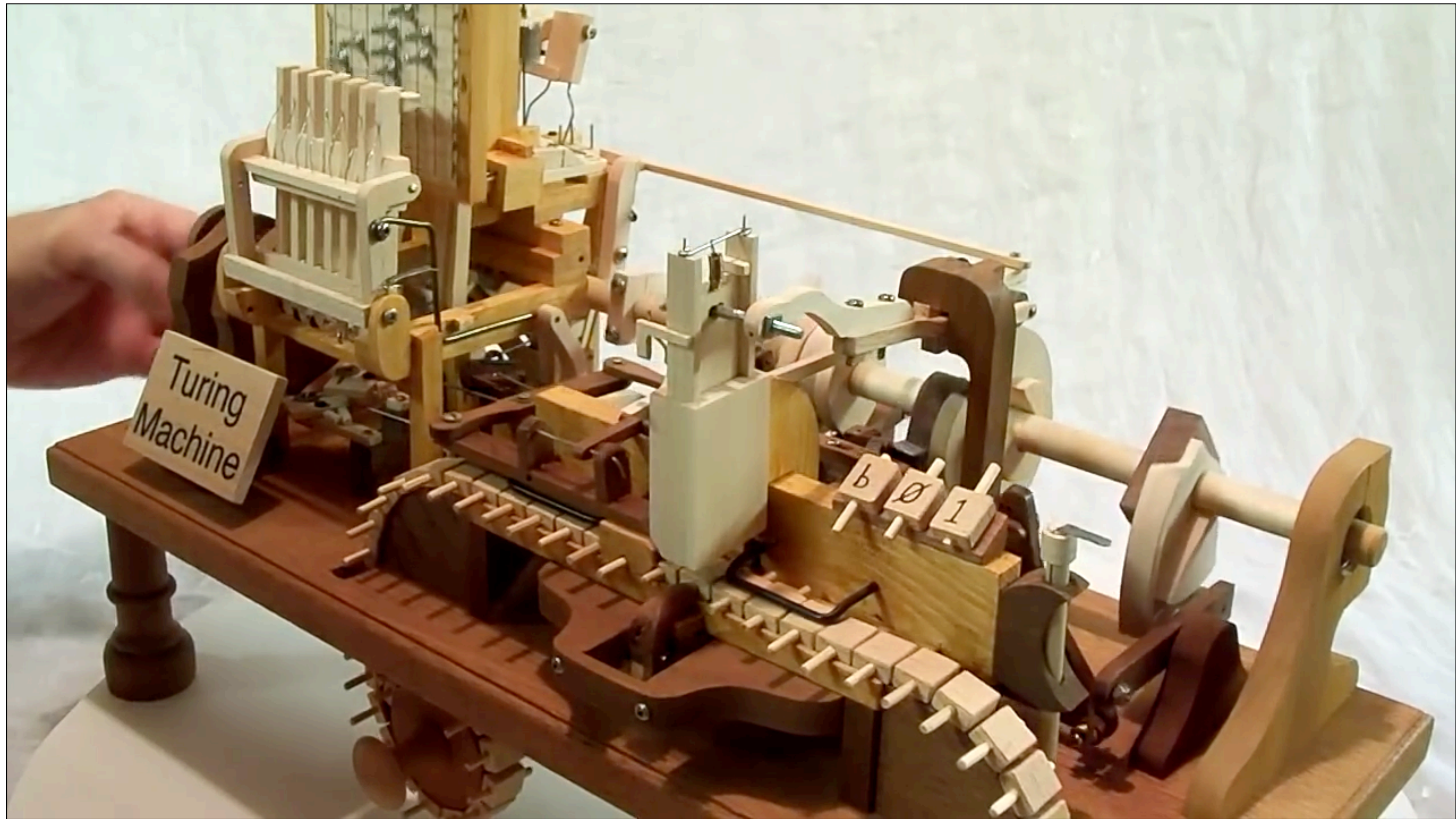




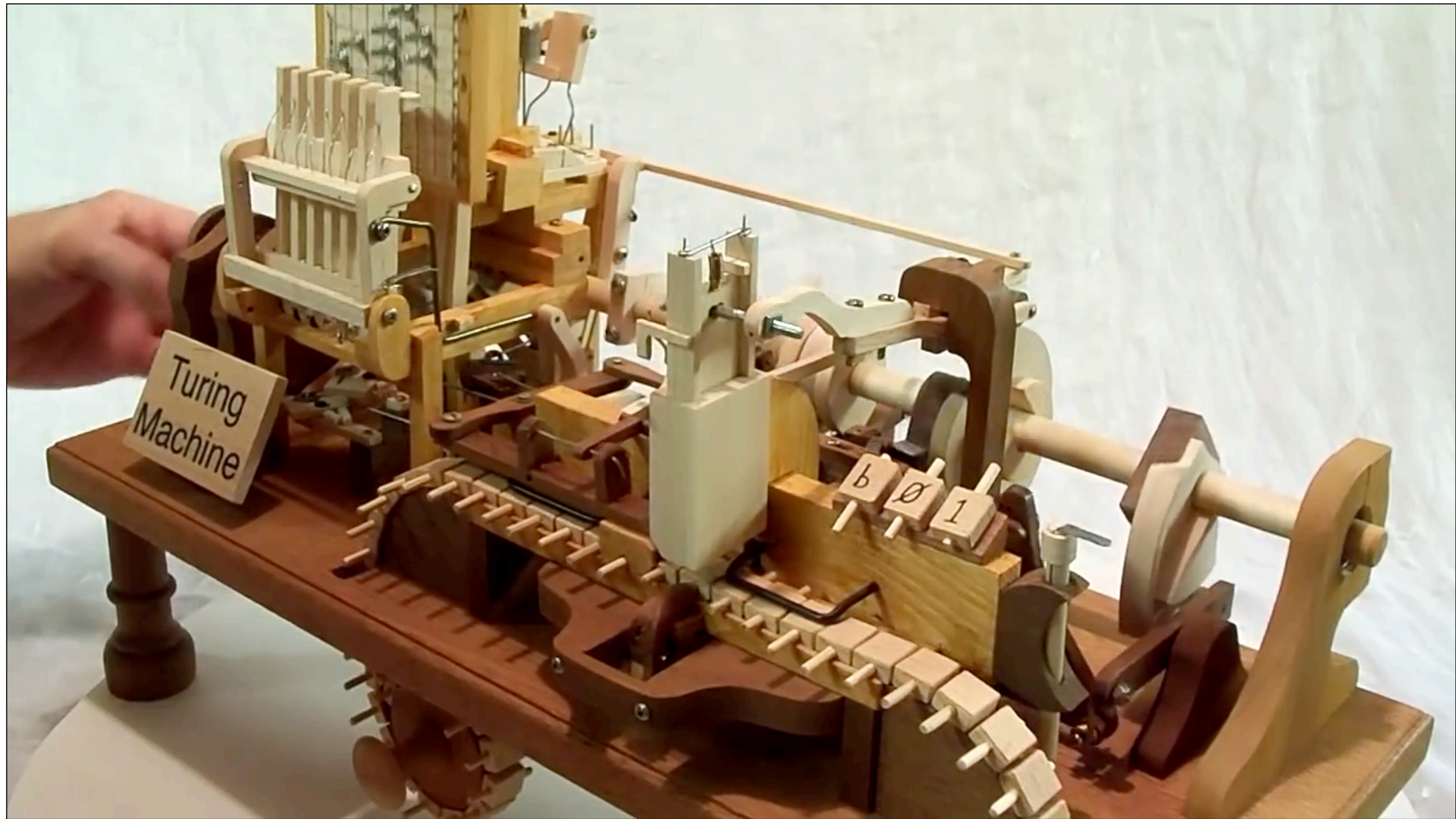
[youtu.be/E3keLeMwfHY](https://youtu.be/E3keLeMwfHY)



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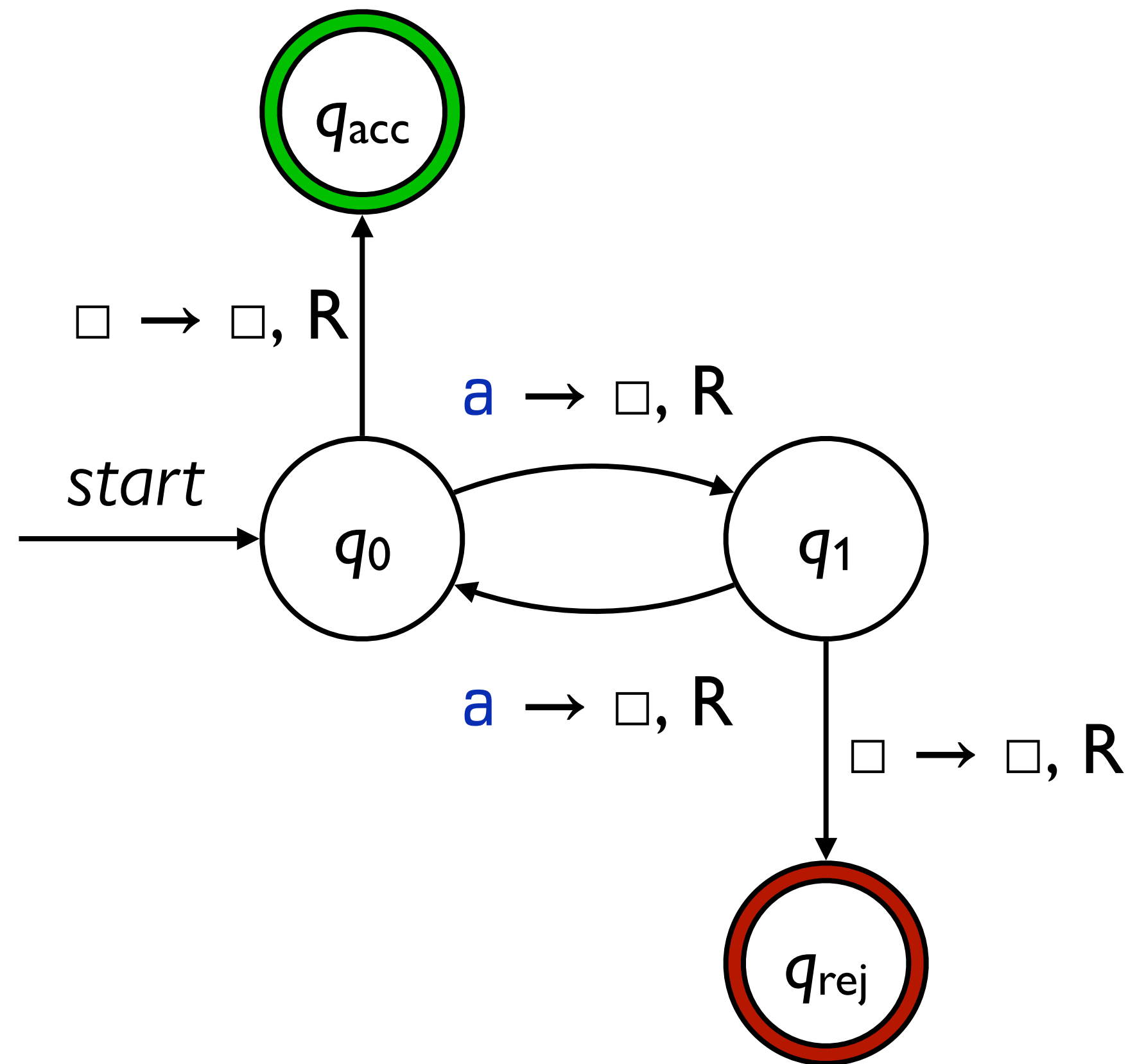


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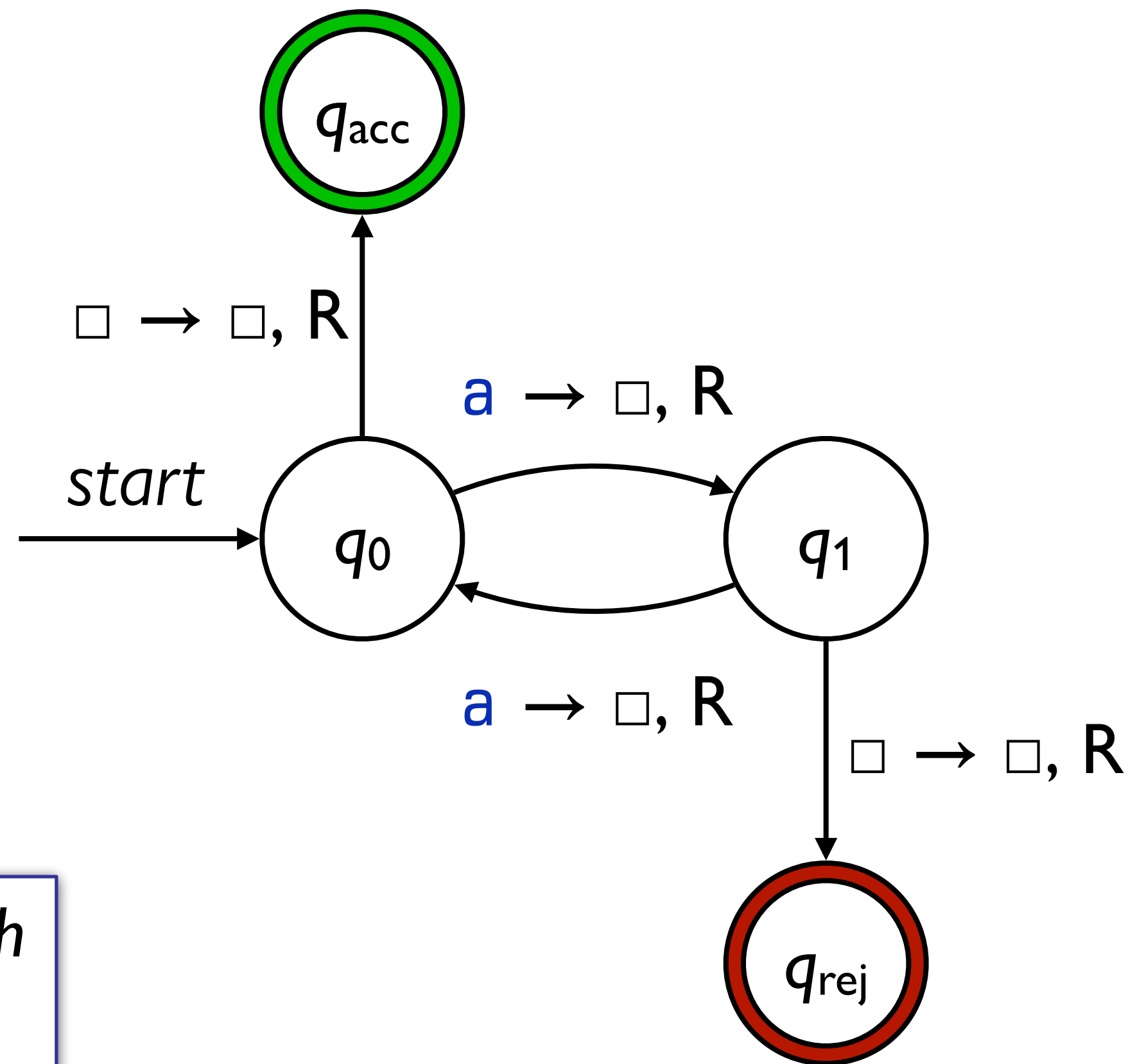


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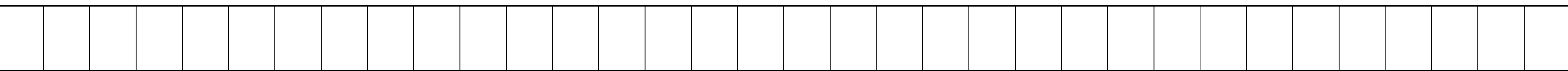


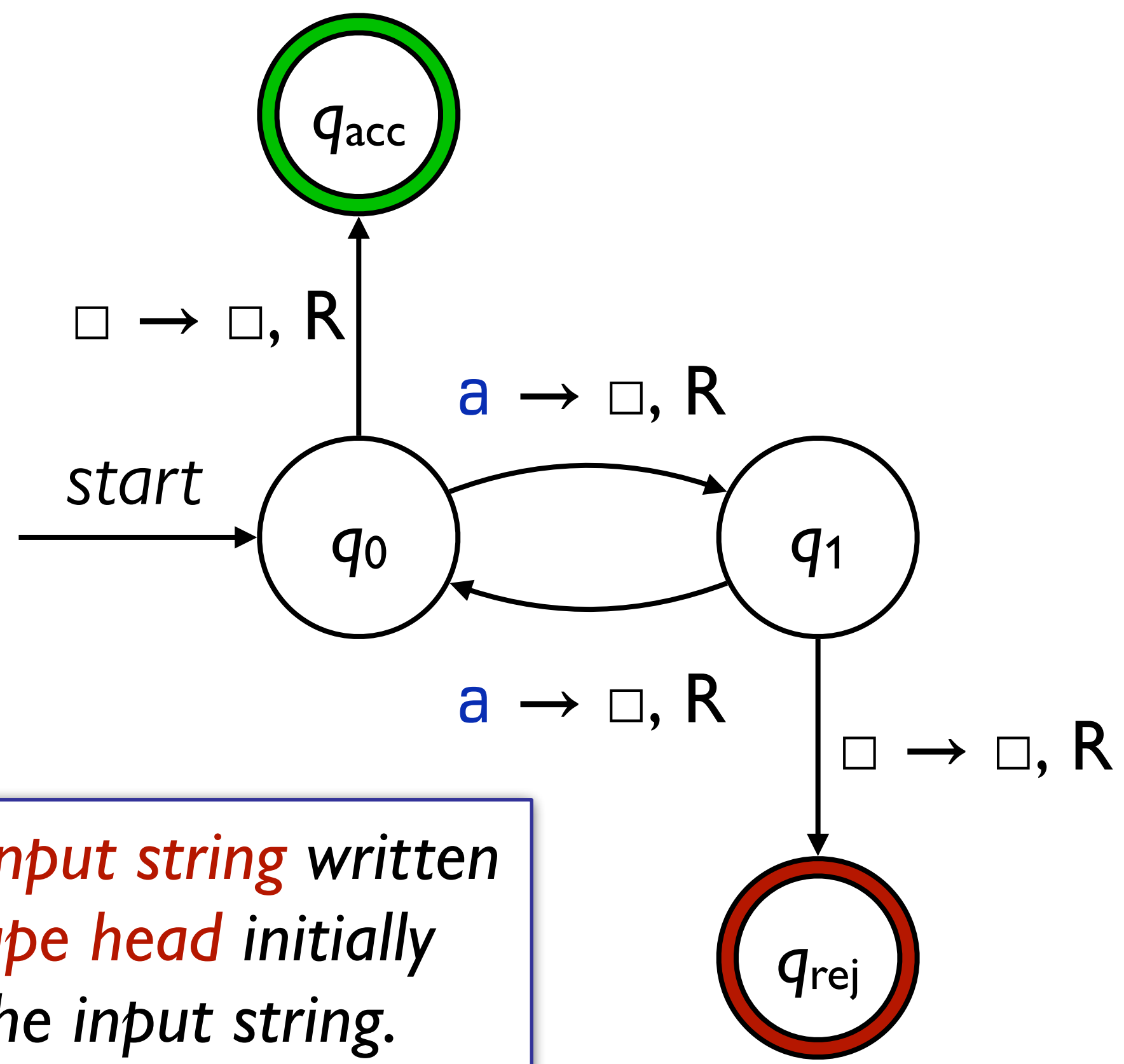


*This is the Turing machine's **finite-state control**. It issues commands that drive the operation of the machine.*

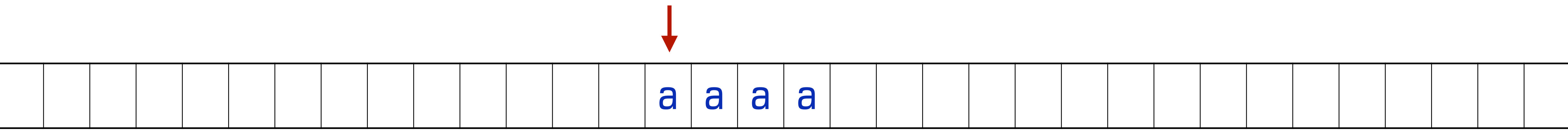


This is the TM's *infinite tape*. Each tape cell holds a *tape symbol*. Initially all tape symbols are blank.

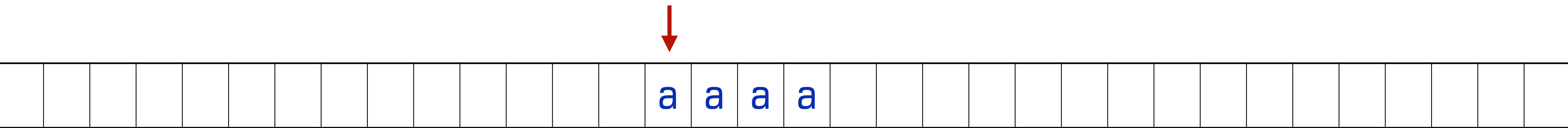
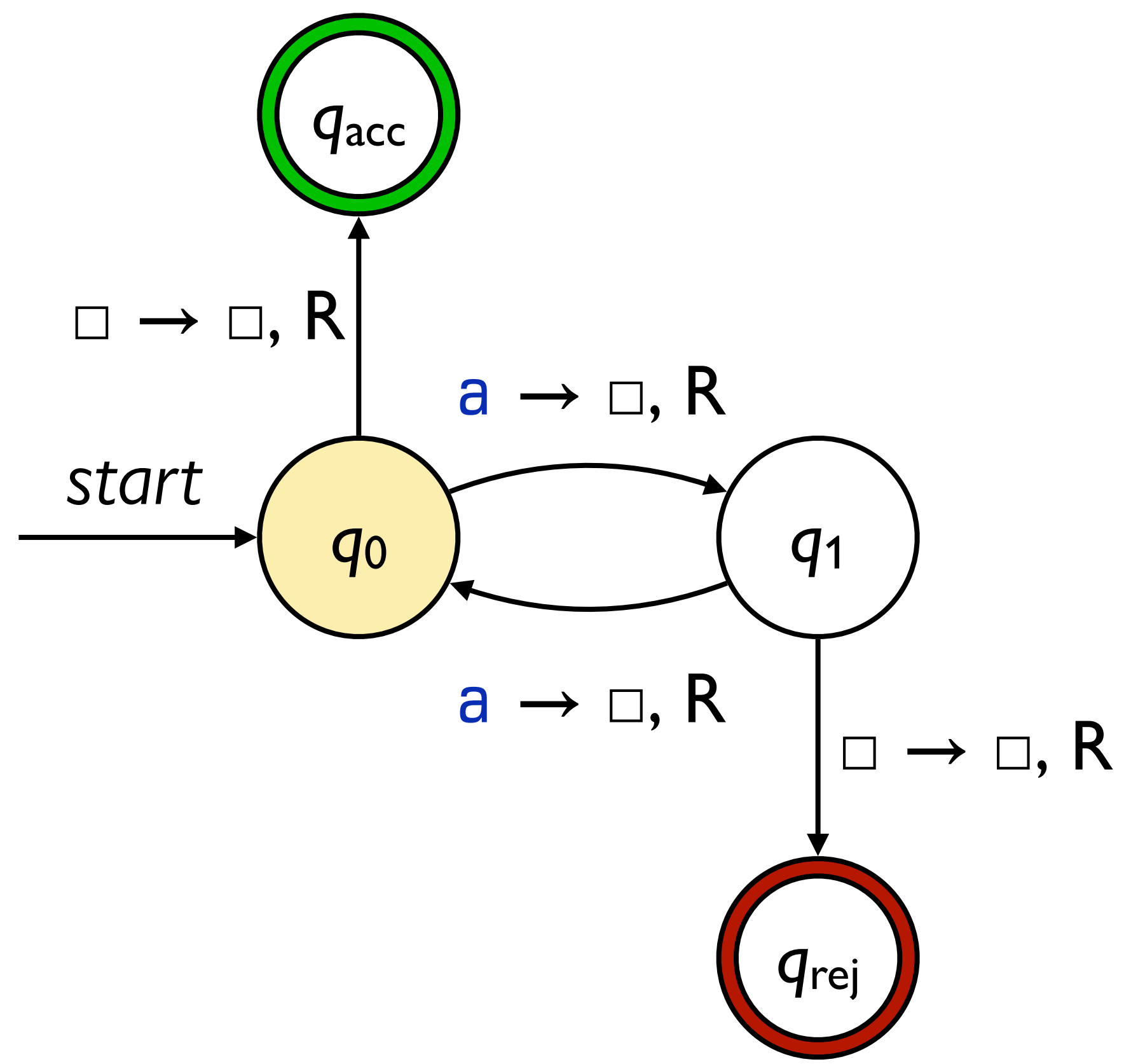


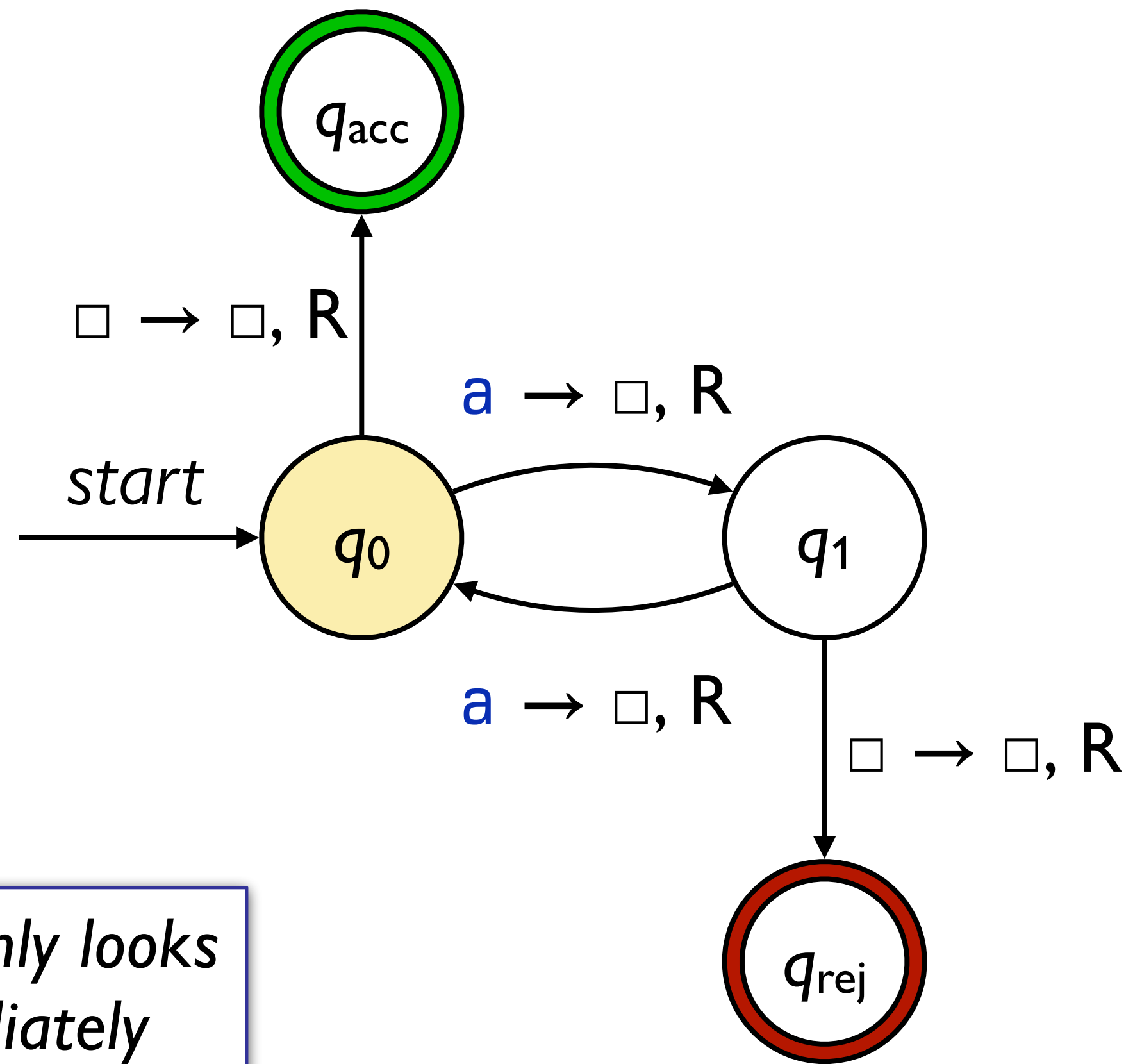


The machine is started with the *input string* written somewhere on the *tape*. The *tape head* initially points to the first symbol of the input string.

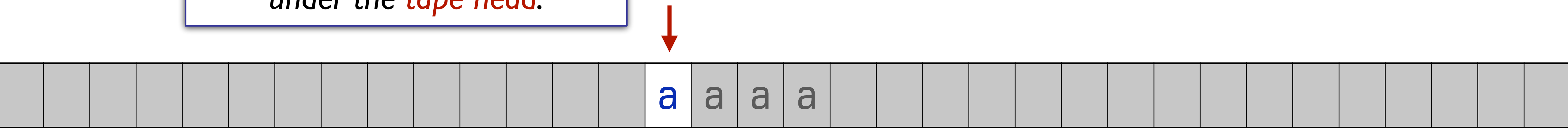


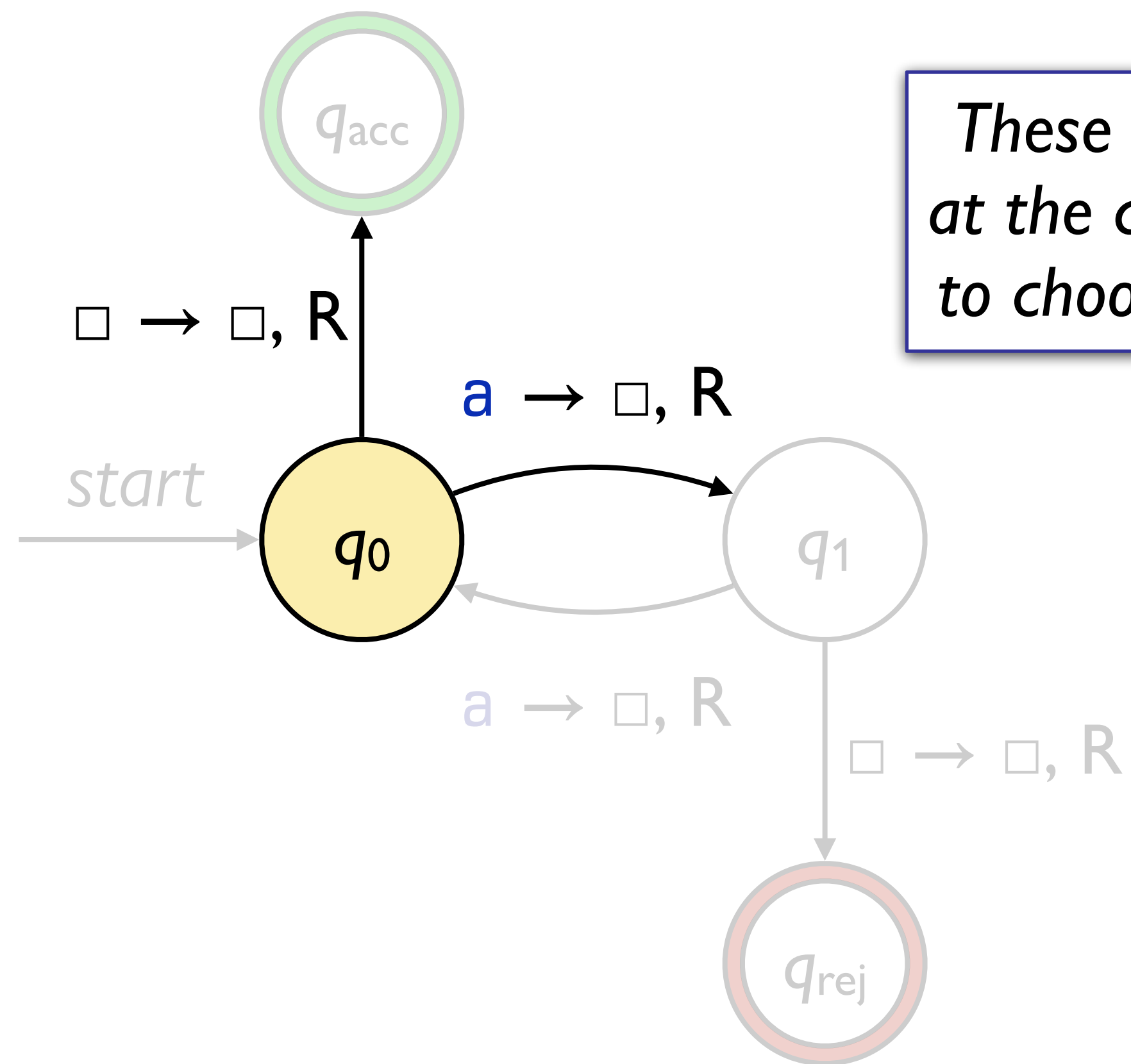
Like other automata, TMs begin execution in their *start state*



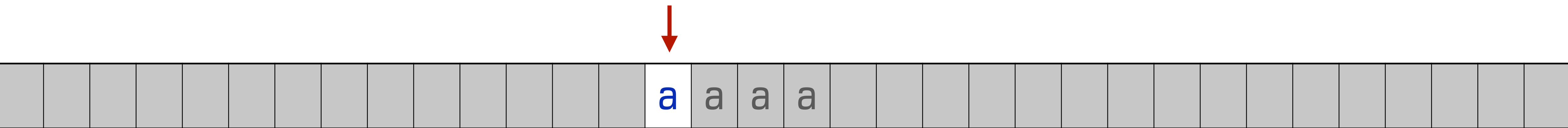


At each step, the TM only looks at the symbol immediately under the *tape head*.

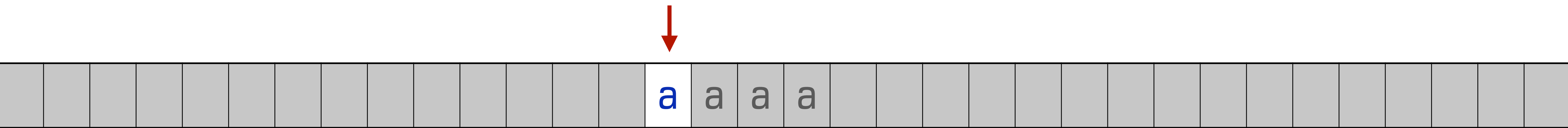
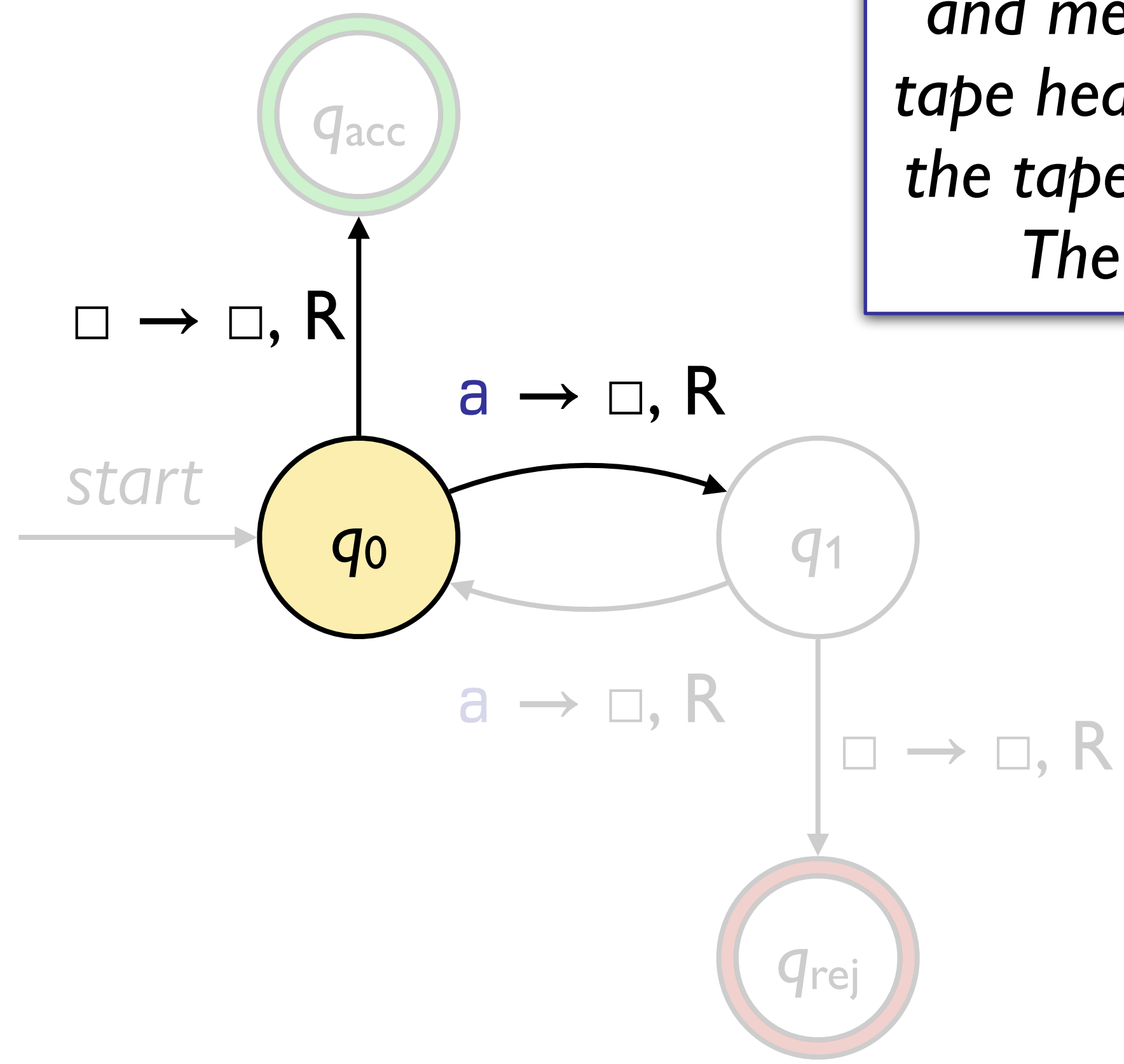




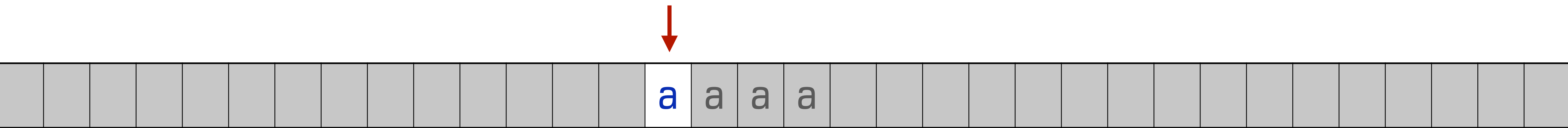
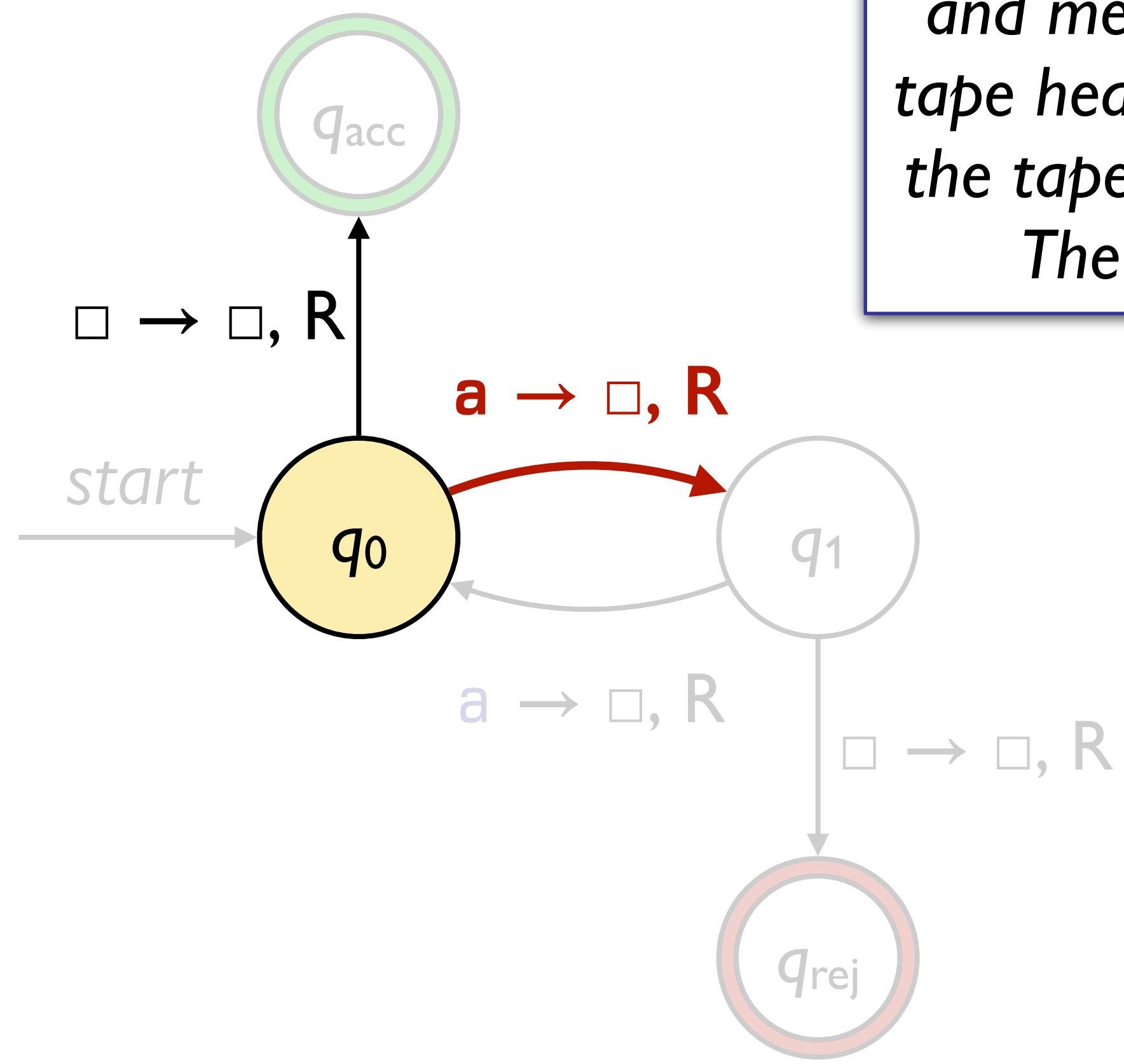
*These two transitions originate at the current state. We're going to choose one of them to follow.*



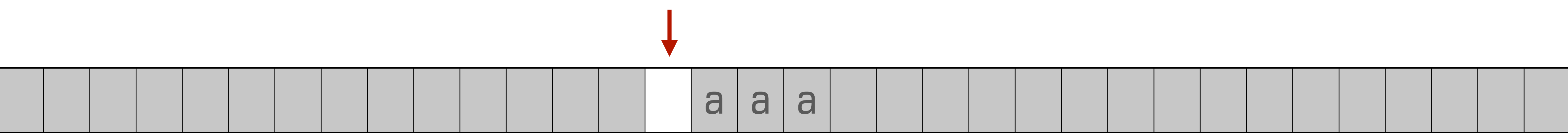
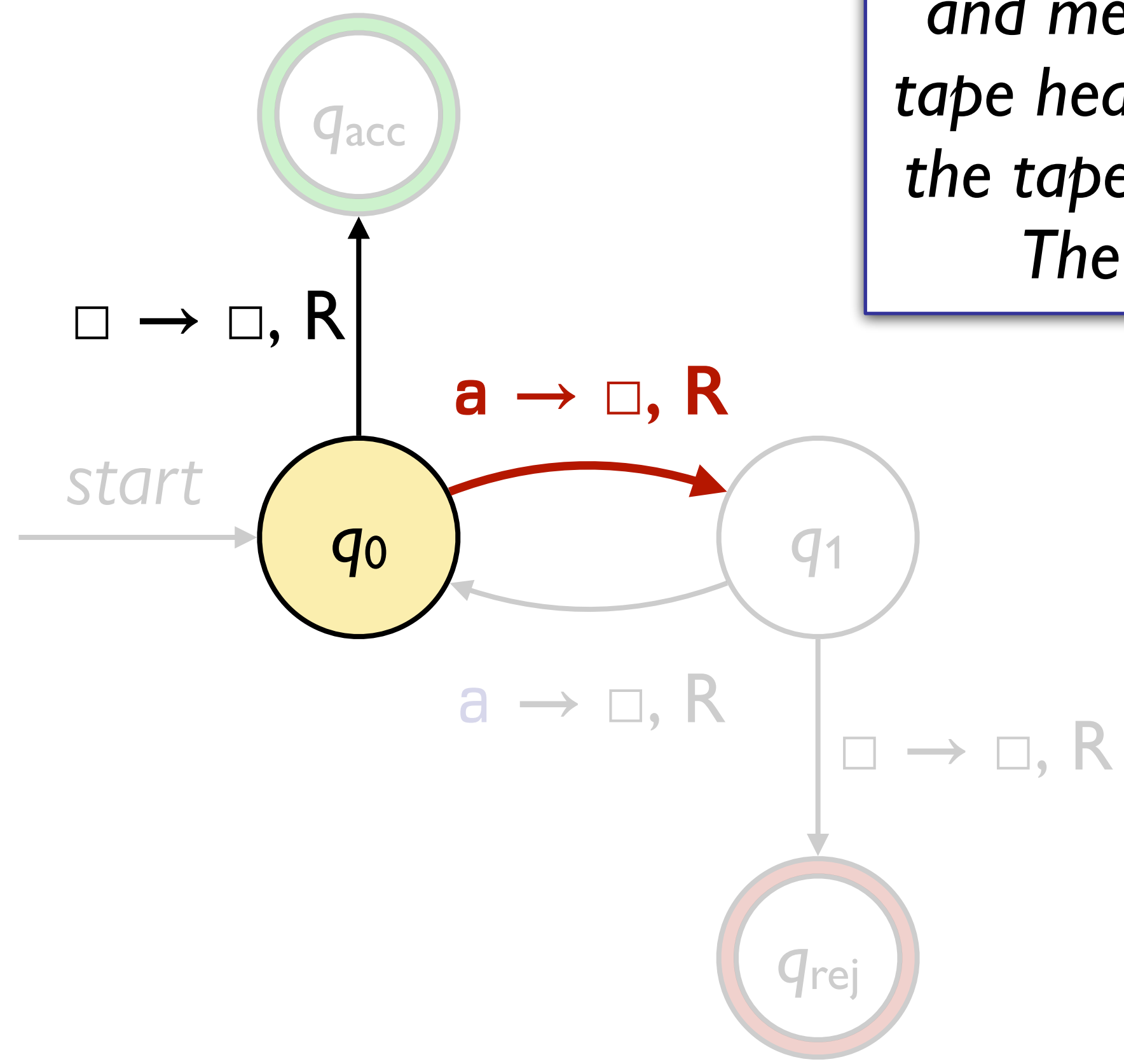
Each transition has the form  $\langle \text{read} \rangle \rightarrow \langle \text{write} \rangle, \langle \text{direction} \rangle$  and means “if symbol  $\langle \text{read} \rangle$  is under the tape head, replace it with  $\langle \text{write} \rangle$  and move the tape head in  $\langle \text{direction} \rangle$  (left or right)”. The  $\square$  symbol denotes a blank cell.



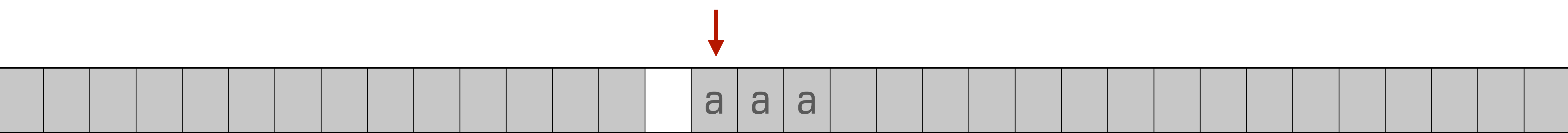
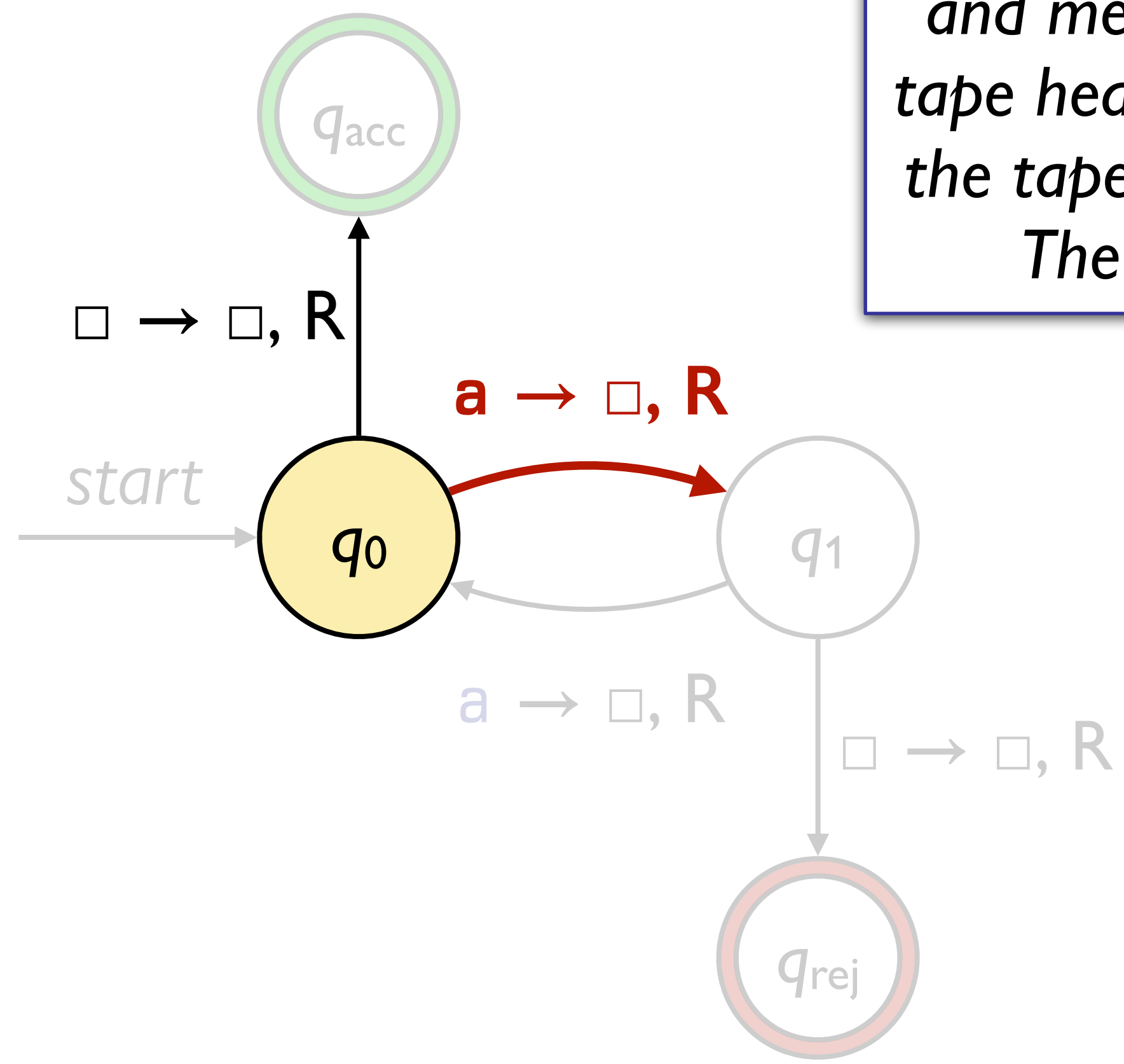
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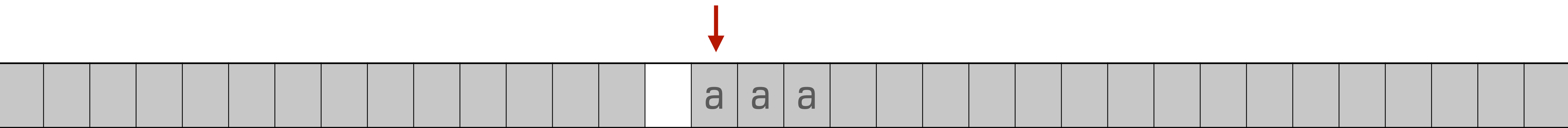
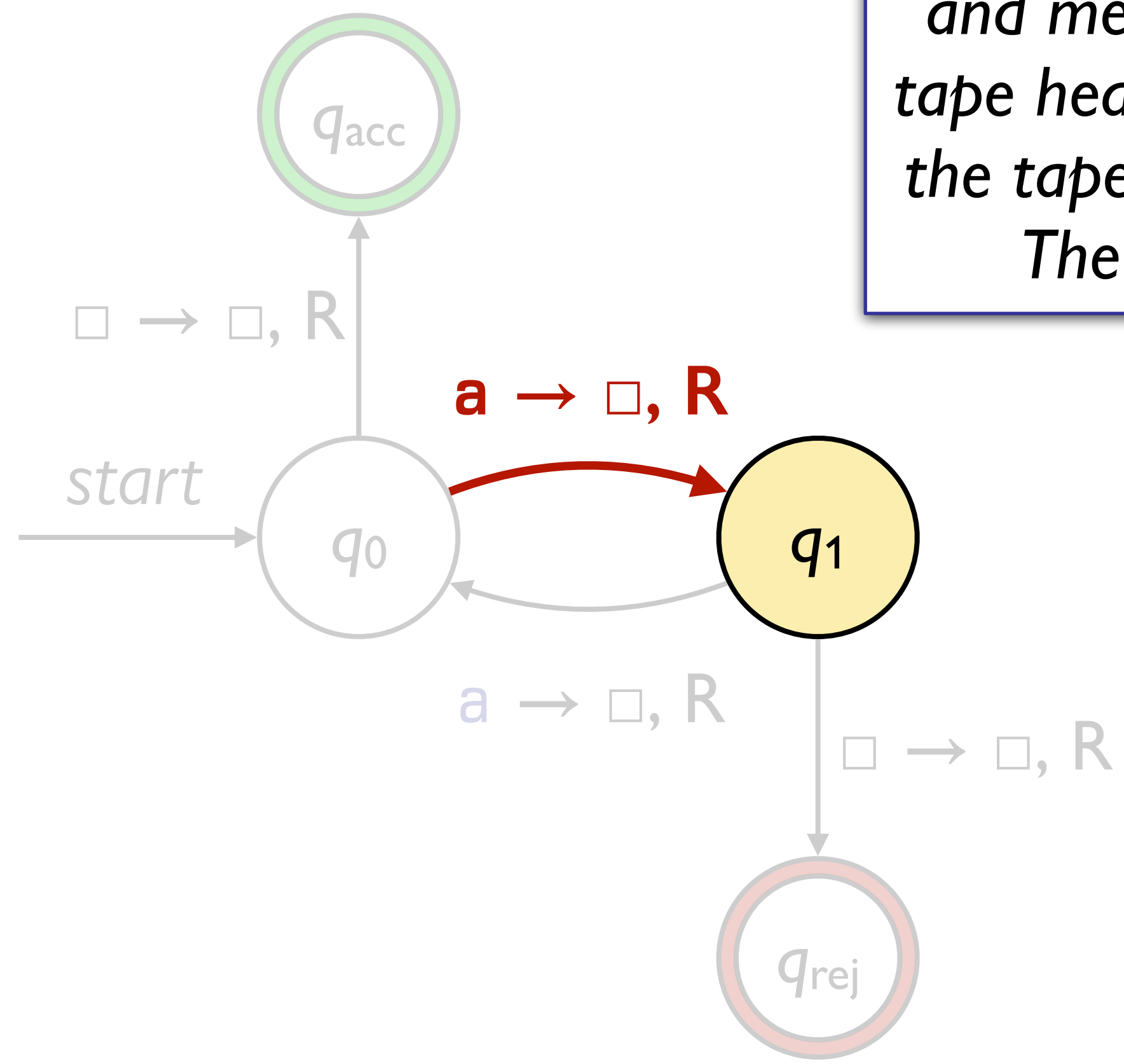
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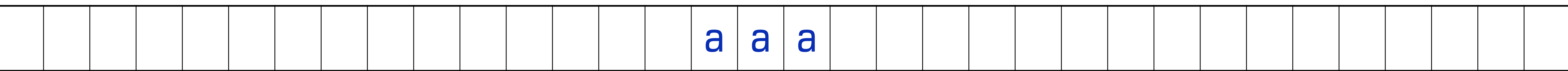
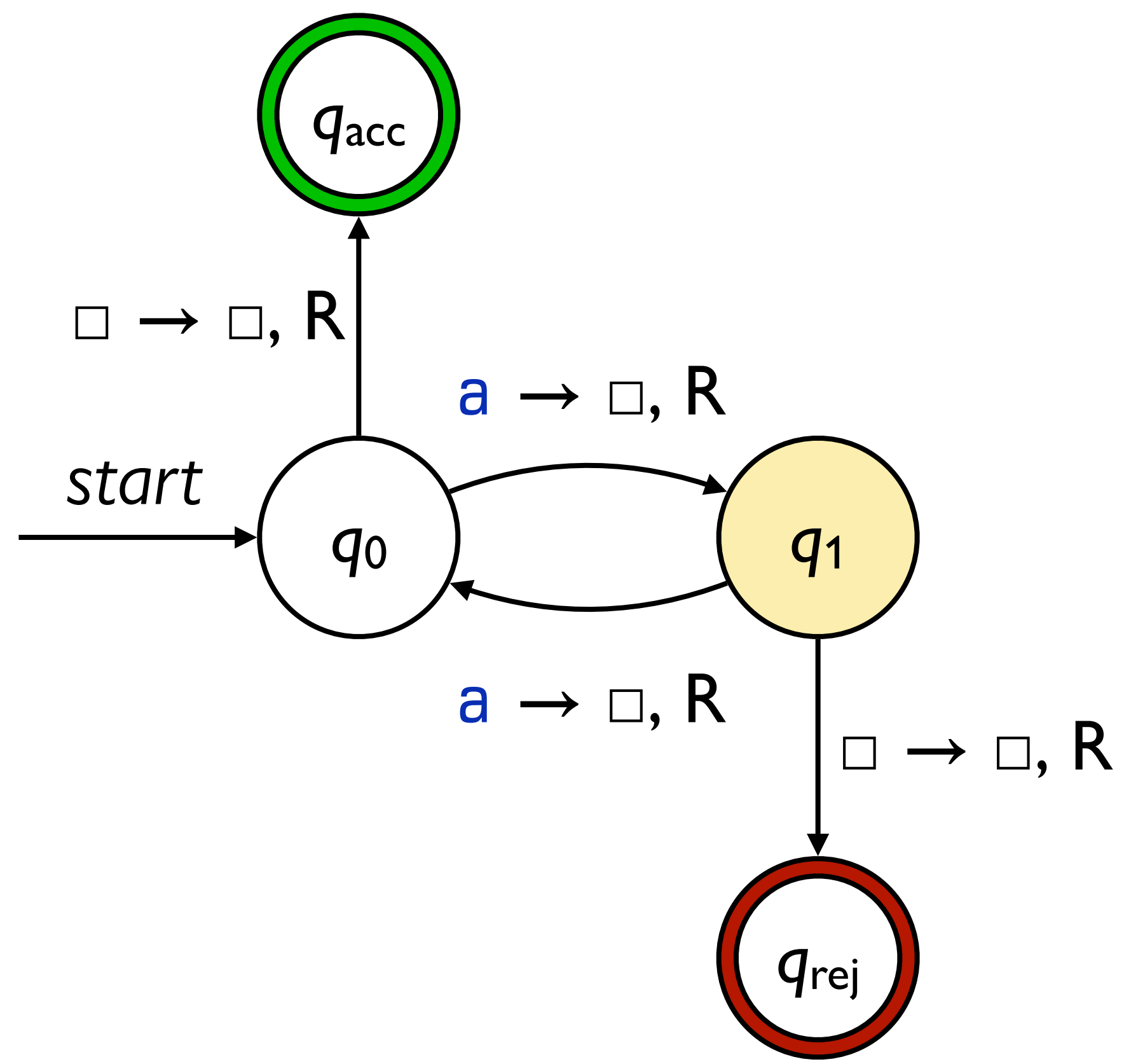


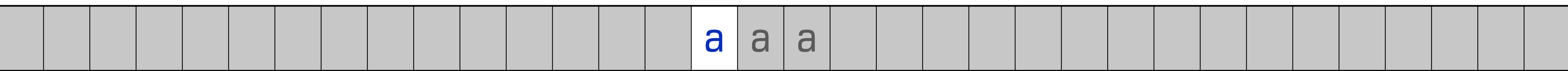
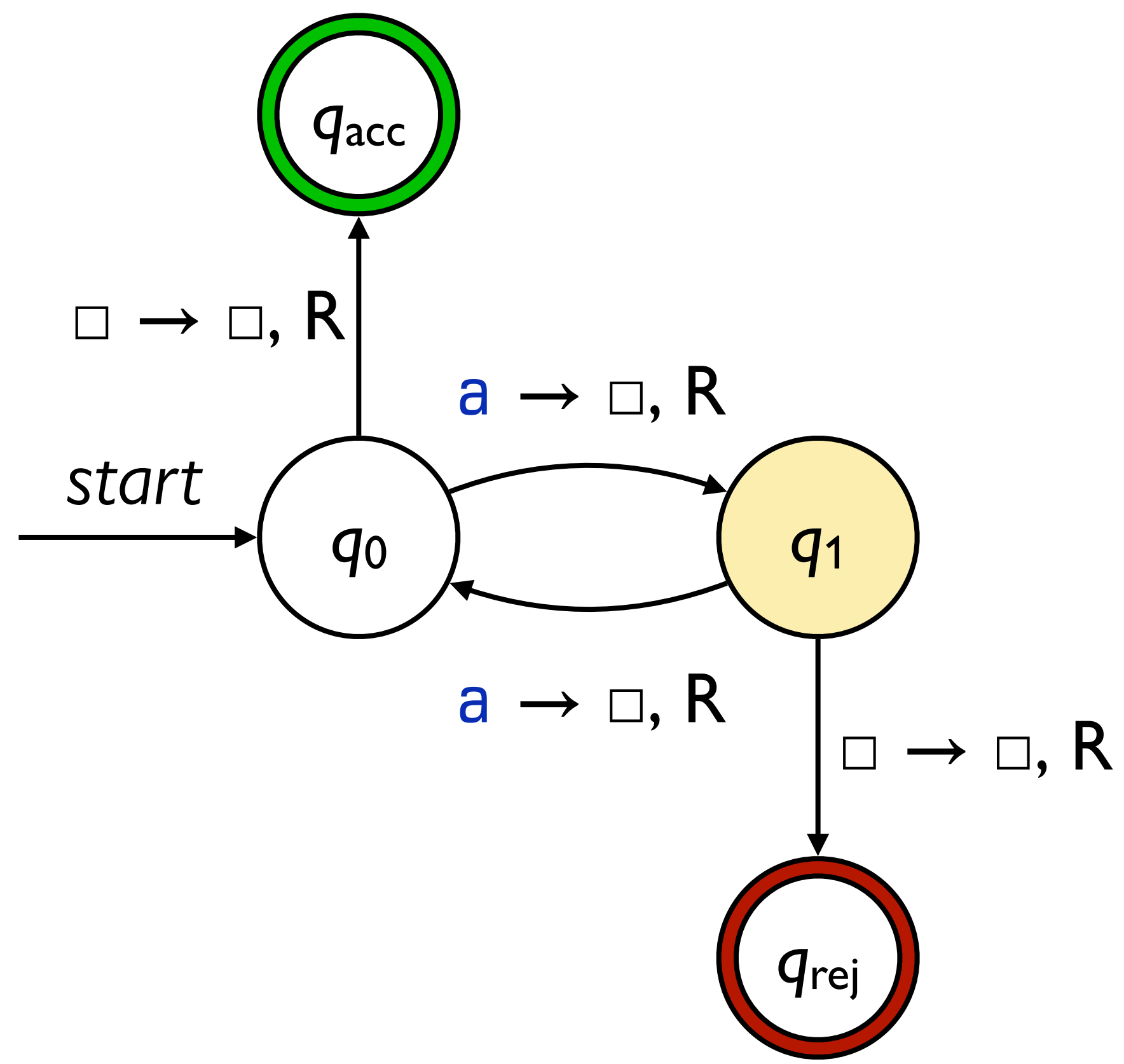
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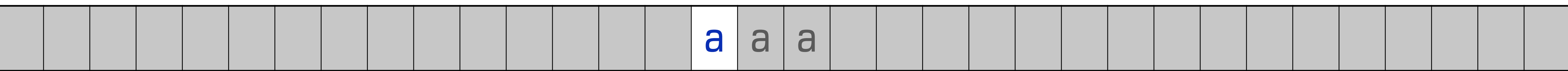
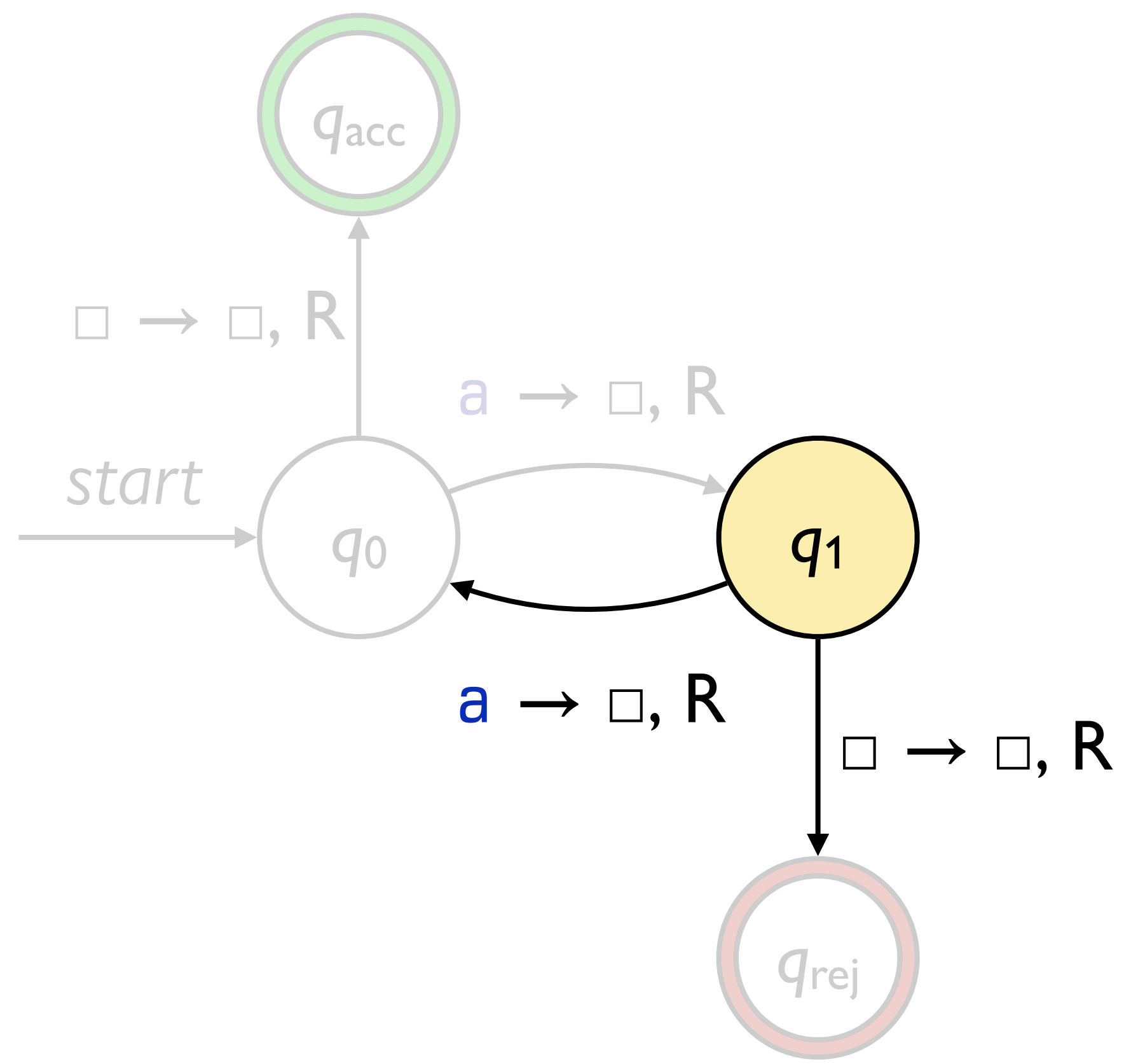


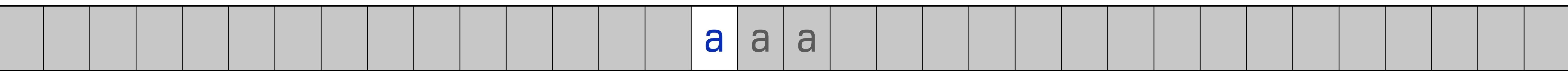
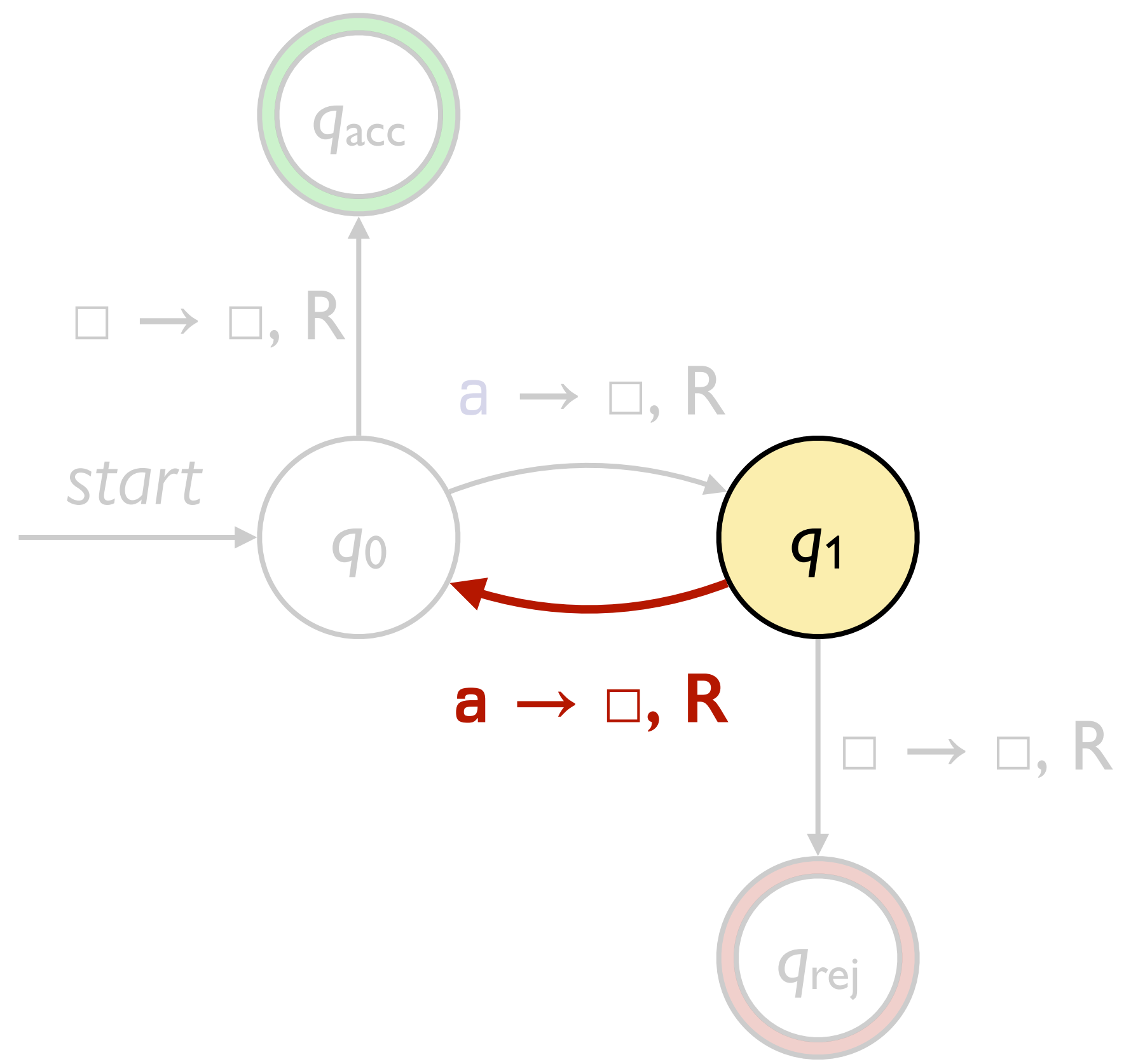
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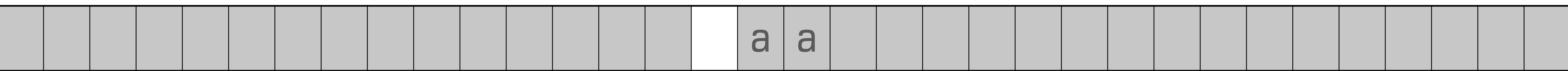
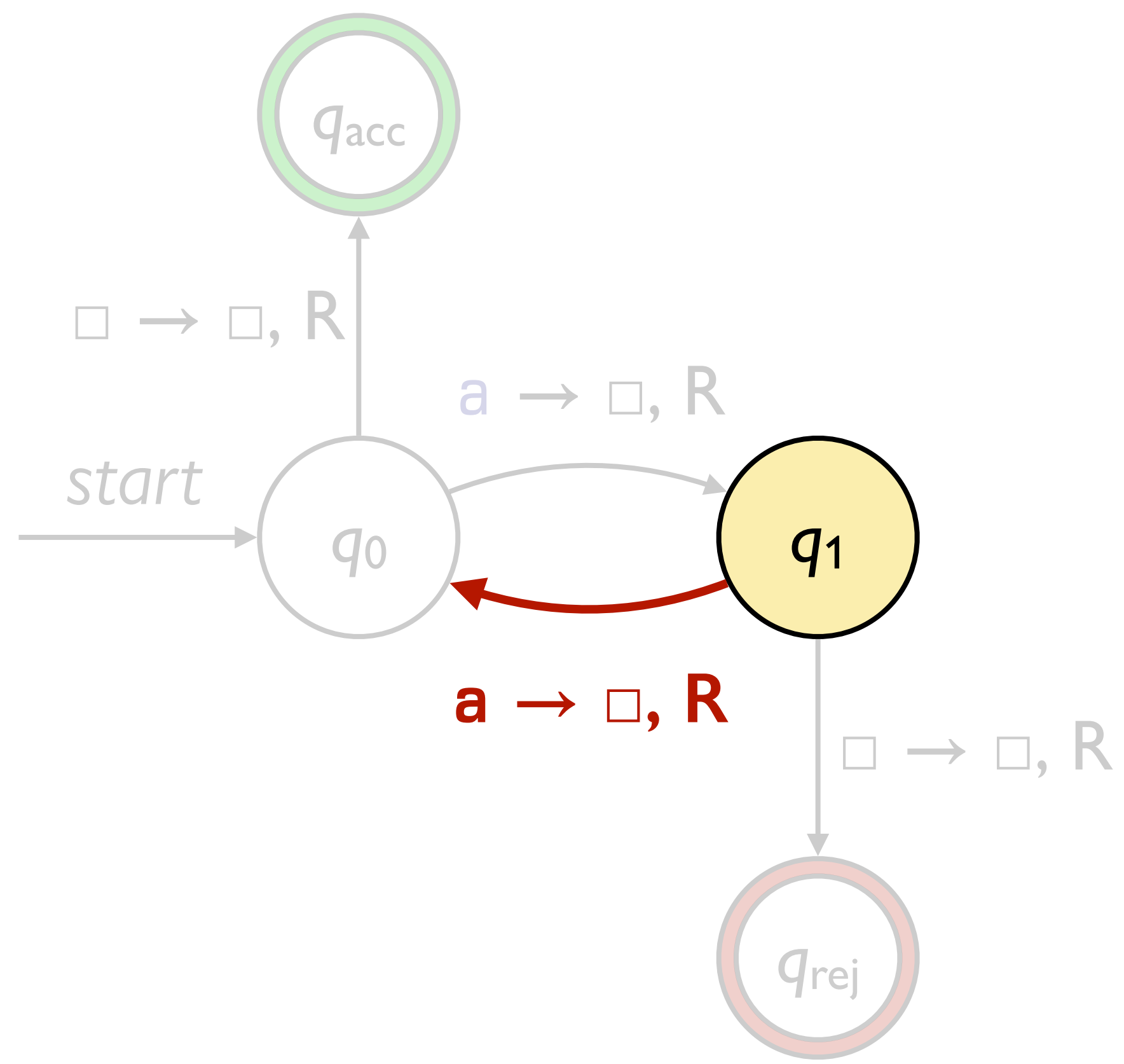


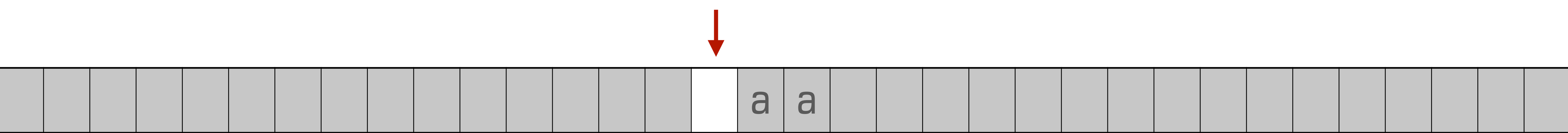
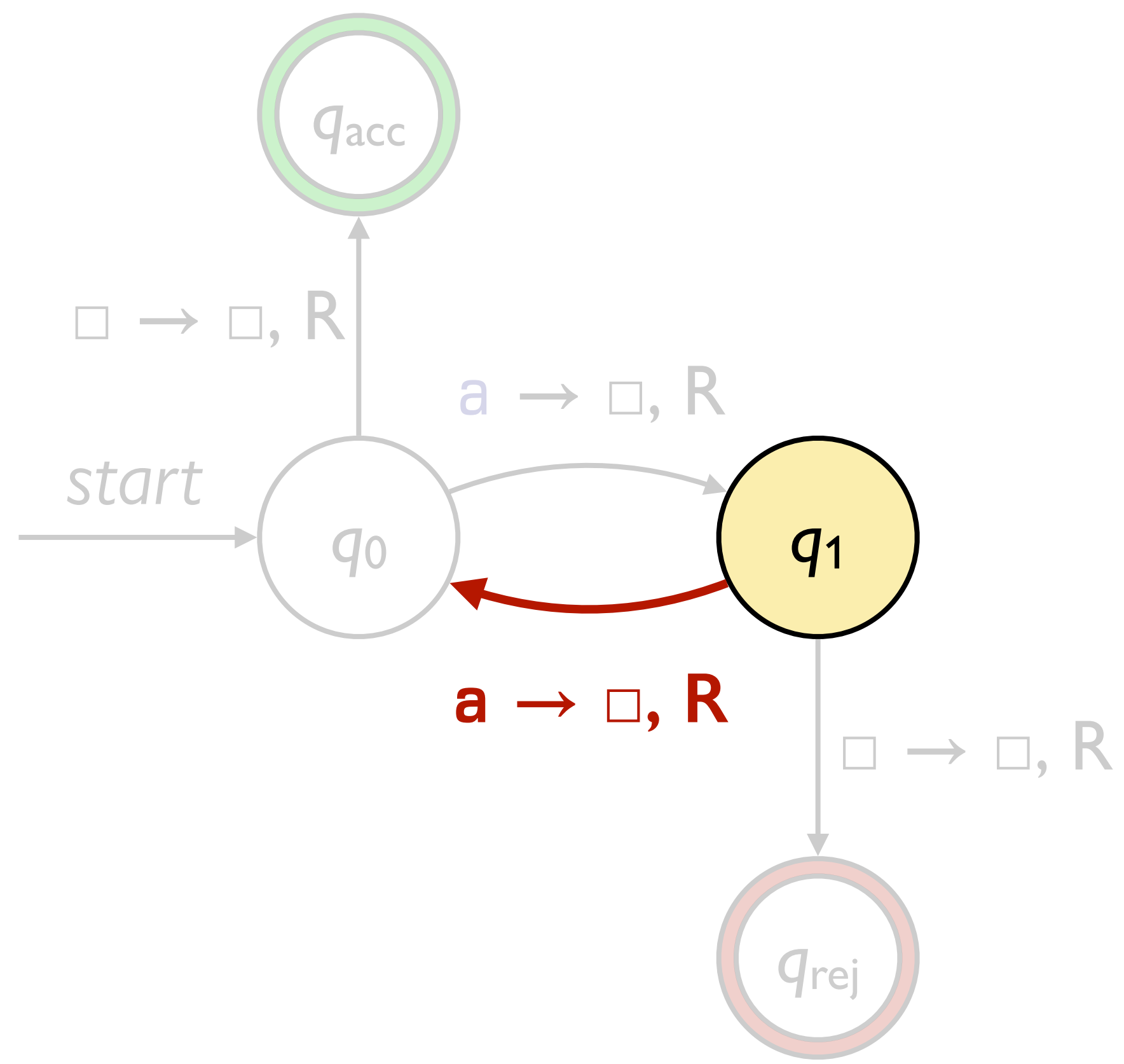


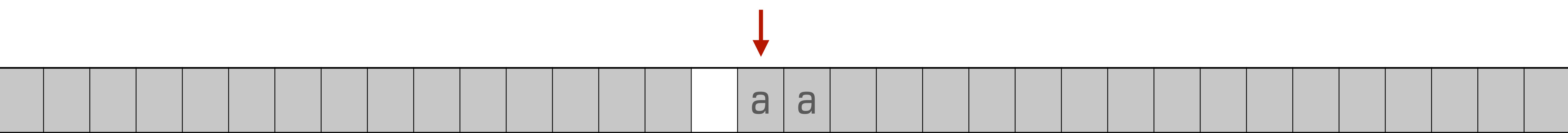
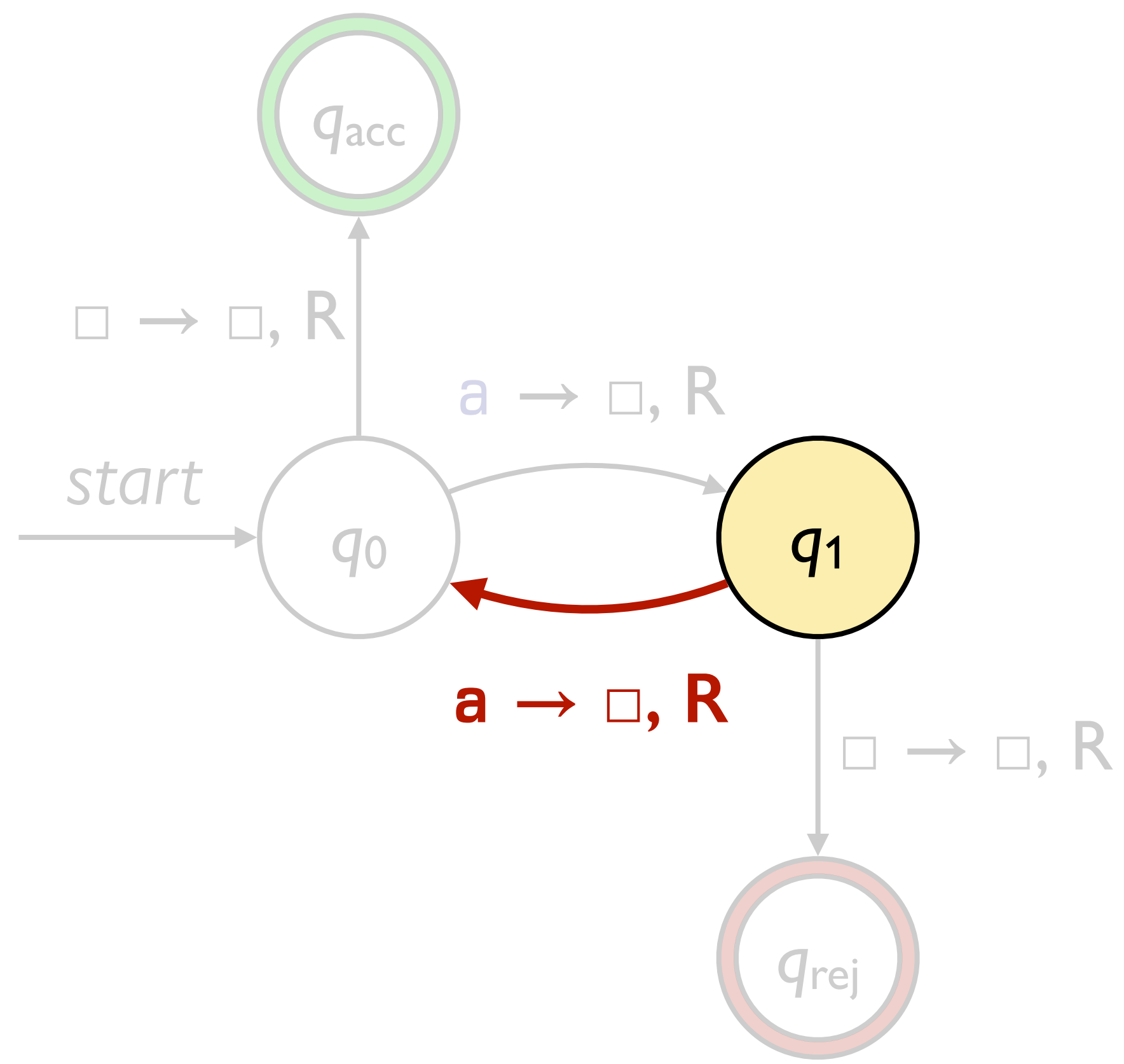


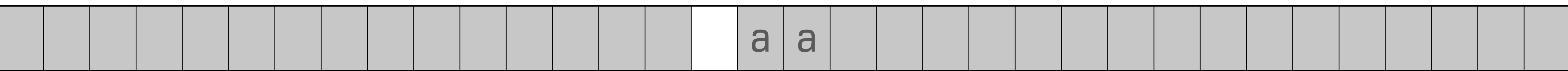
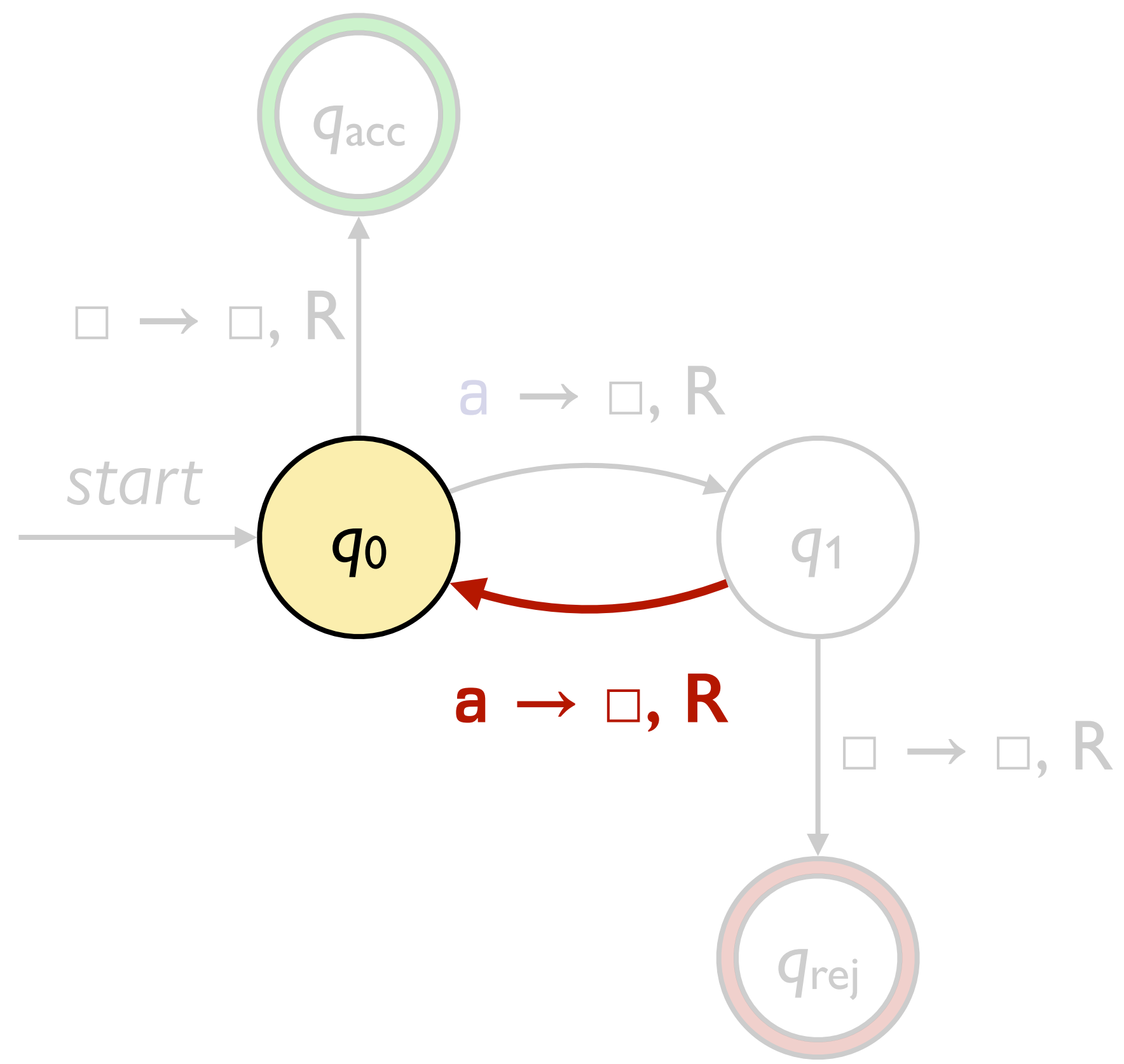


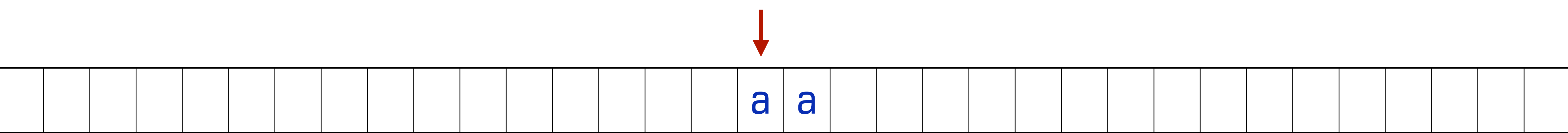
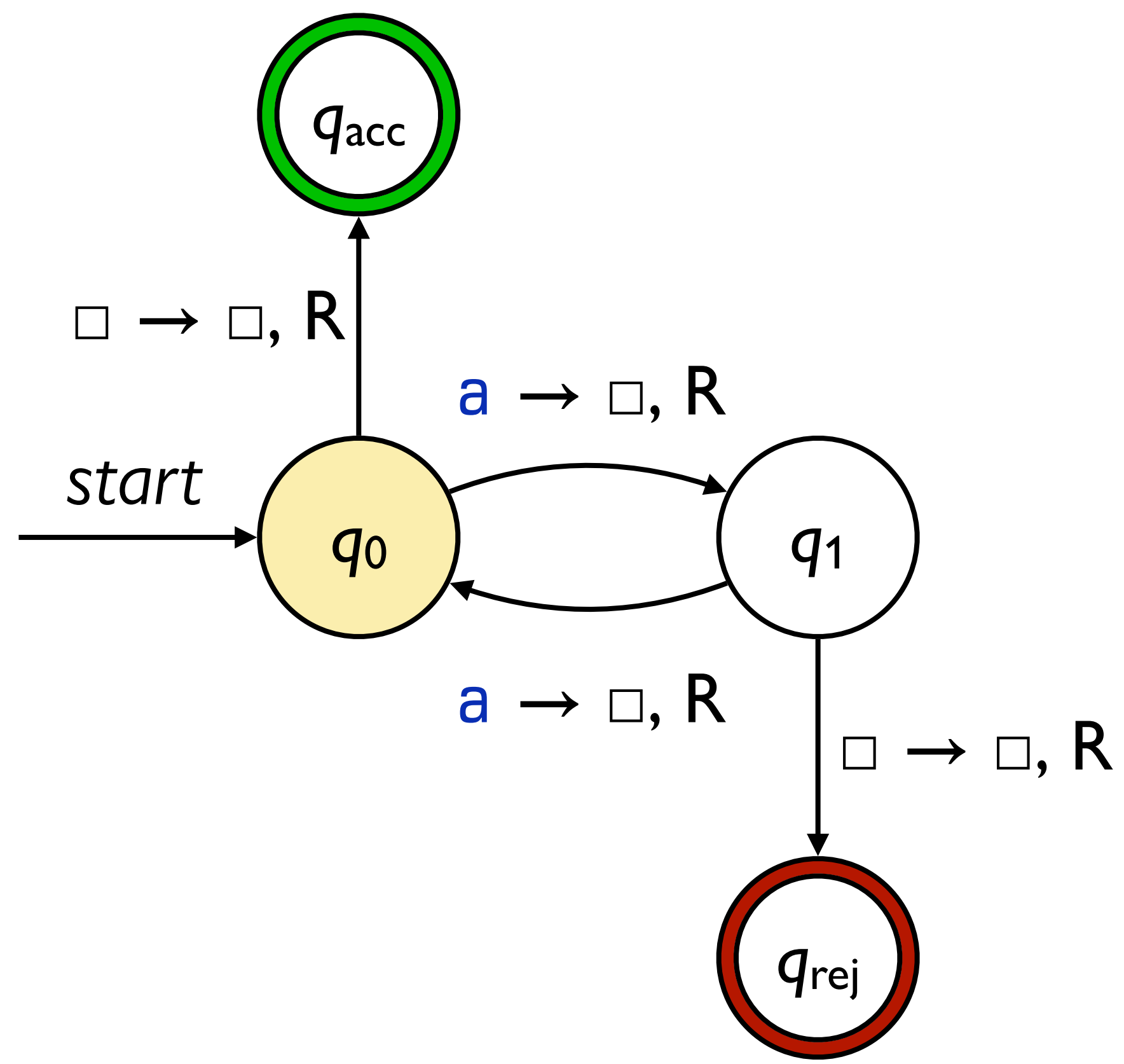


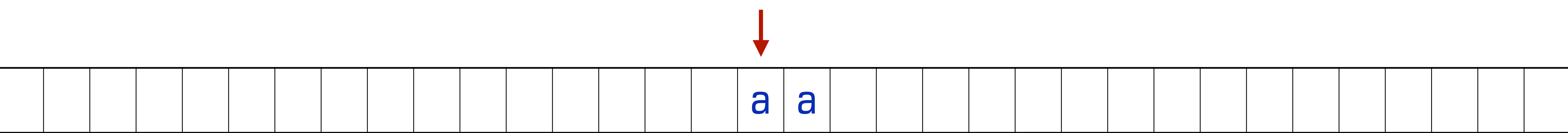
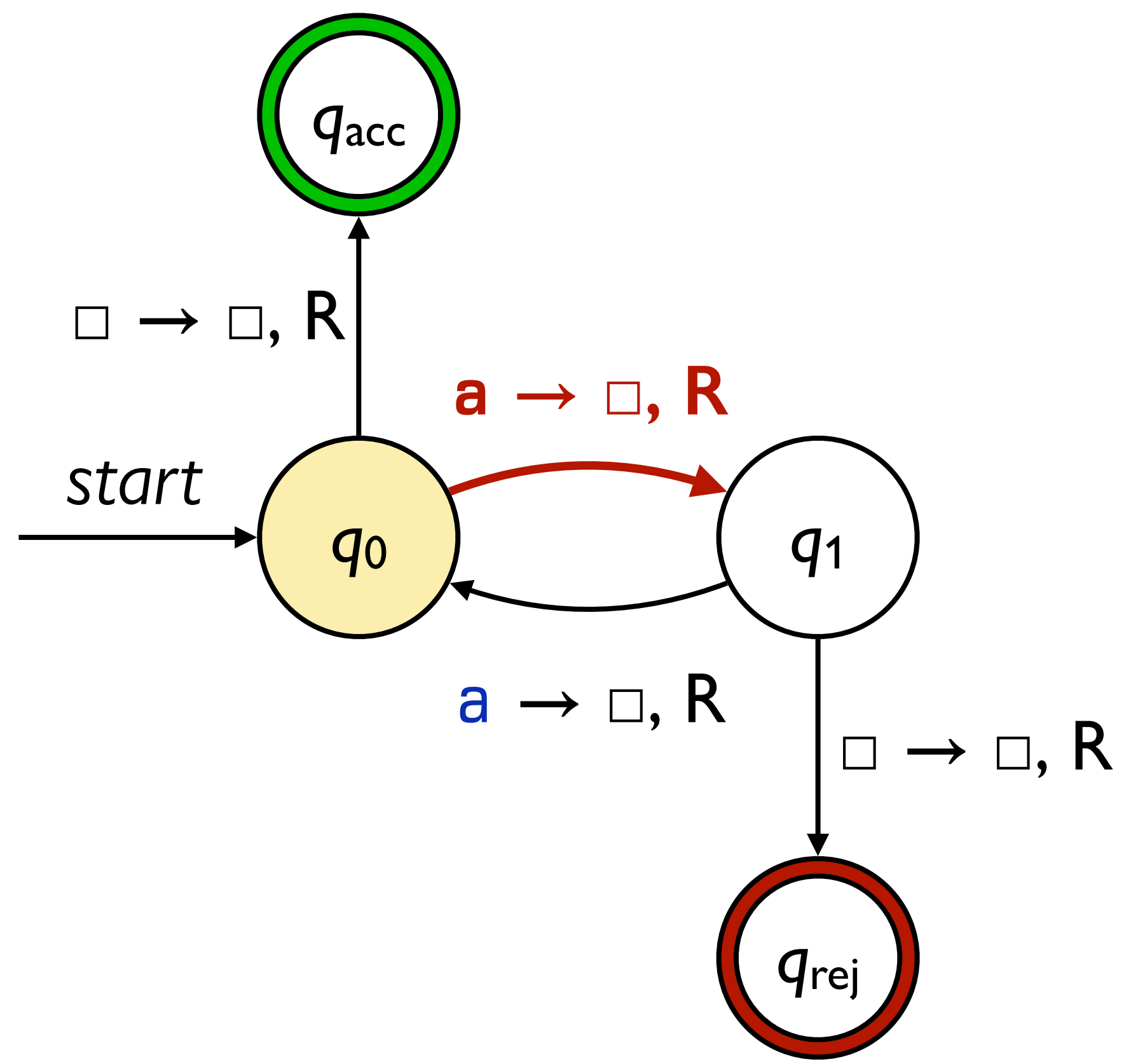


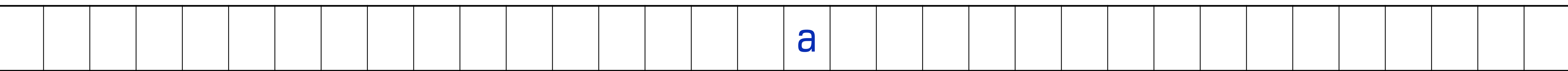
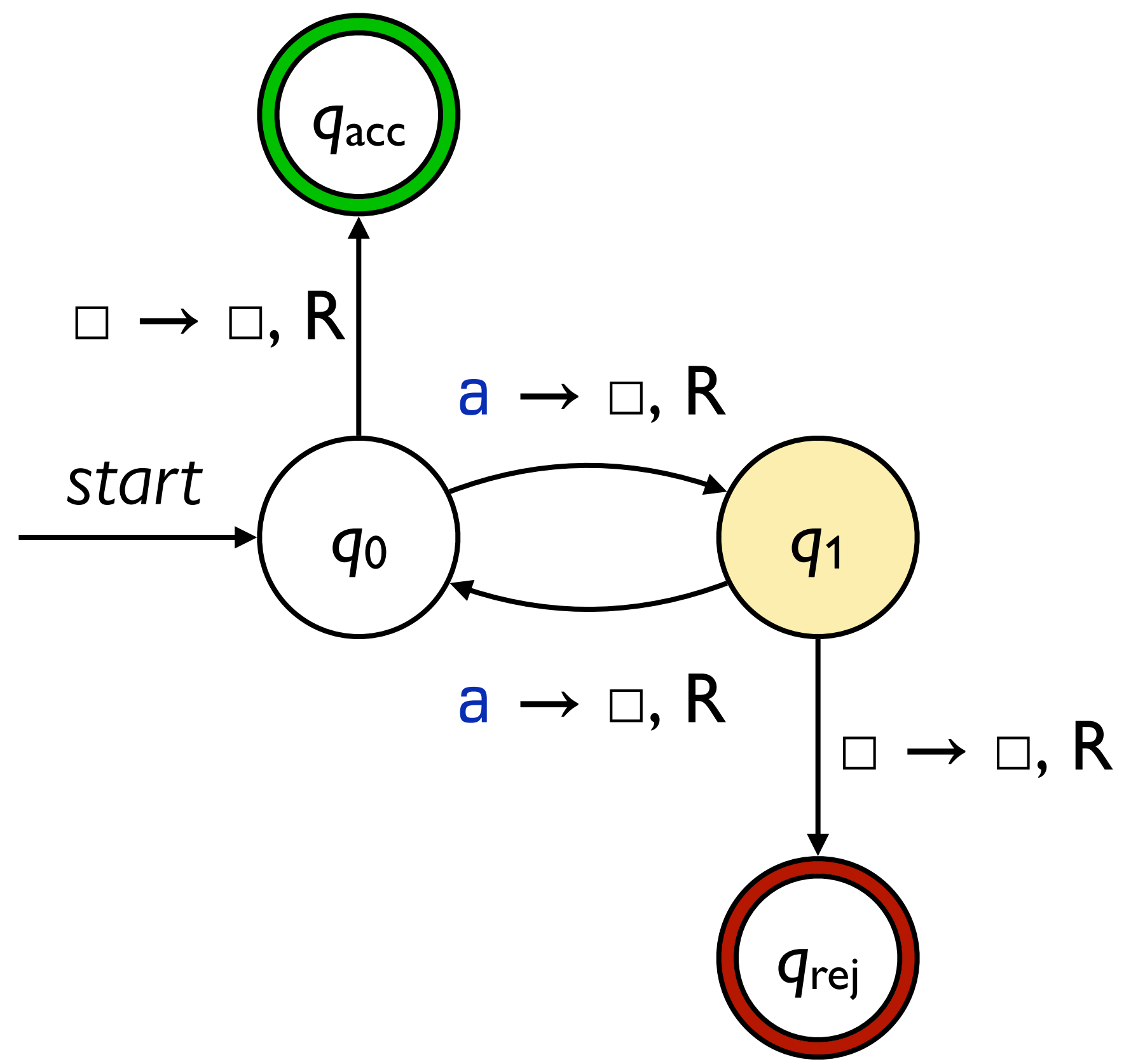


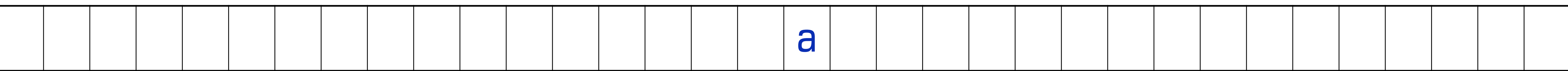
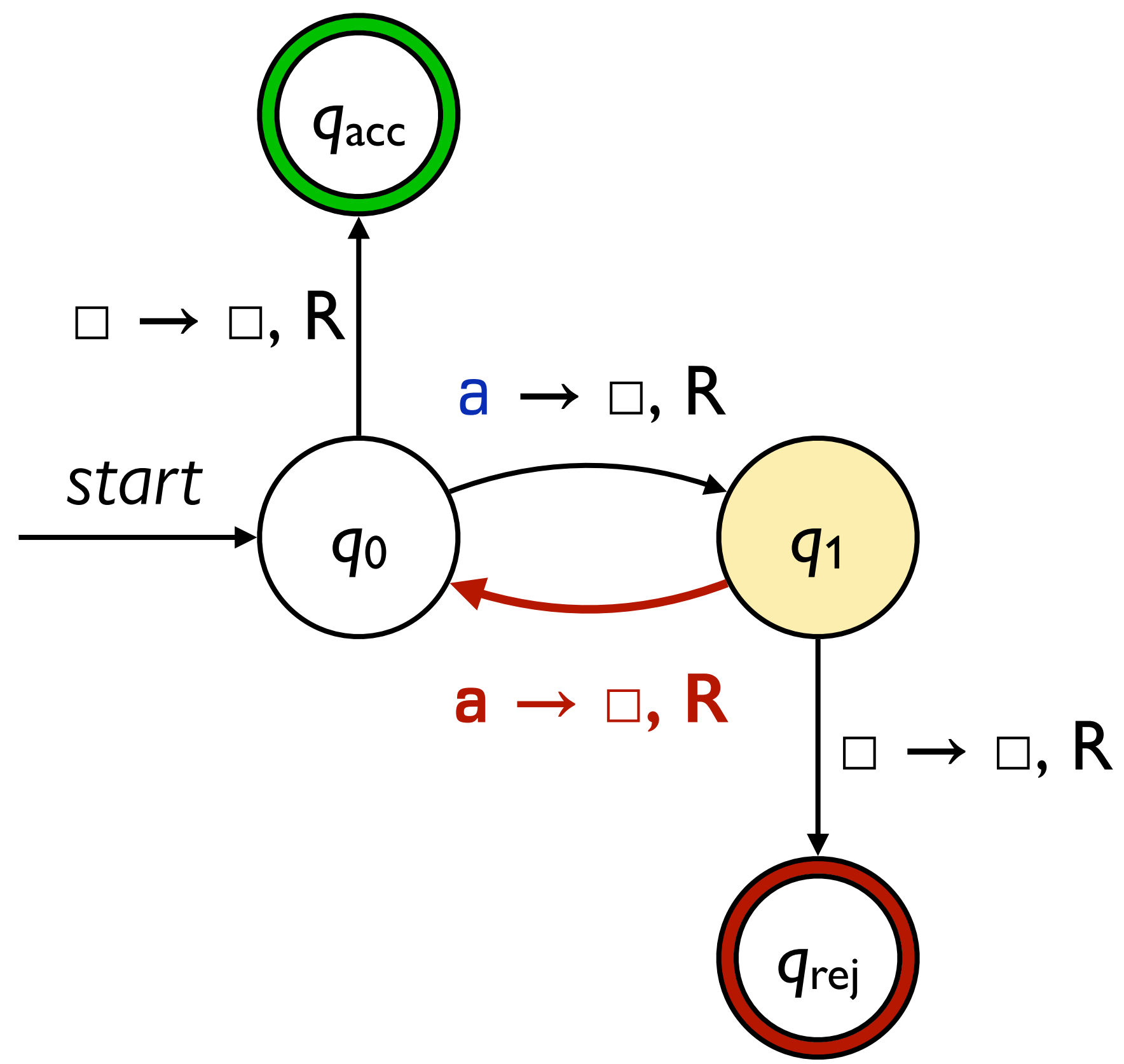


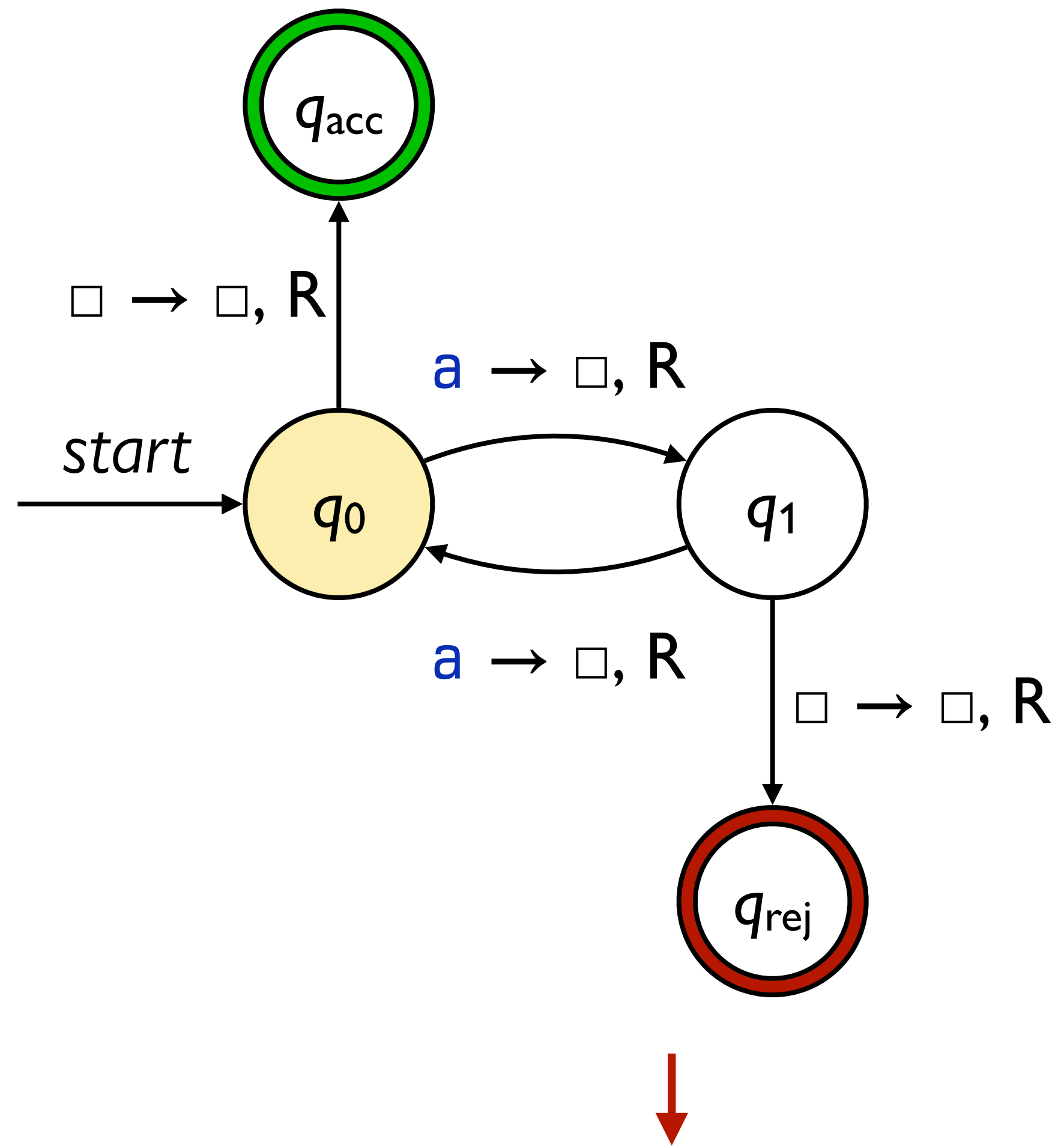




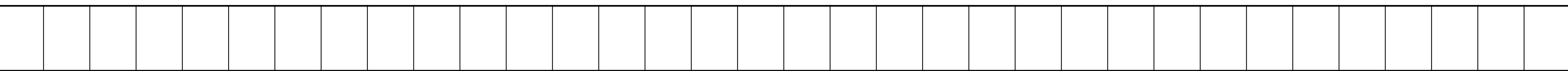
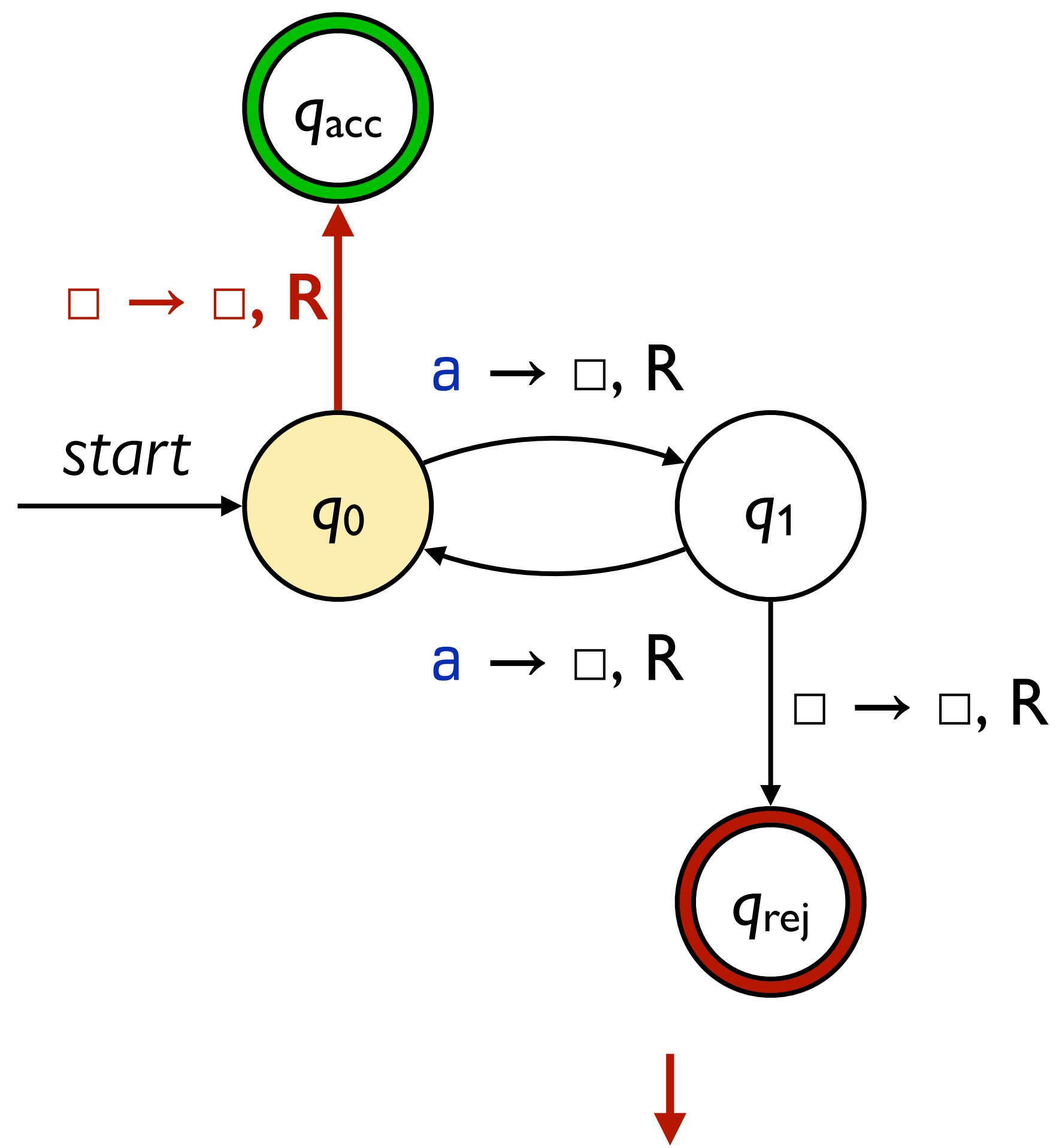




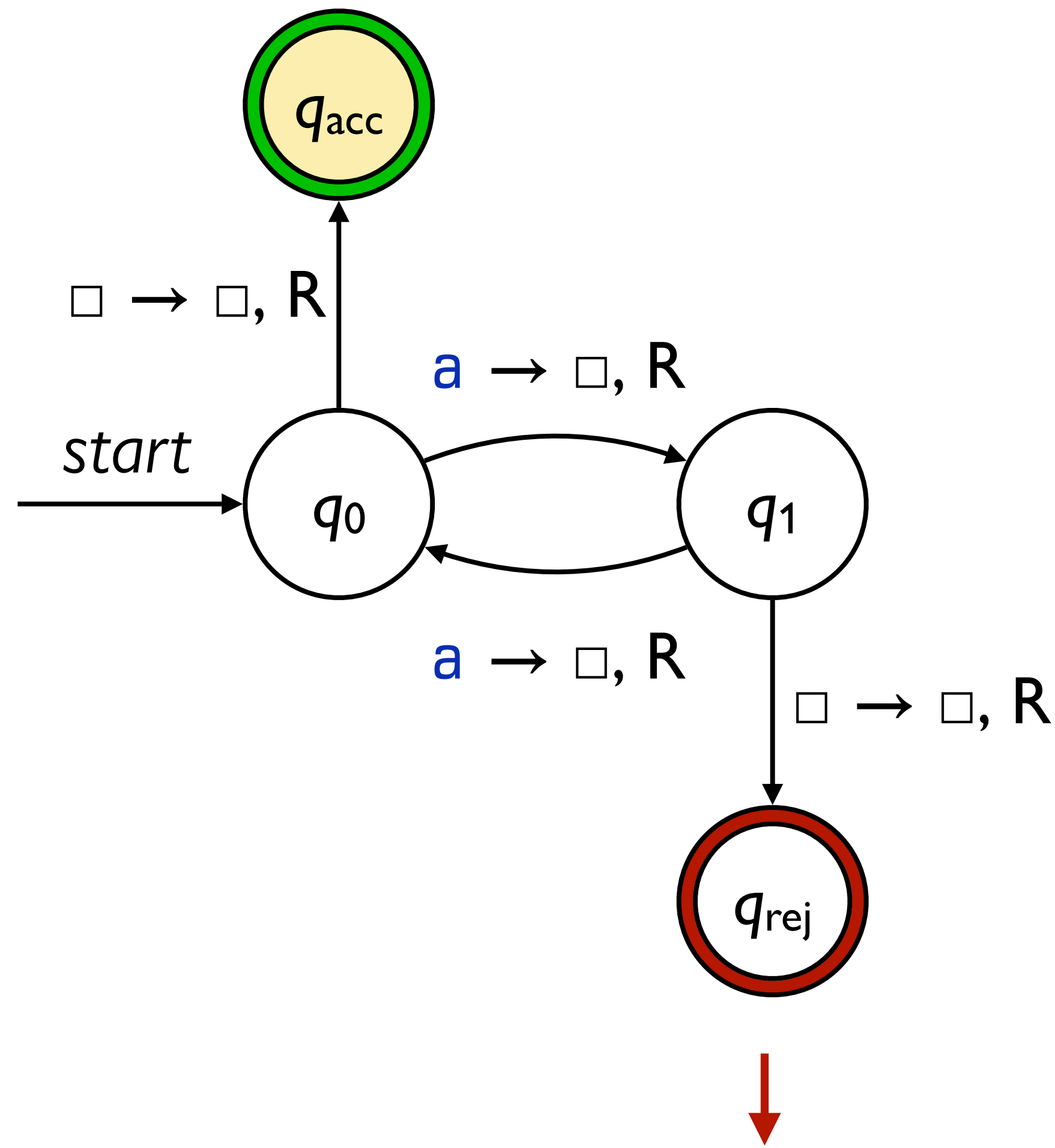




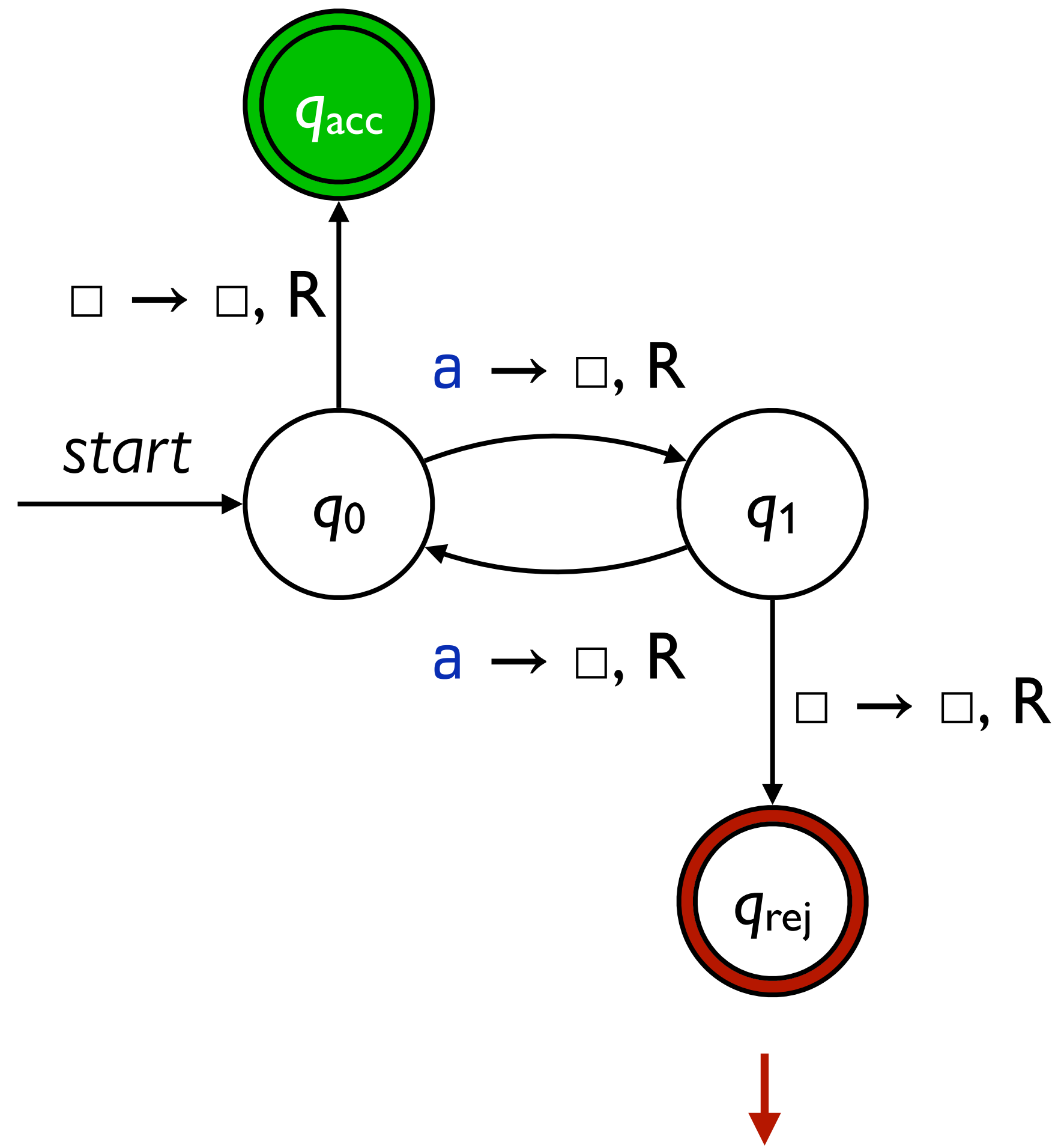
*Unlike a DFA or NFA, a TM doesn't stop after reading all the input characters. We keep running until the machine explicitly says to stop.*

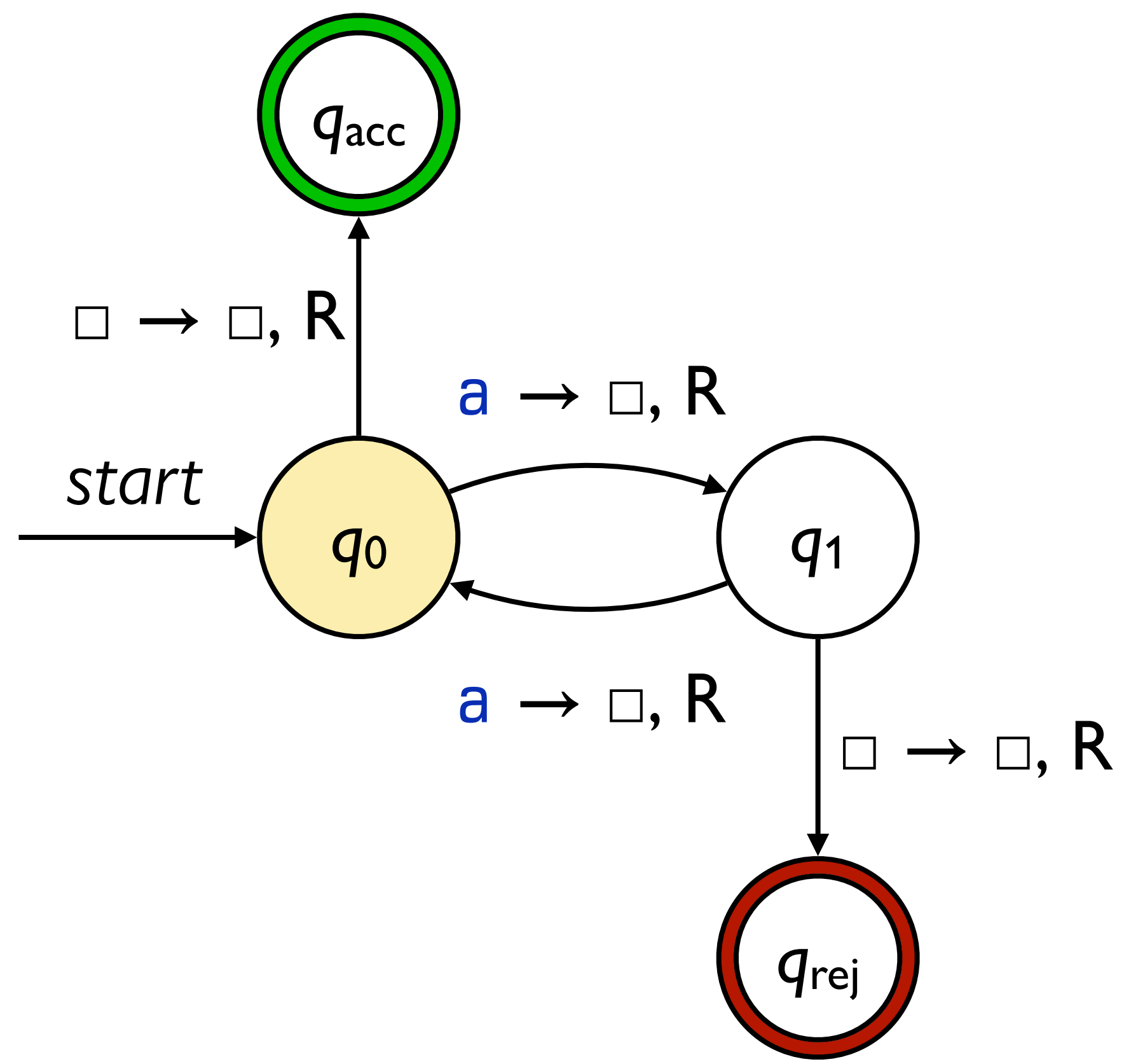


This special state is an **accept state**. When a TM enters an accept state, it **immediately stops running and accepts whatever the original input string was** – in this case, **aaaa**.

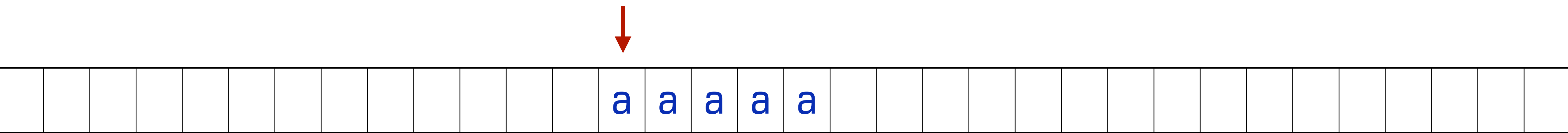


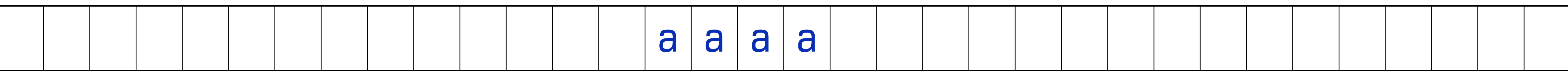
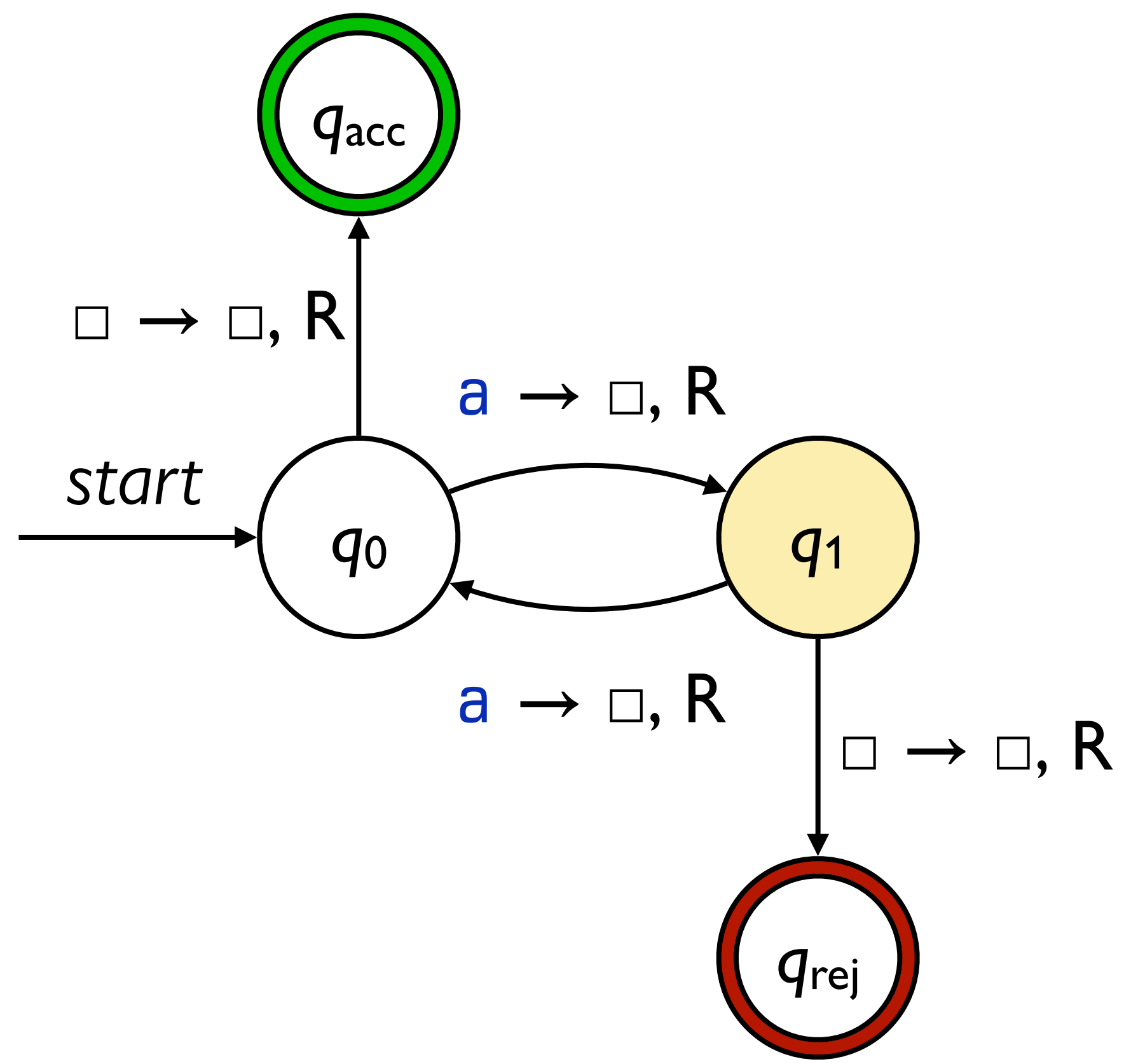
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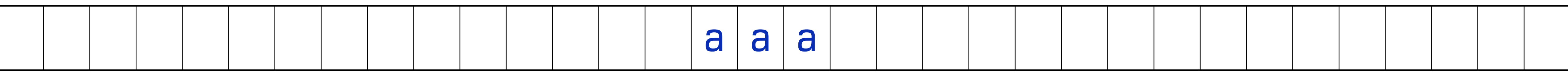
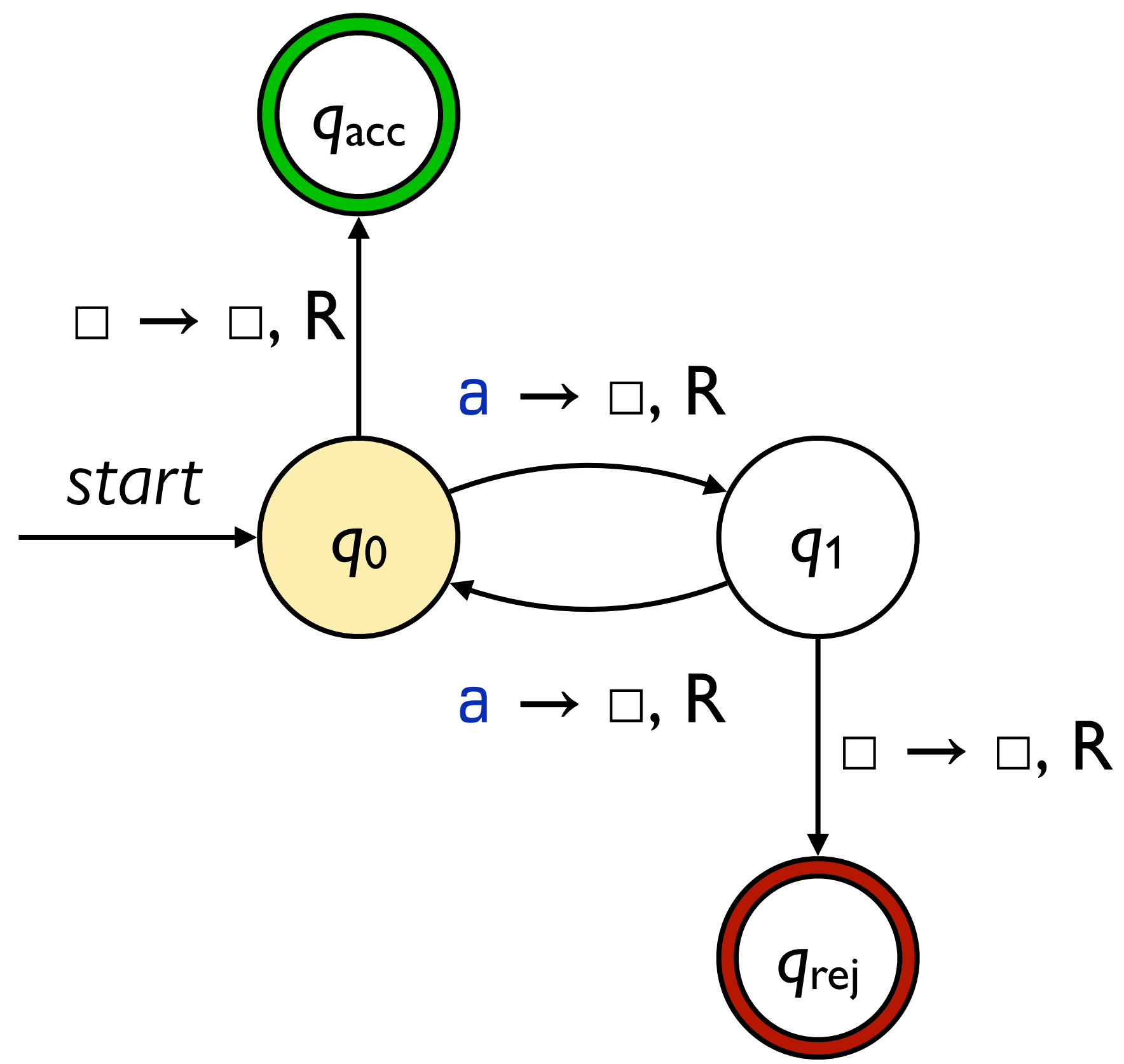


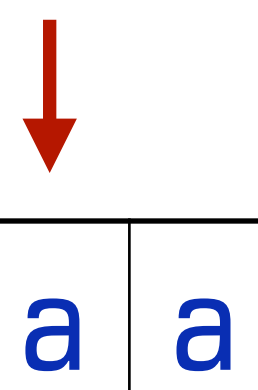
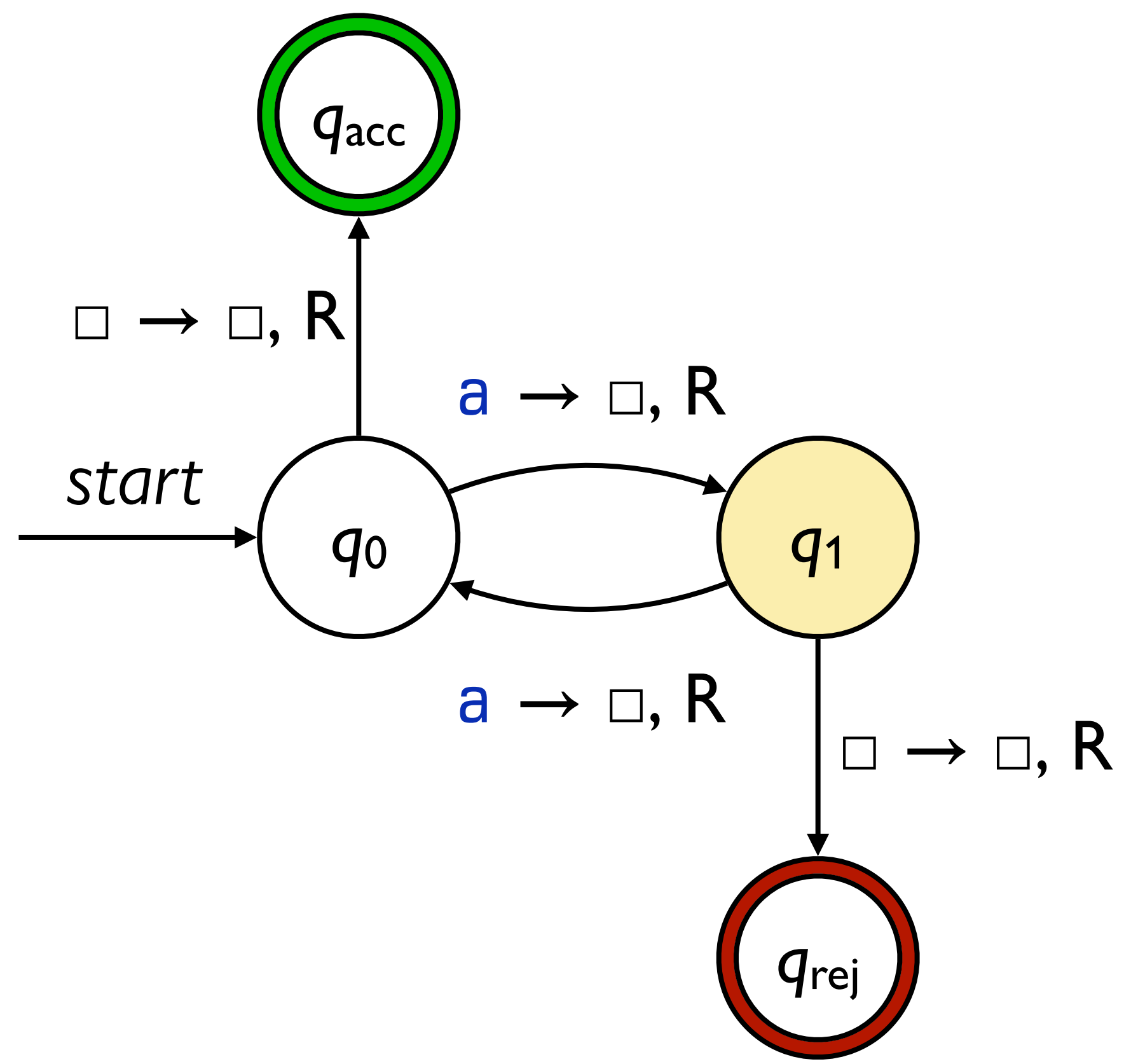


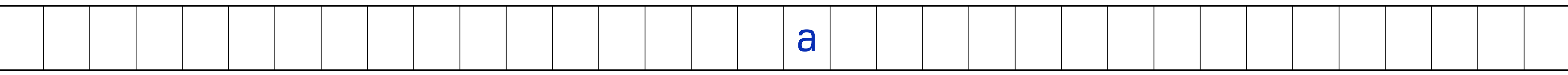
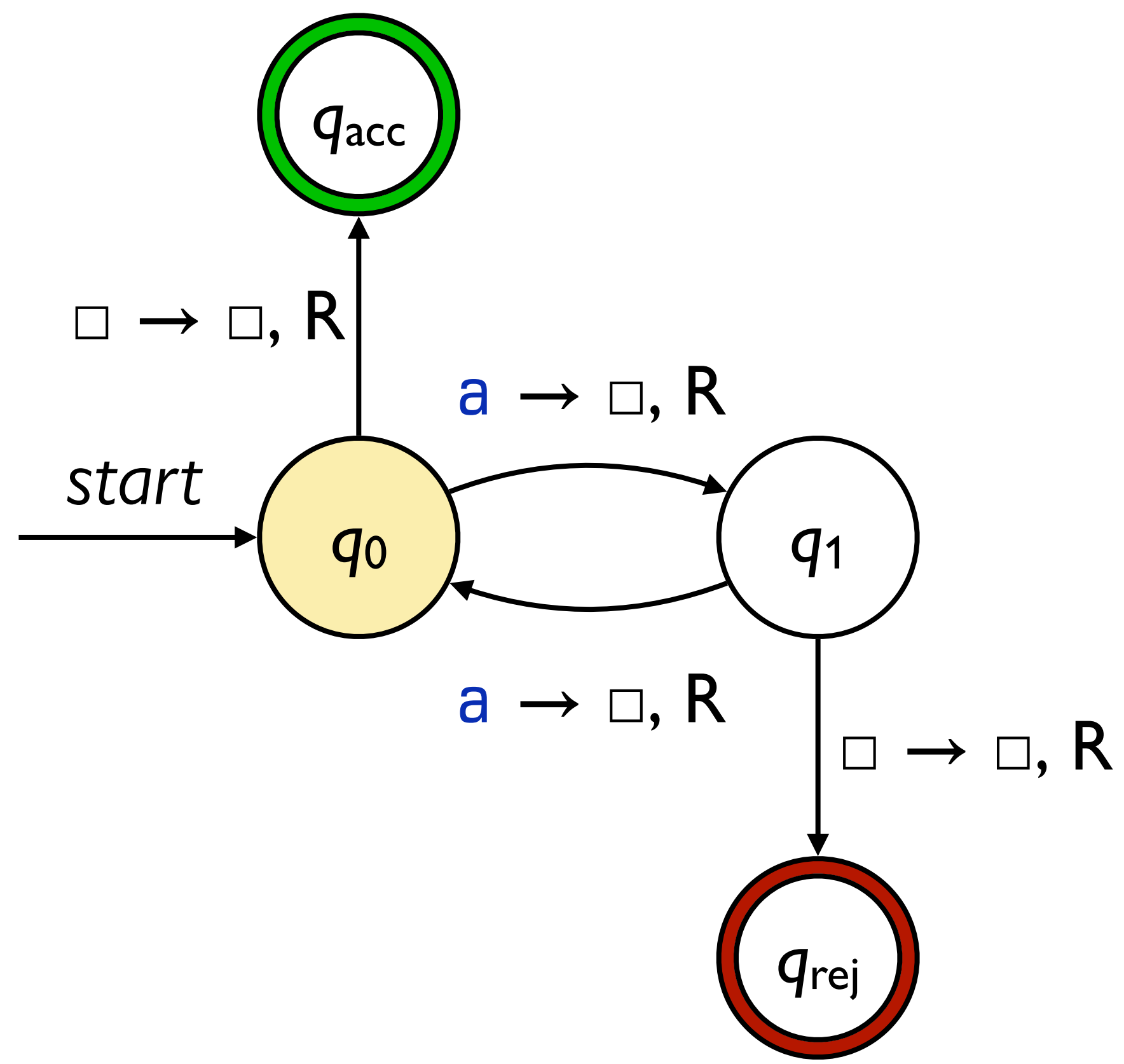
Would the TM accept this string?

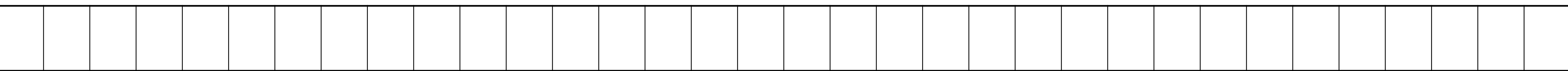
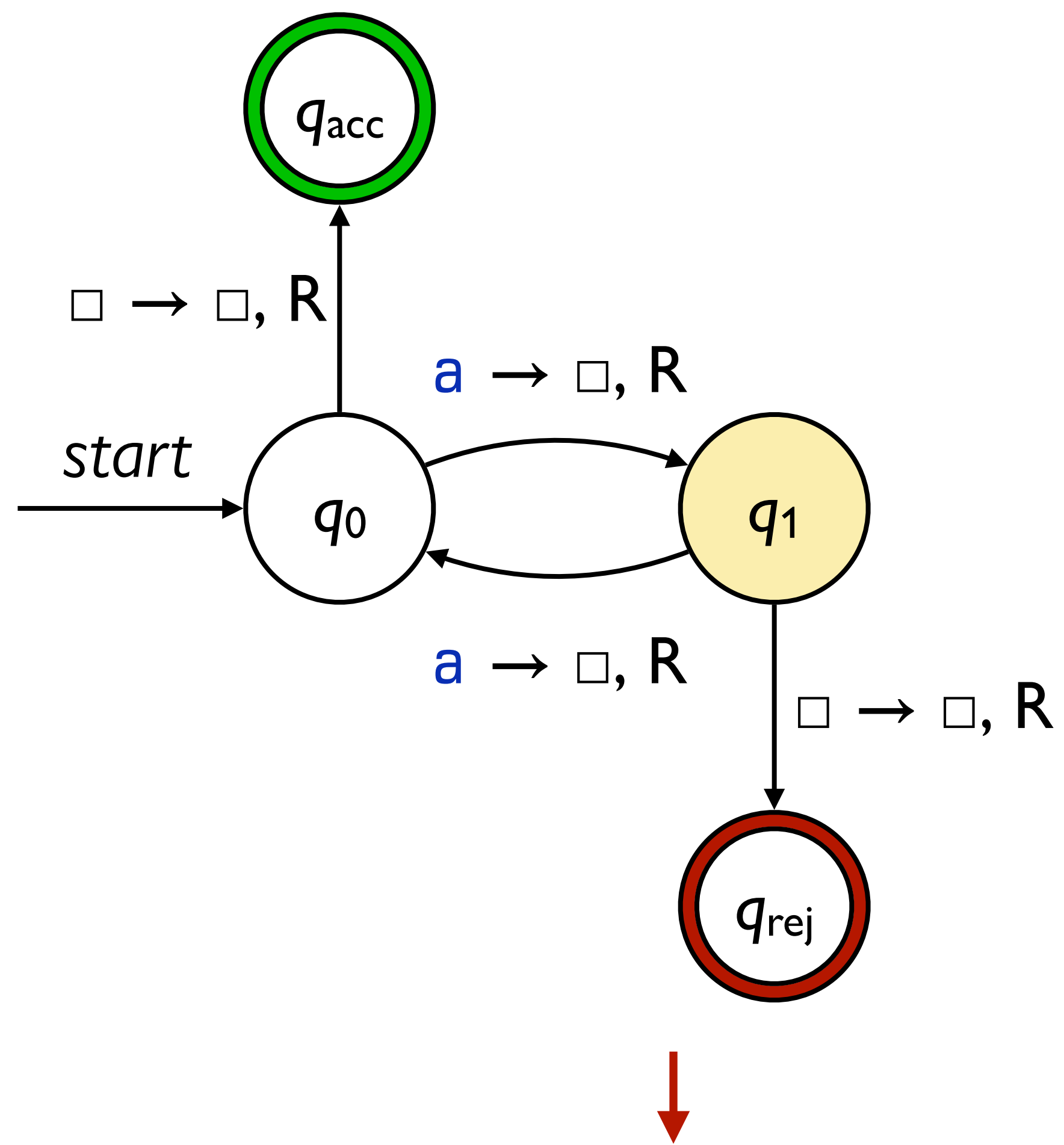


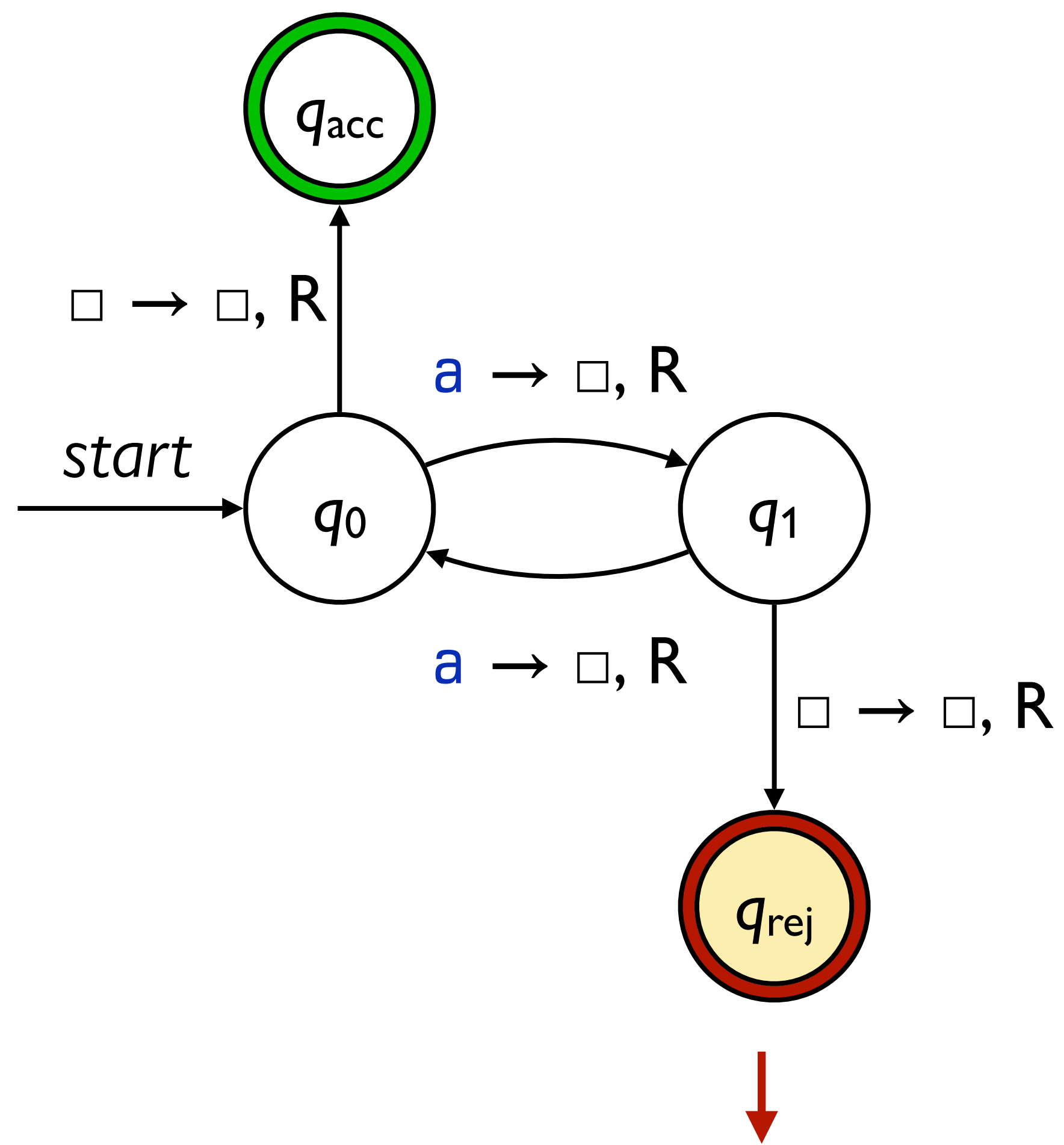




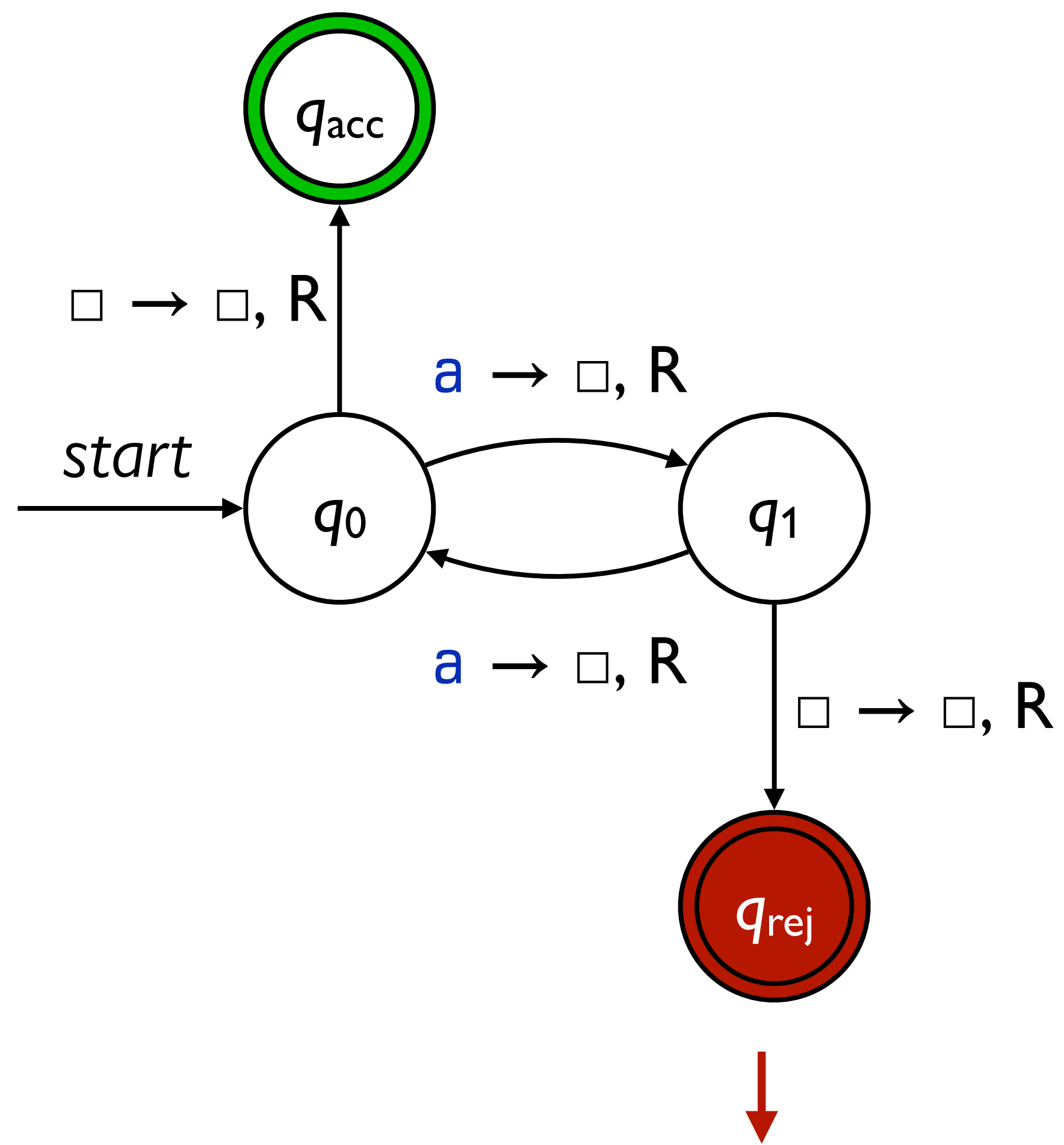






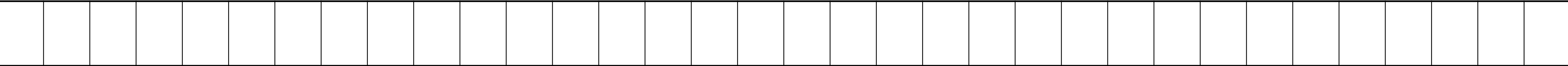
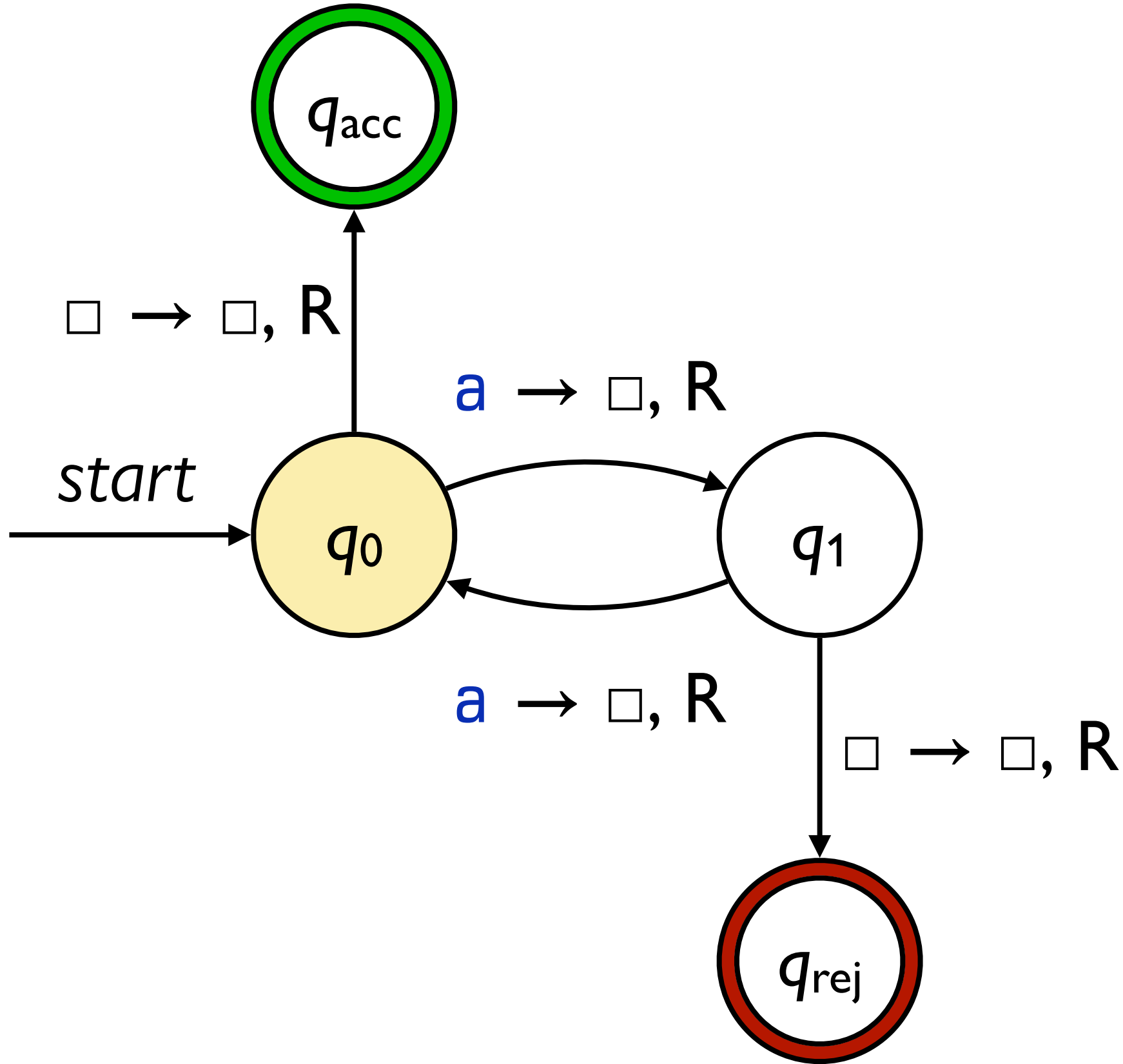


This special state is a **reject state**. When a TM enters a reject state, it **immediately** stop running and rejects whatever the original input string was – in this case, **aaaaa**.

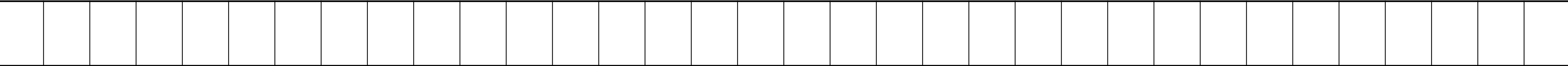
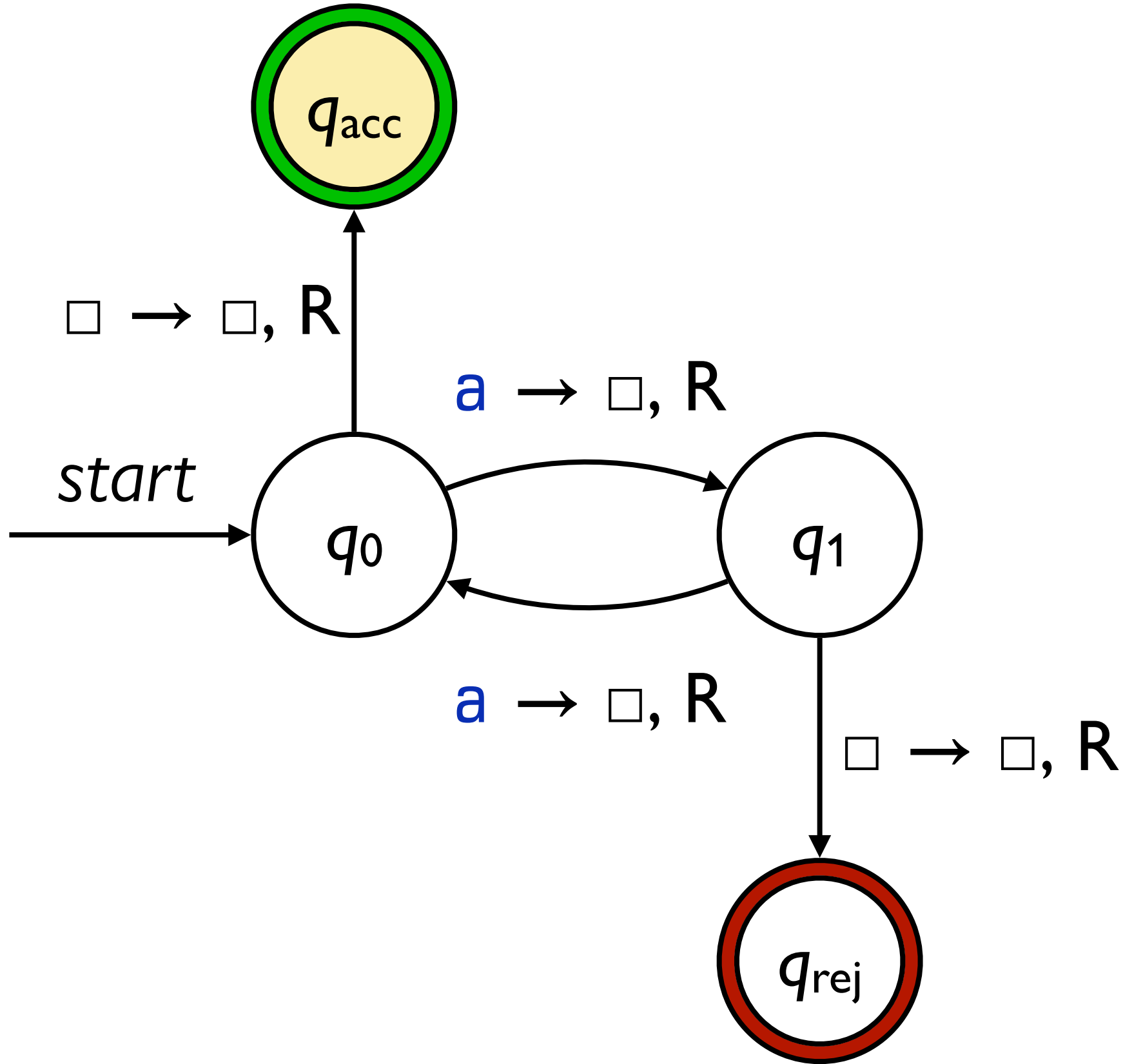


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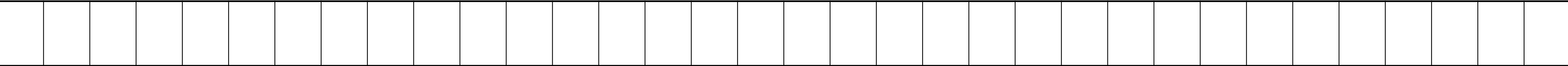
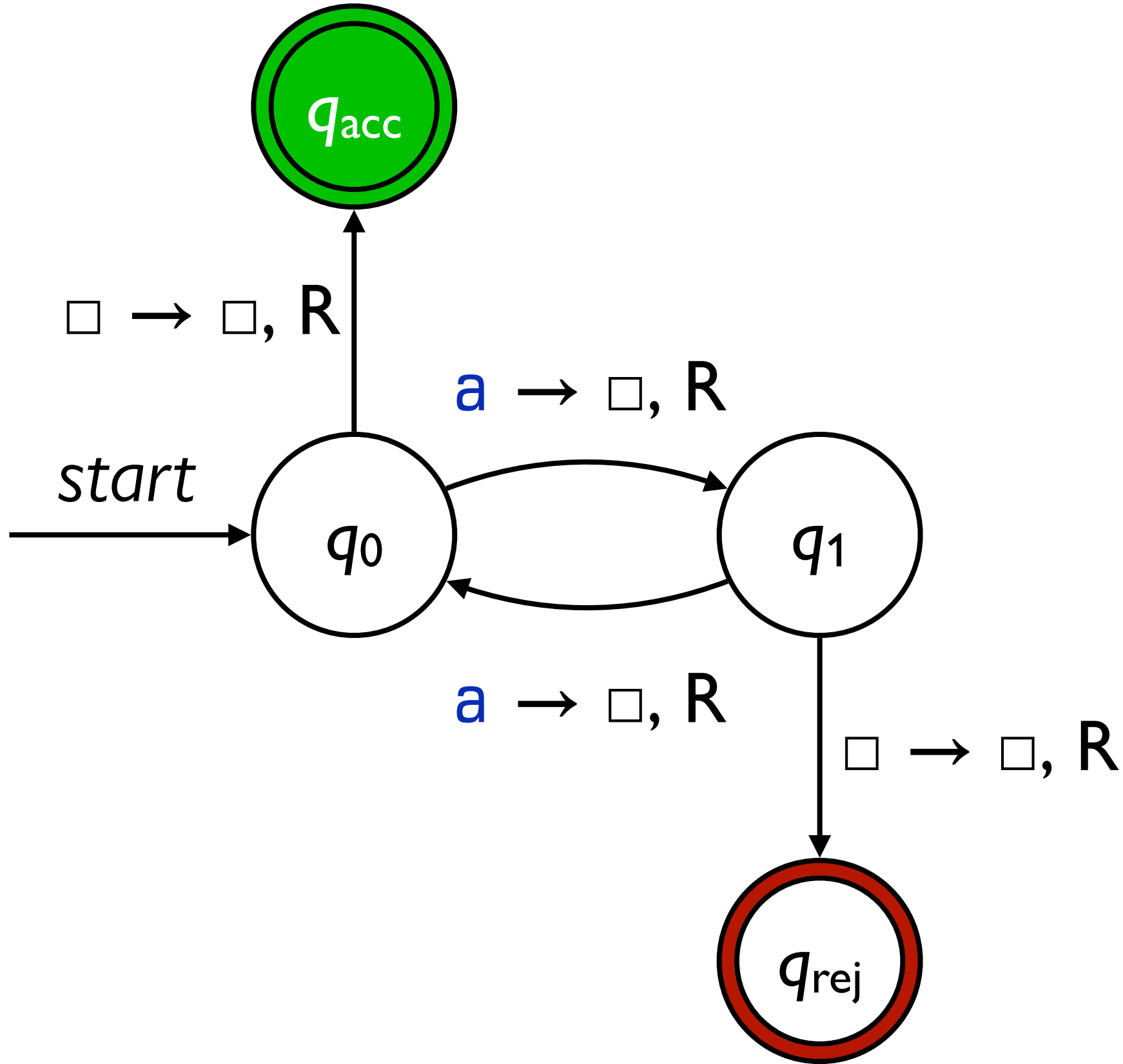
If the TM is started on the empty string  $\epsilon$ , the entire tape is blank and the tape head is positioned at some arbitrary location on the tape.



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A Turing machine is  $(Q, \Sigma, \Gamma, \delta, q_o, q_{acc}, q_{rej})$ , where

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$\Sigma$  is a finite *input alphabet*,

$\Gamma$  is a finite *tape alphabet*, with  $\Sigma \subset \Gamma$ ,

*The tape alphabet  $\Gamma$  can contain any number of symbols but always includes at least one **blank symbol**,  $\square \notin \Sigma$ .*

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$\delta : Q \times \Gamma \rightarrow Q \times \Gamma \times \{L, R\}$  is the *transition function*,

*At each step, the Turing machine*

*Writes a symbol to the tape cell under the tape head,*

*Changes state, and*

*Moves the tape head left or right.*

A Turing machine is  $(Q, \Sigma, \Gamma, \delta, q_0, q_{\text{acc}}, q_{\text{rej}})$ , where

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$\delta : Q \times \Gamma \rightarrow Q \times \Gamma \times \{L, R\}$  is the *transition function*,

$q_0 \in Q$  is the *start state*,

$q_{\text{acc}} \in Q$  is the *accept state*, and

$q_{\text{rej}} \in Q, q_{\text{rej}} \neq q_{\text{acc}}$  is the *reject state*.

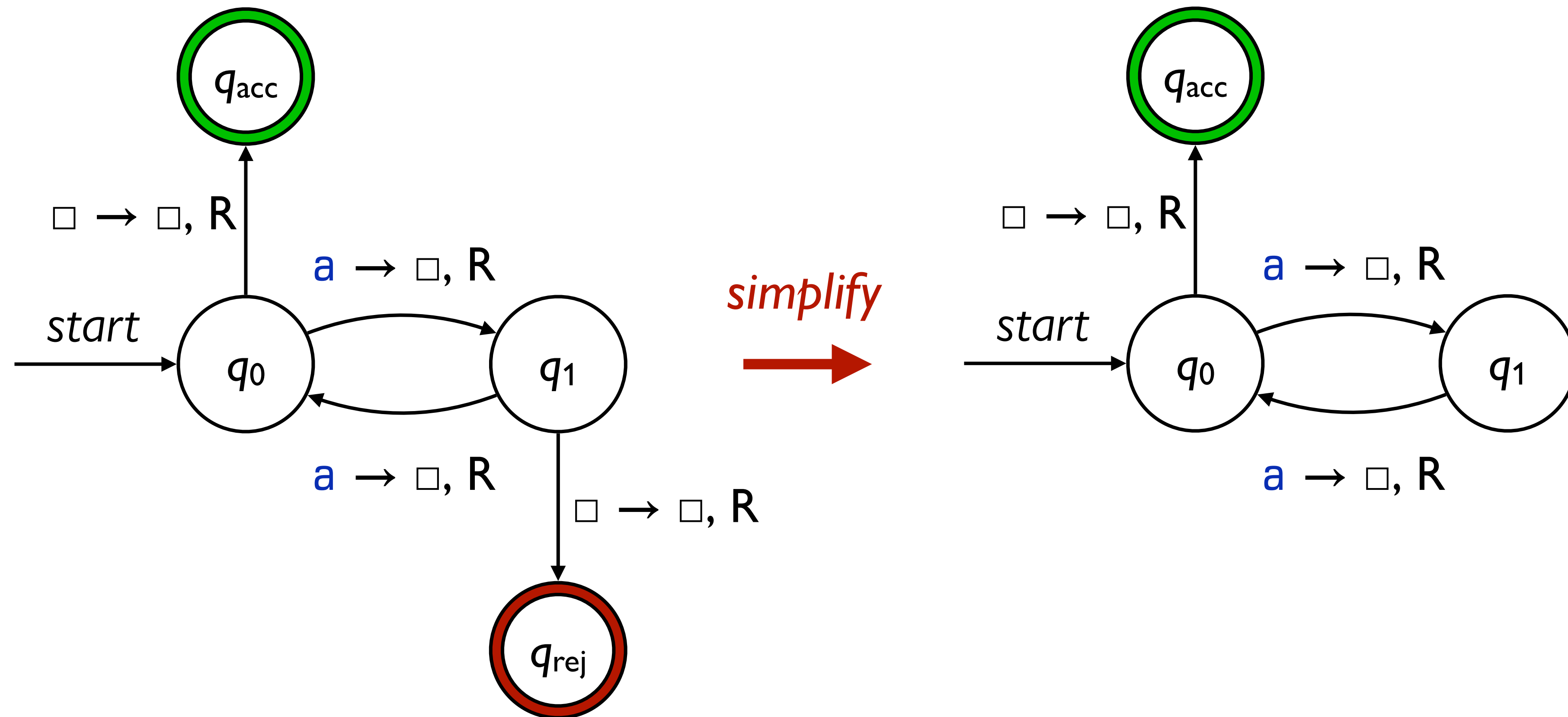
Unlike the automata we've seen before, Turing machines can revisit characters from the input, and the machine doesn't need to read the entire input.

Turing machines decide when – and if! – they will accept or reject their input.

Turing machines are *deterministic*.

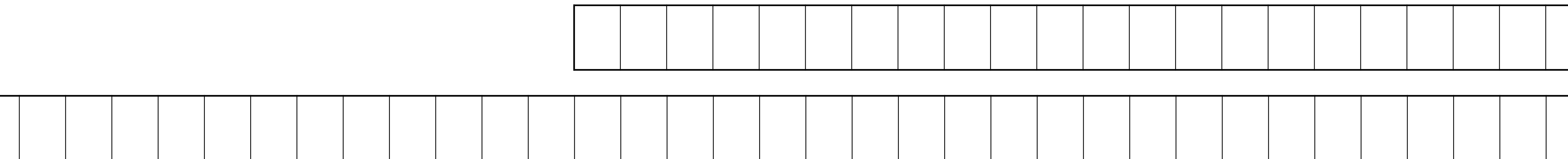
For every combination of a (non-accepting, non-rejecting) state  $q$  and a tape symbol  $a \in \Gamma$ , there must be exactly one transition defined for that combination of  $q$  and  $a$ .

Any transitions that are missing implicitly go straight to a rejecting state. We'll use this later to simplify our designs.

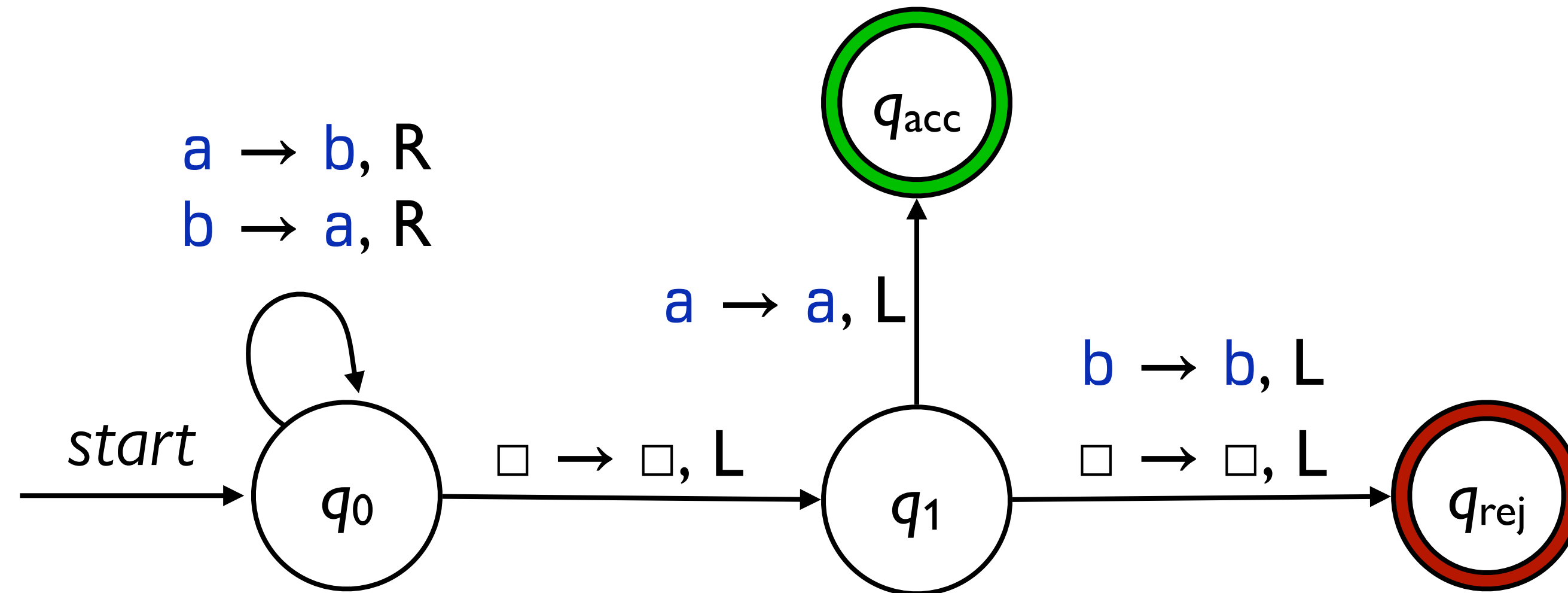


*Warning:* The textbook treats the tape of a Turing machine as only being infinite to the right.

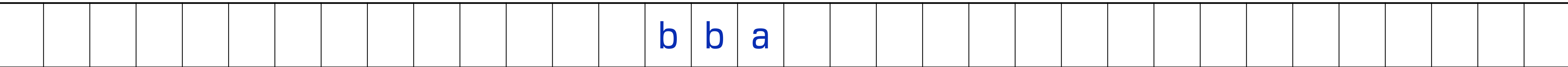
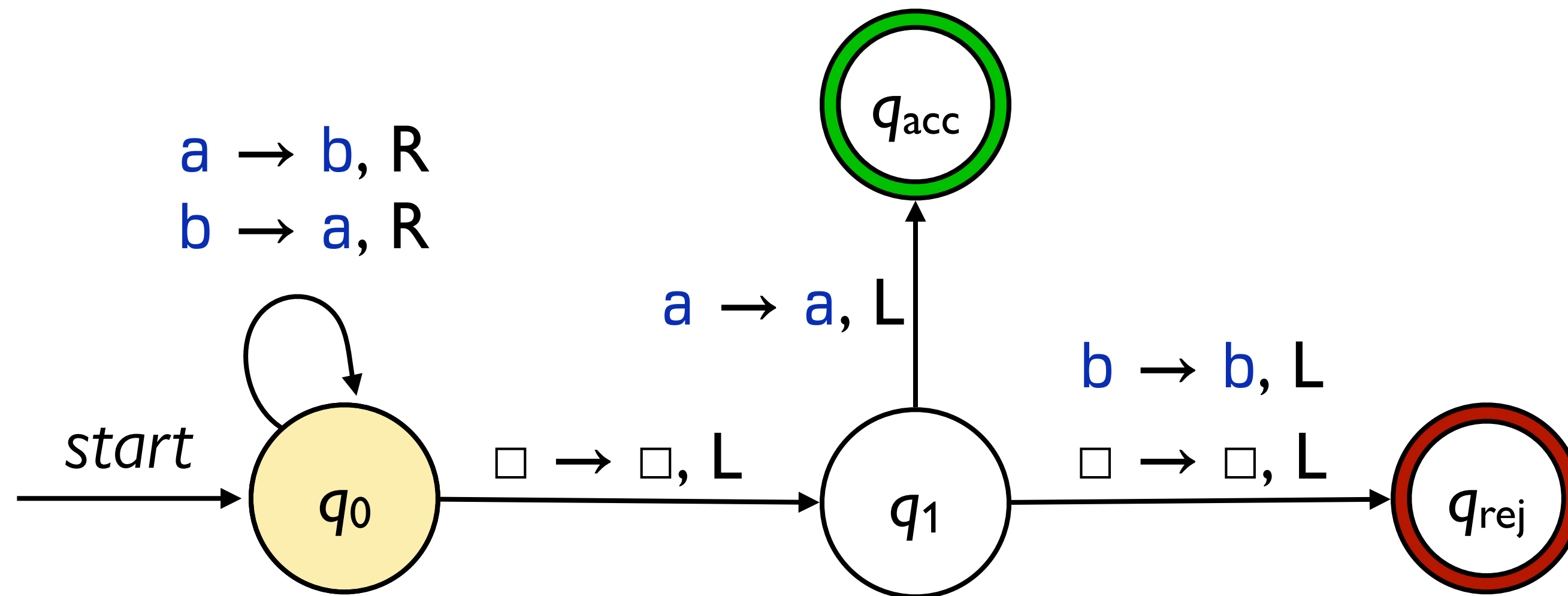
This is no less powerful than a tape that's infinite in both directions, but it's more annoying to use.

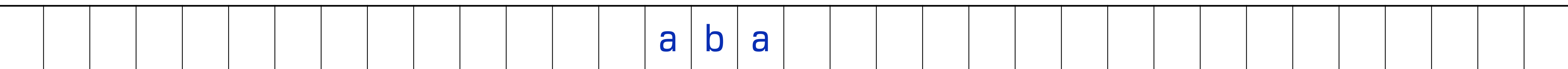
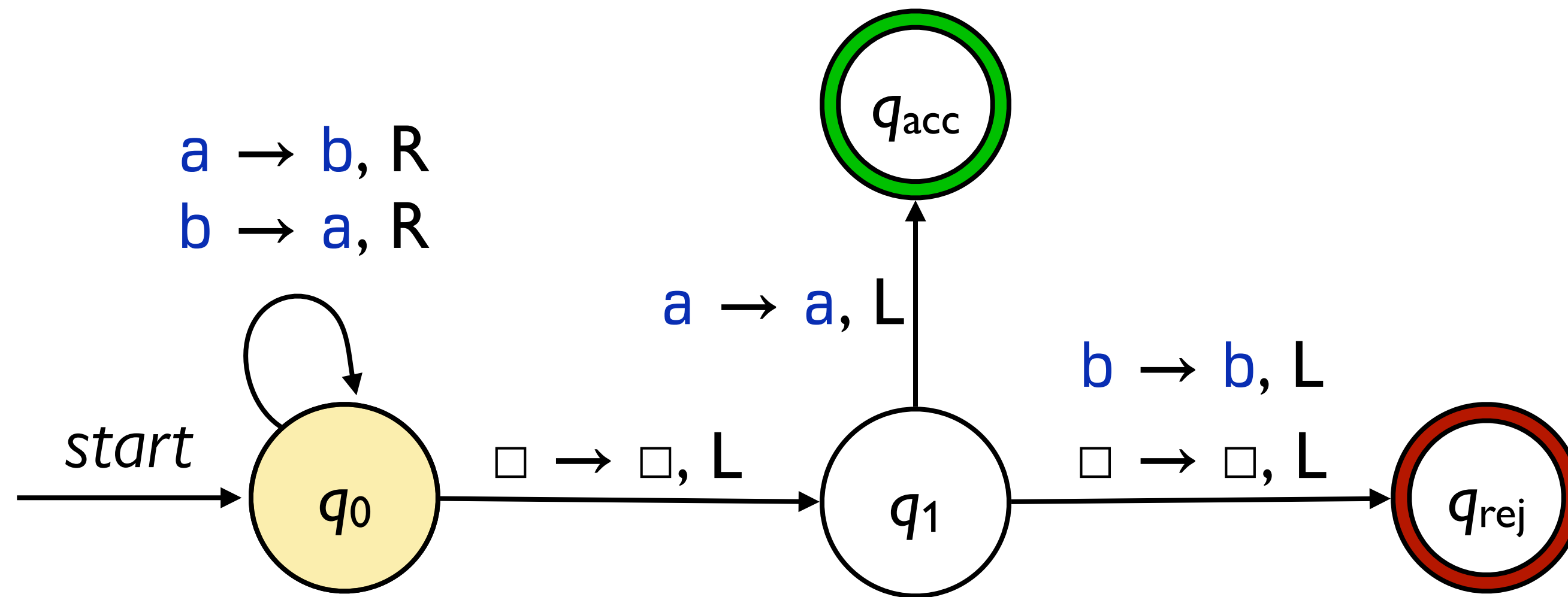


The language of a Turing machine



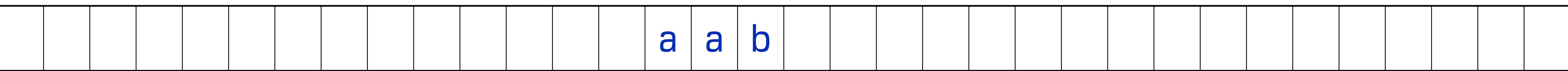
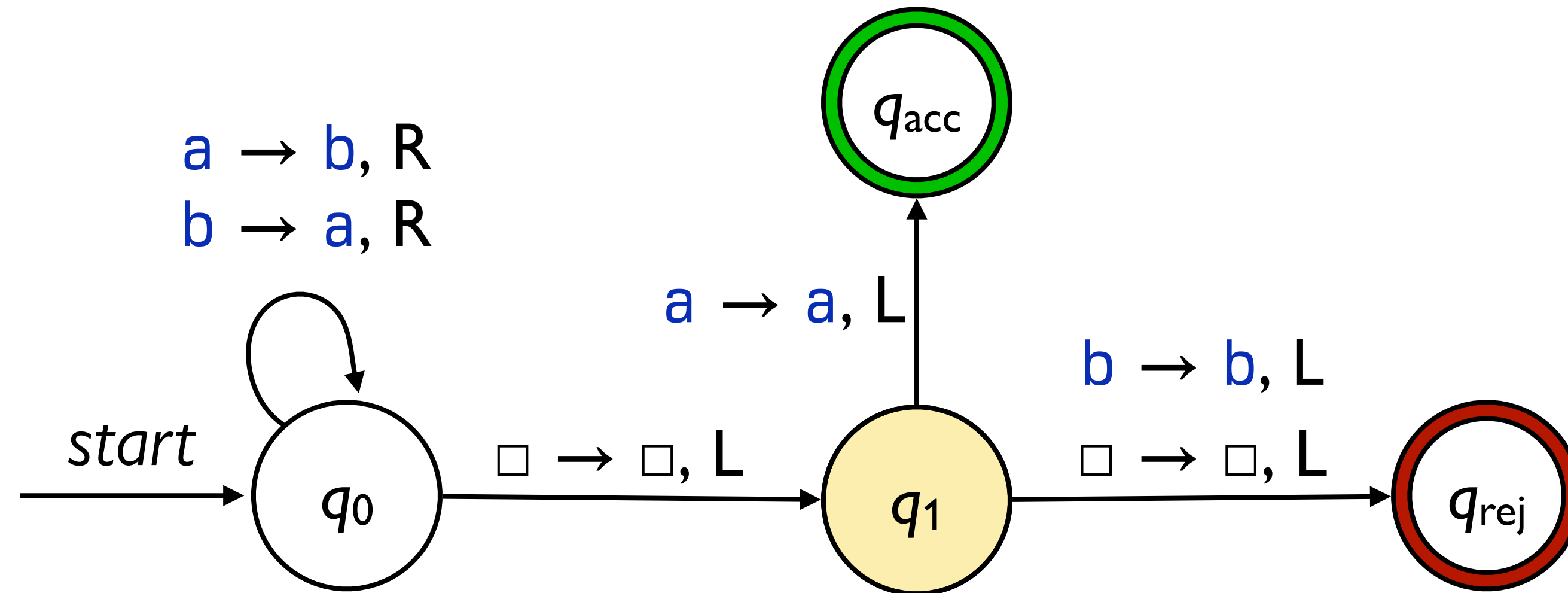
Run the TM shown above on the input string **bba**.  
 What will the tape look like when the TM finishes running?

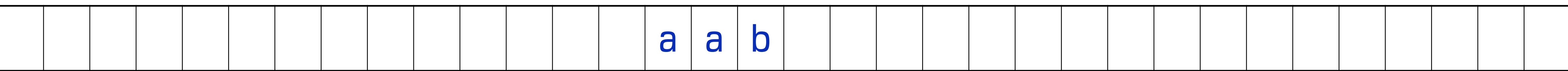
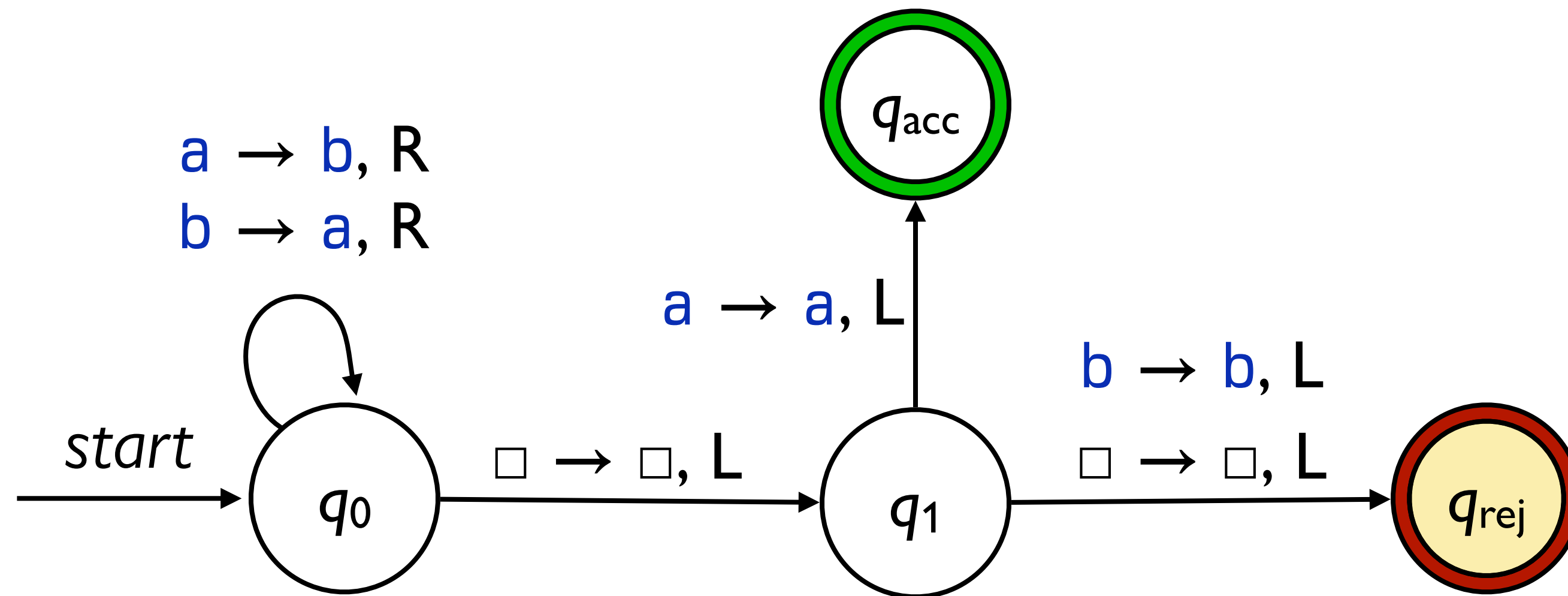


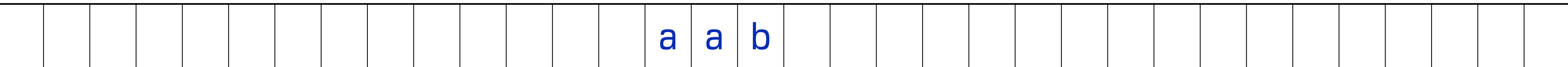
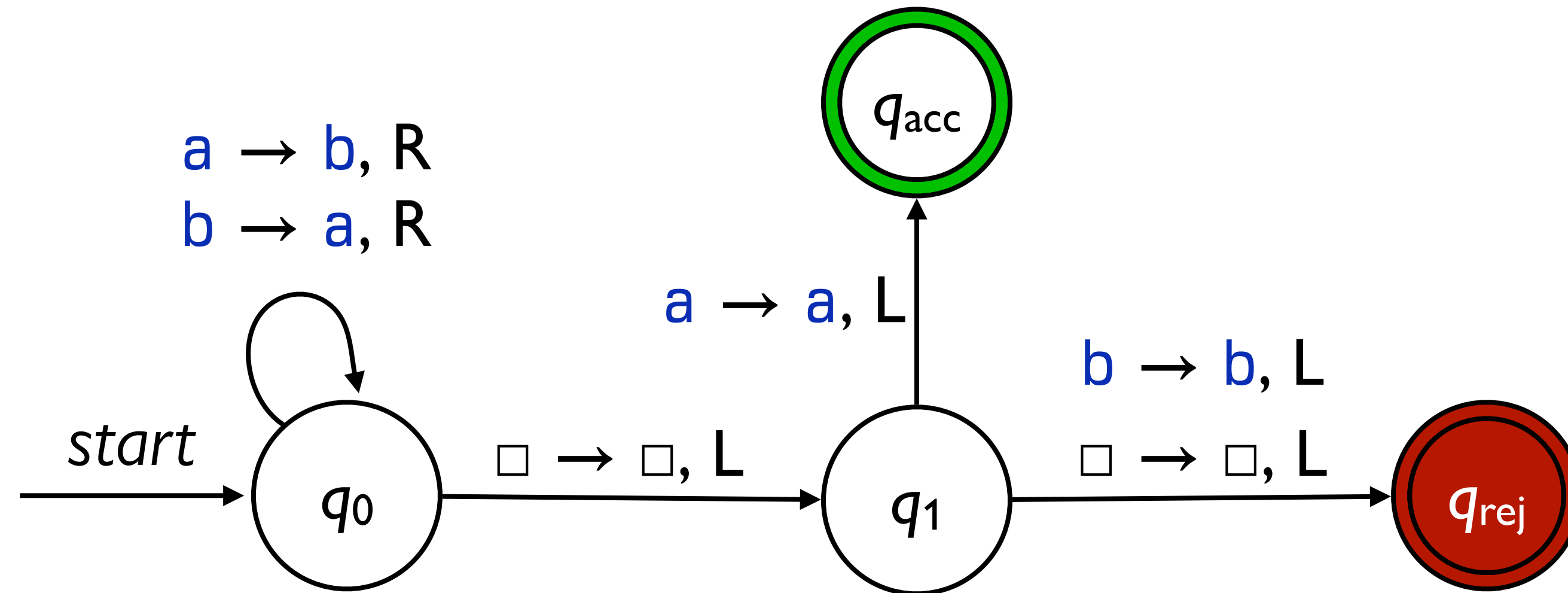


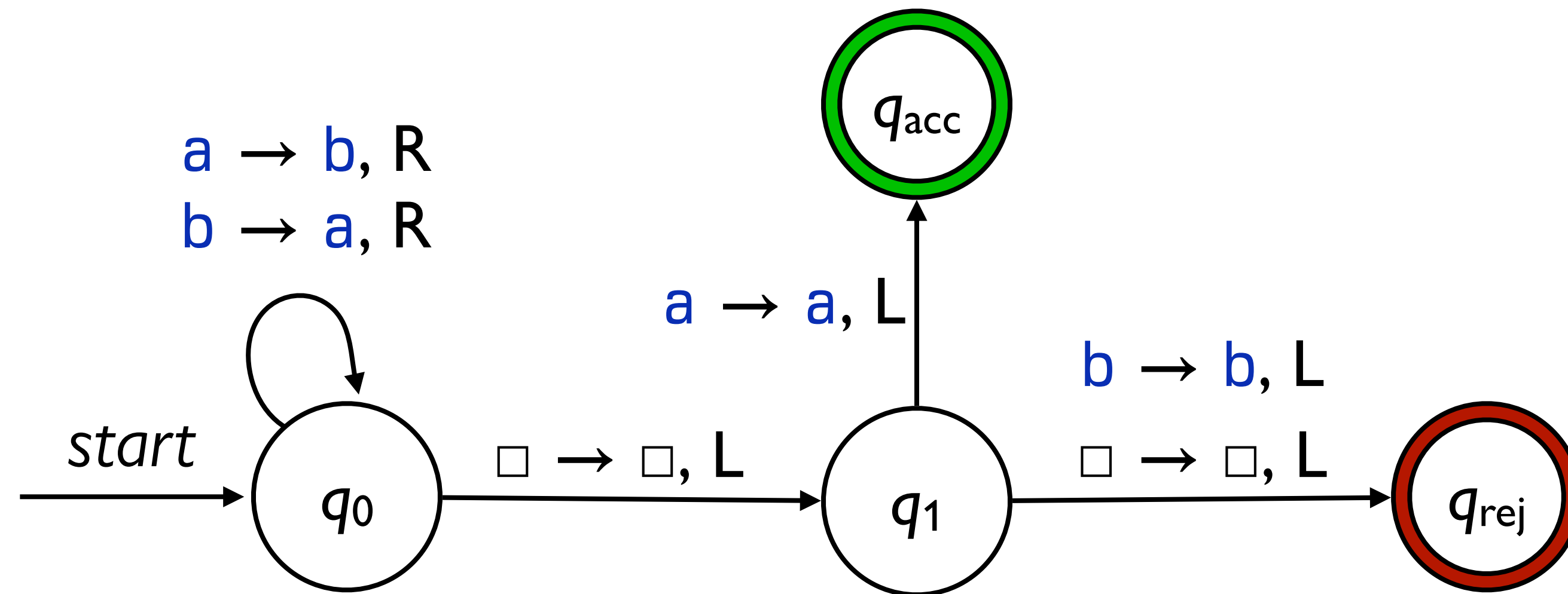








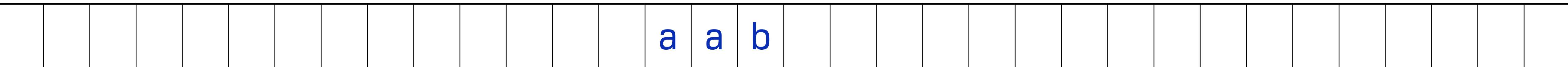
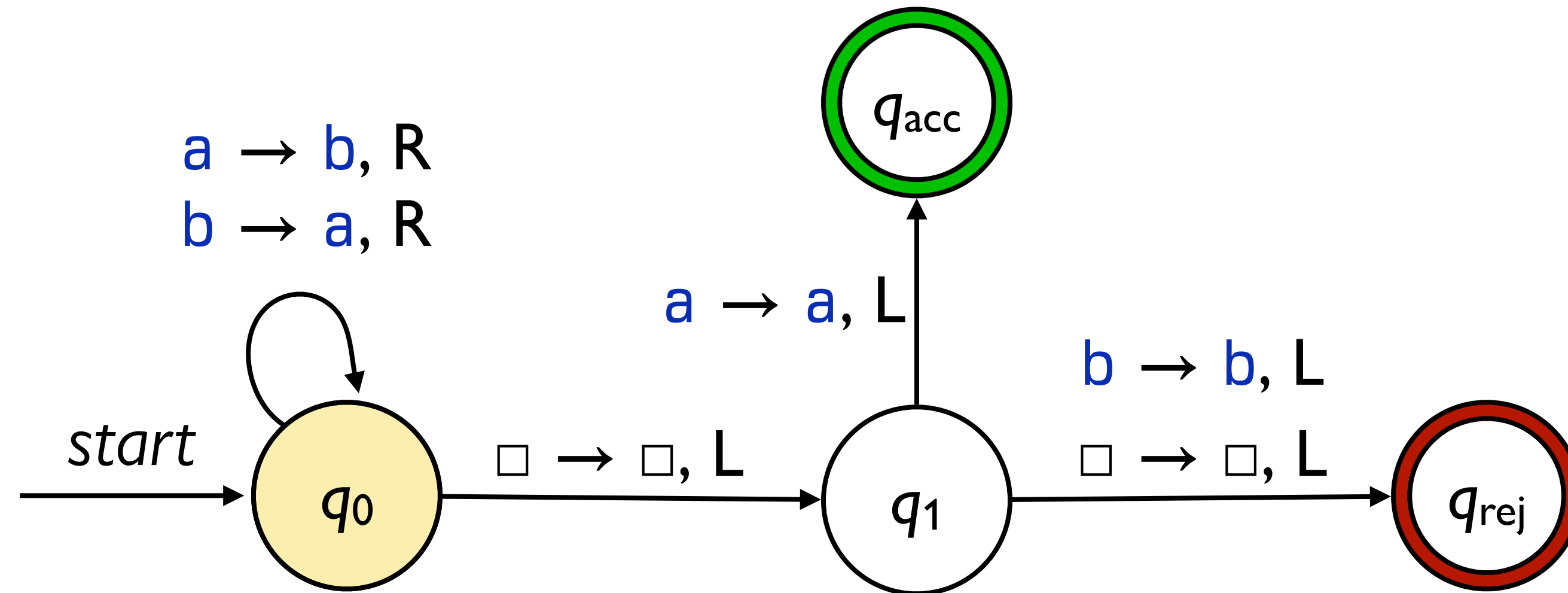


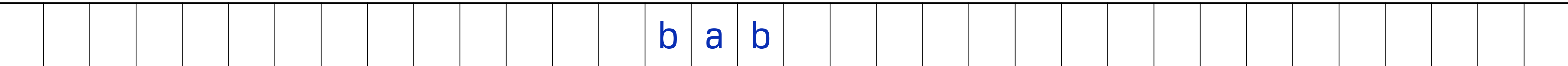
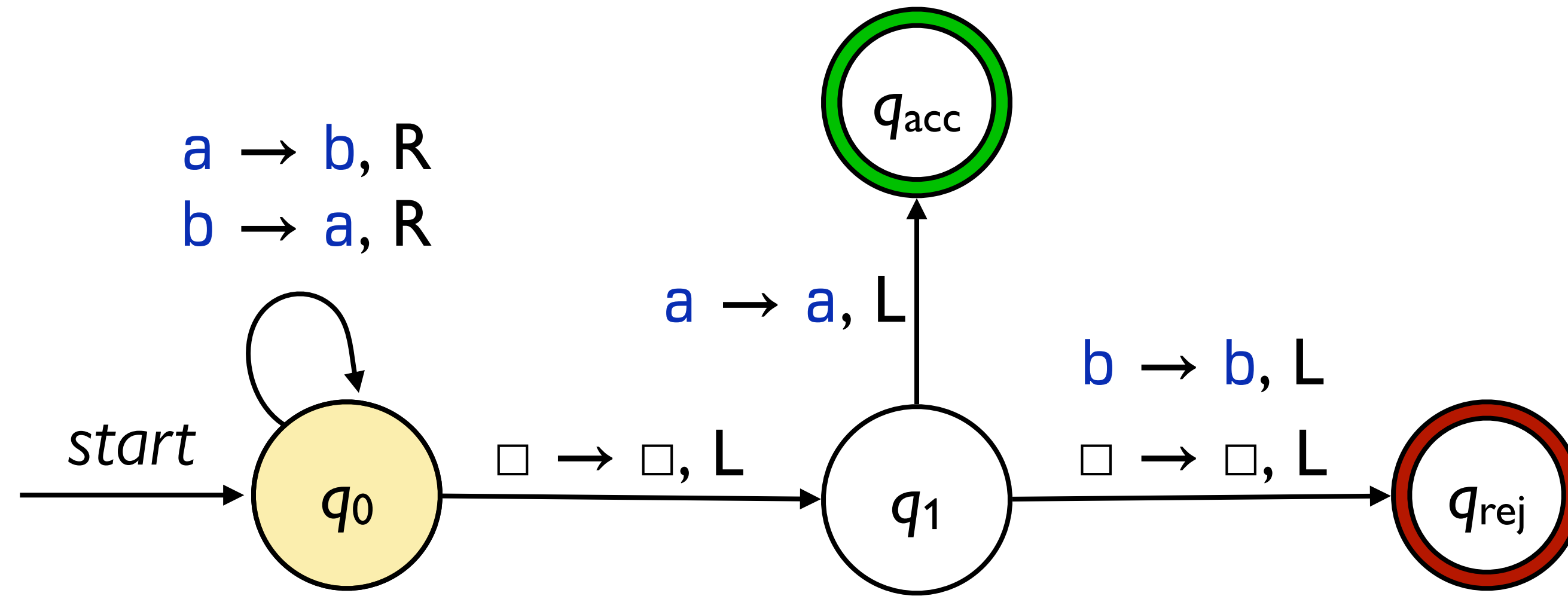


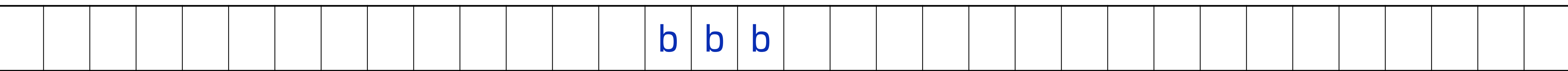
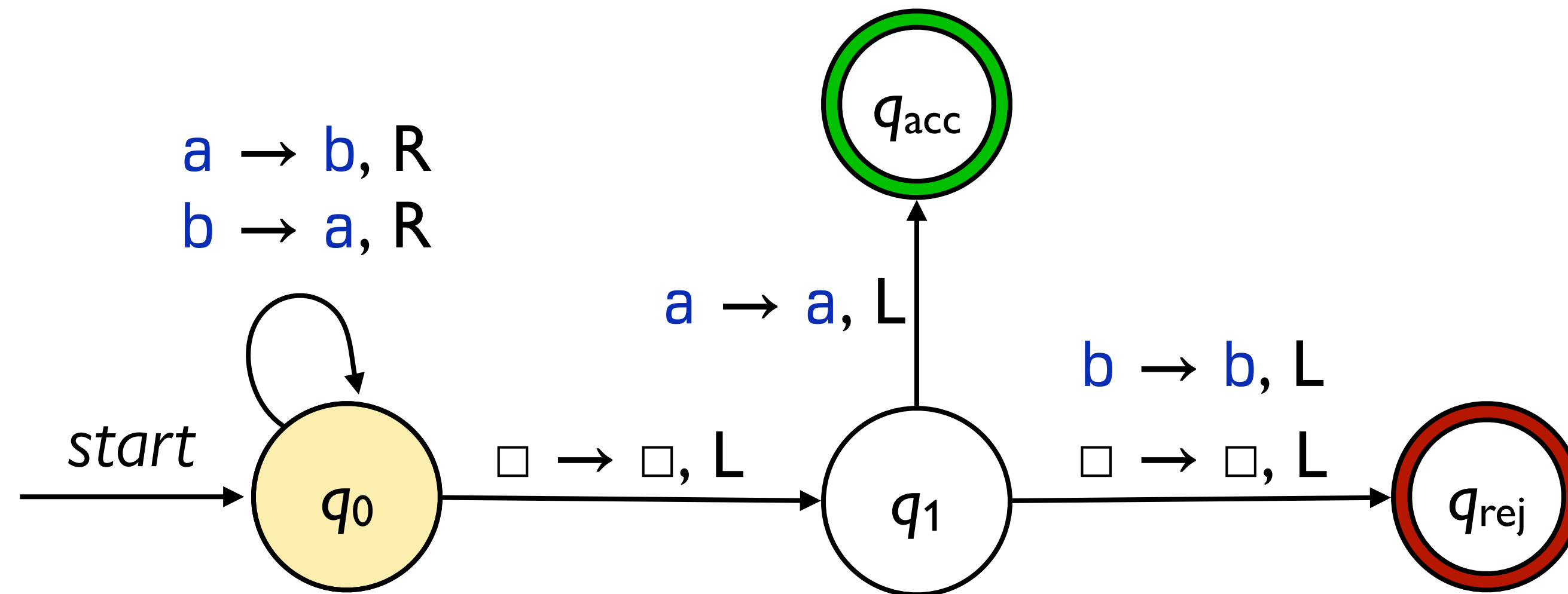
If  $M$  is a Turing machine with input alphabet  $\Sigma$ , then the **language of  $M$** , denoted  $L(M)$ , is the set

$$L(M) = \{w \in \Sigma^* \mid M \text{ accepts } w\}$$

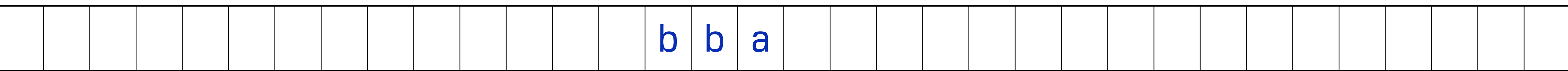
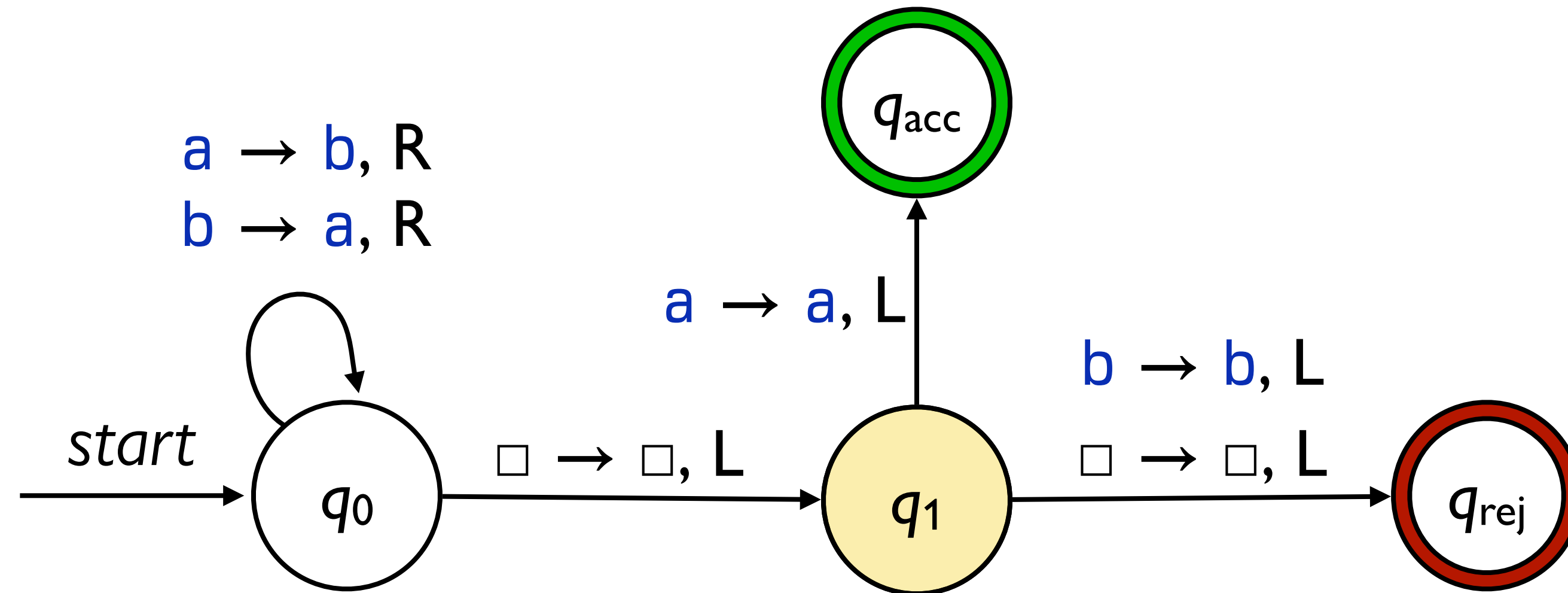
What is  $L(M)$ , where  $M$  is the above TM?

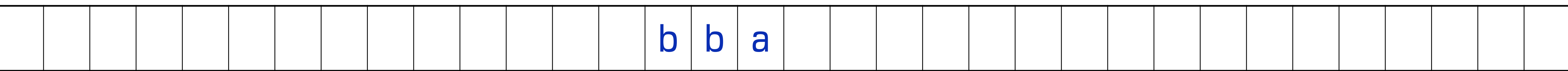
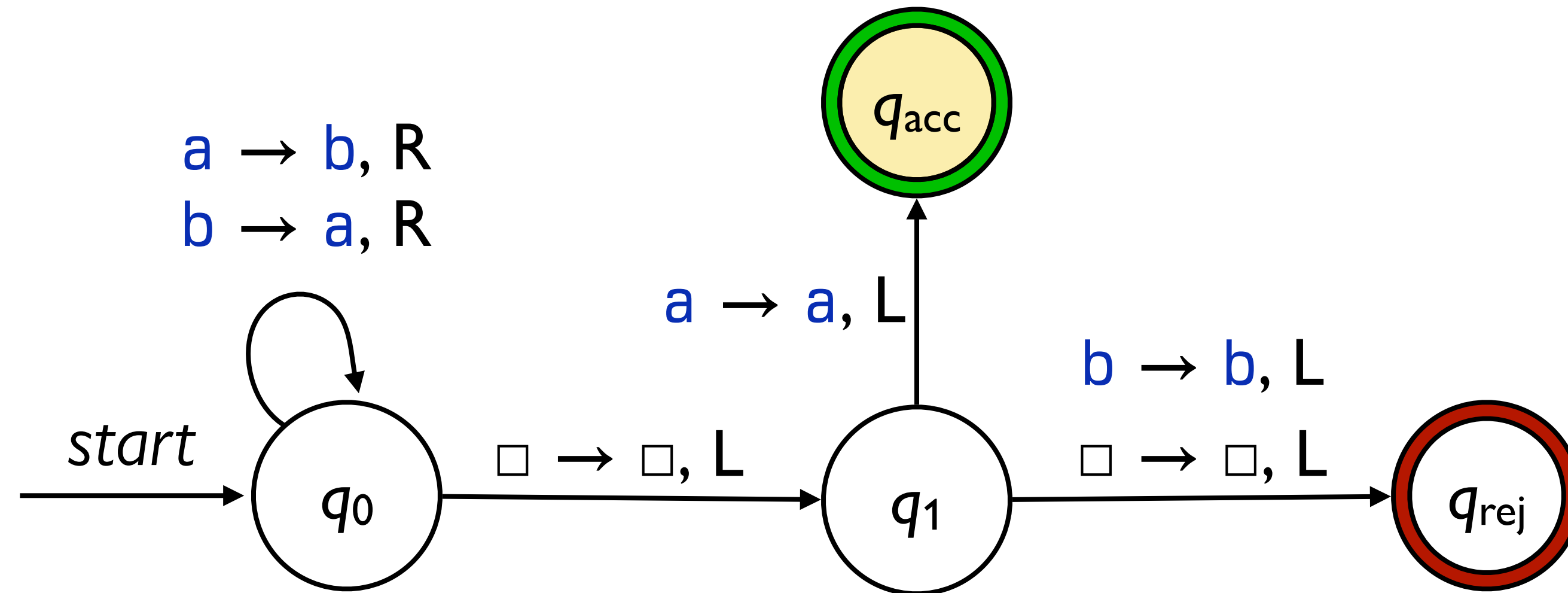
















Let  $\Sigma = \{0, 1\}$  and consider the language  $L = \{0^n 1^n \mid n \in \mathbb{N}_0\}$ .

We know that  $L$  is context-free.

How could we build a Turing machine for it?



# A recursive approach

The string  $\epsilon$  is in  $L$ .

The string  $0w1$  is in  $L$  iff  $w$  is in  $L$ .

Any string starting with  $1$  is not in  $L$ .

Any string ending with  $0$  is not in  $L$ .









































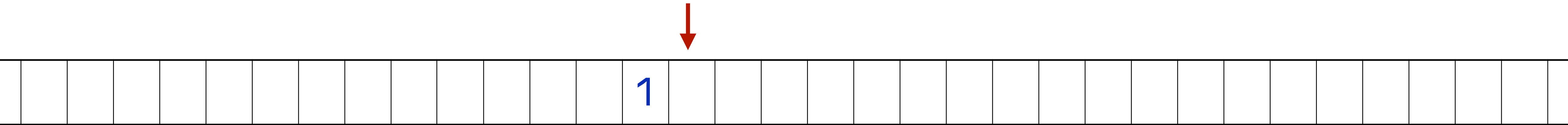








# A sketch of the TM

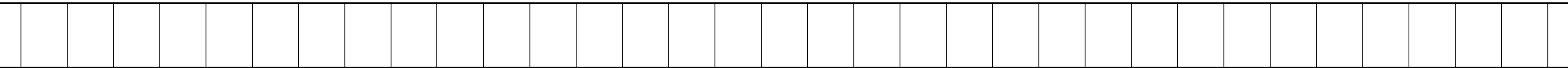




# A sketch of the TM



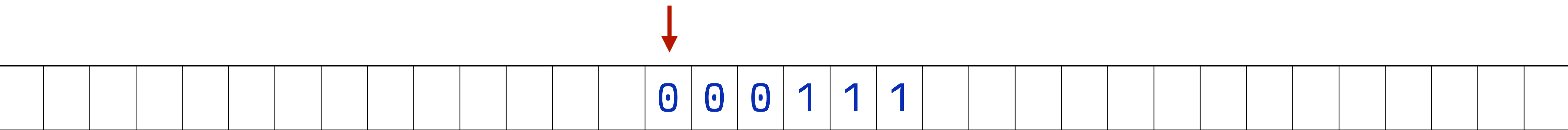
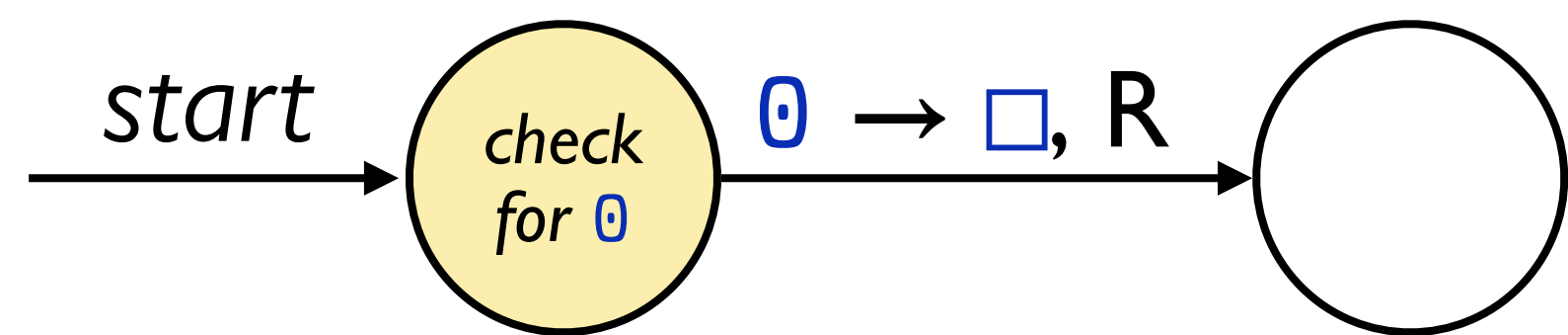
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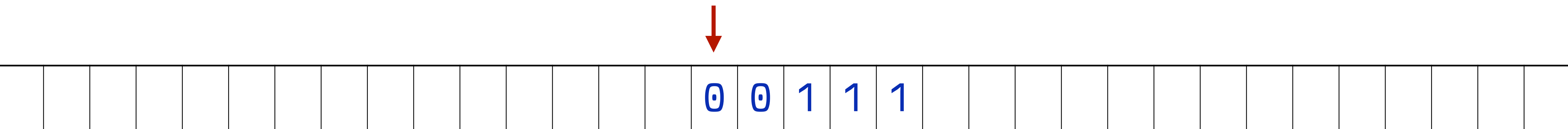
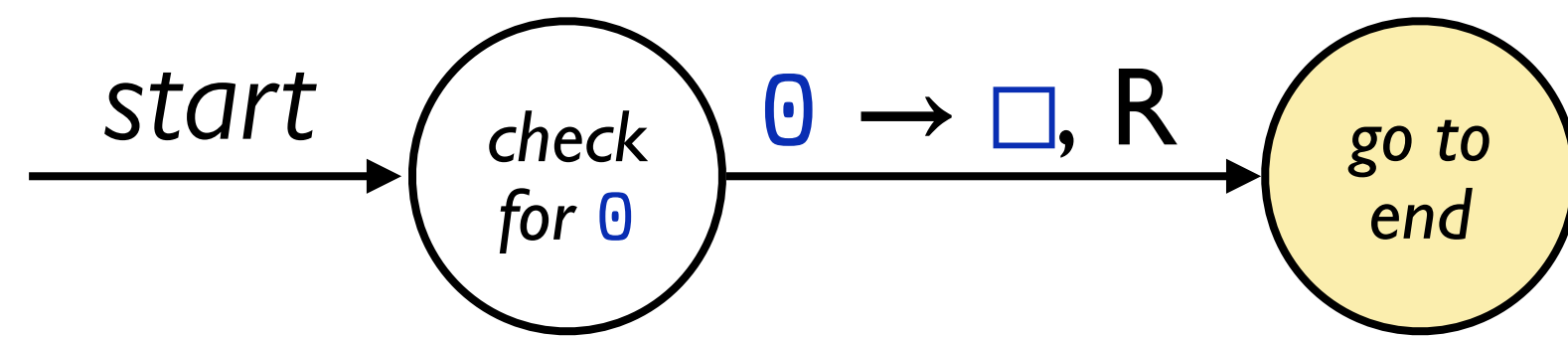










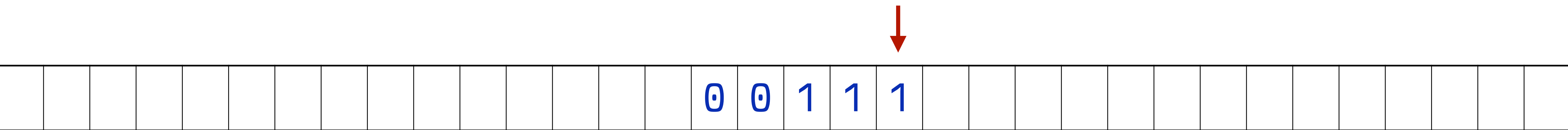
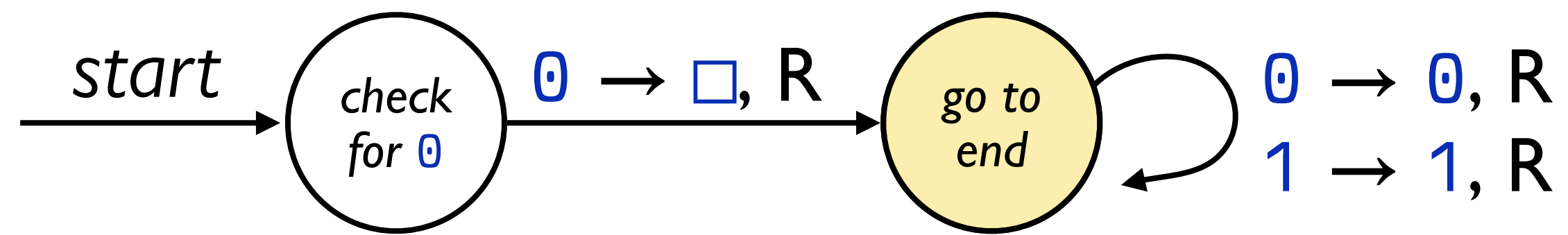






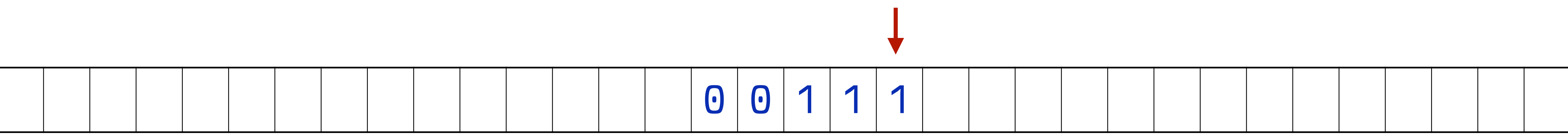
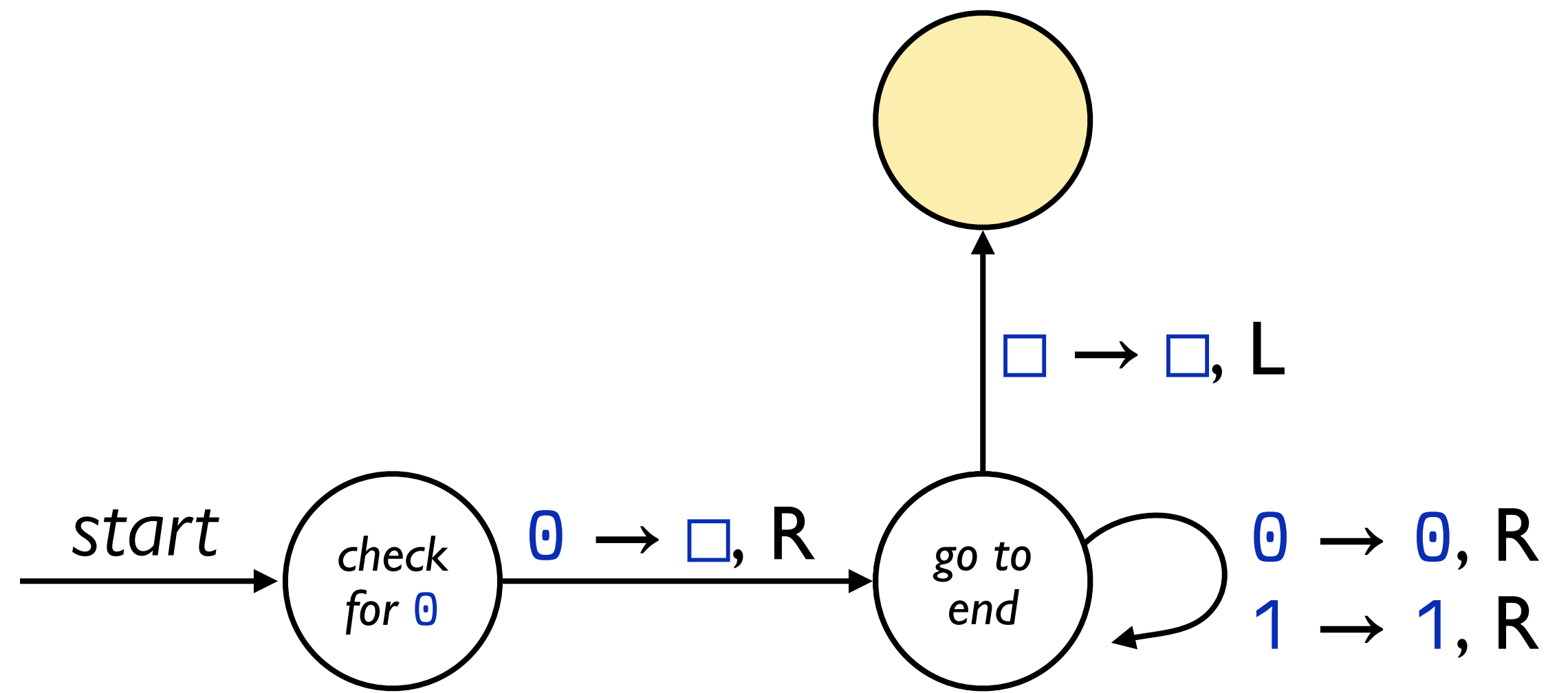


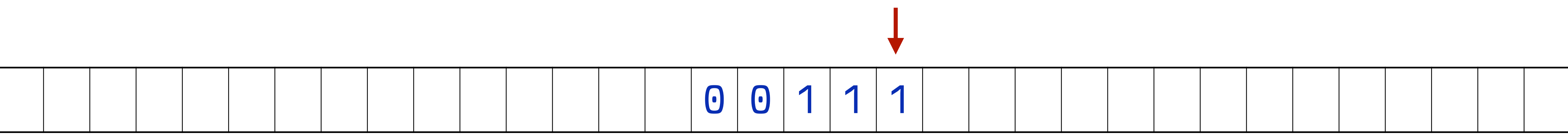
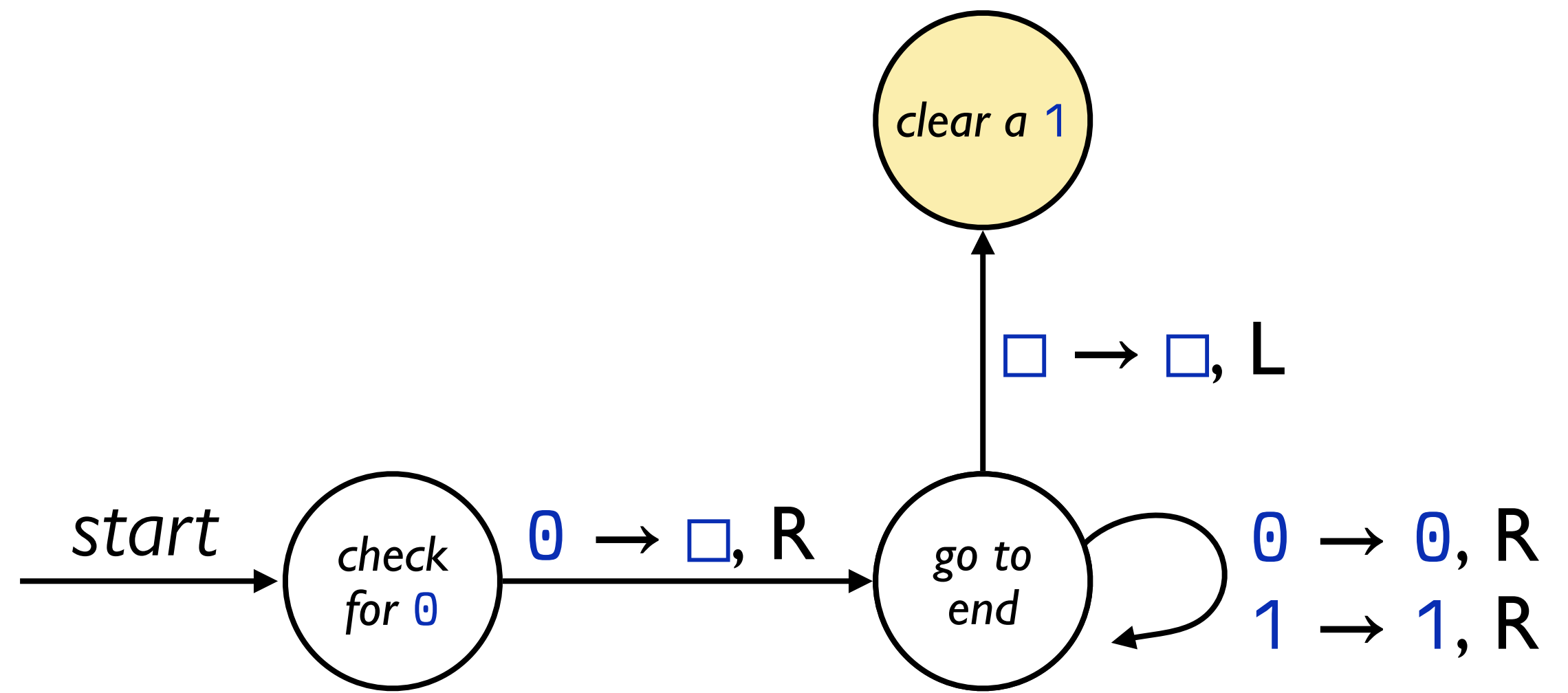




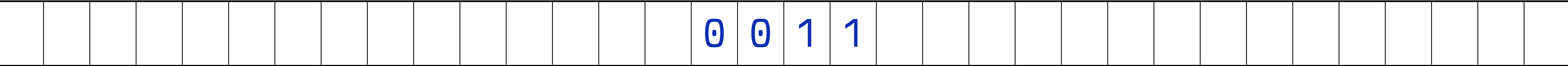
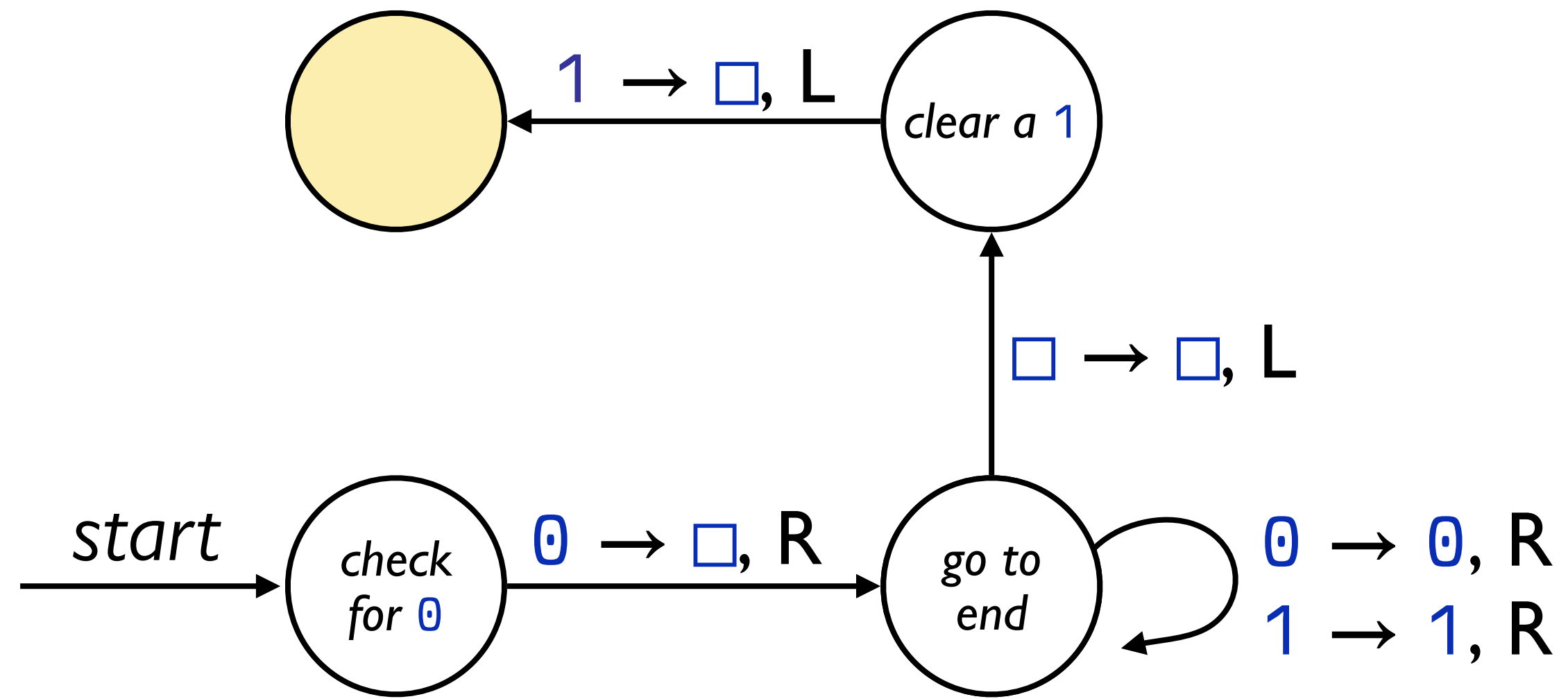


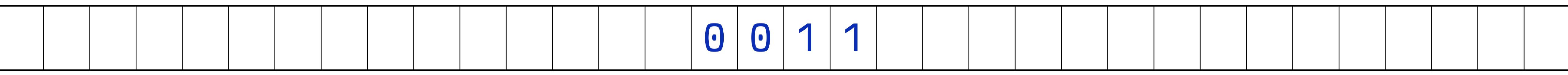
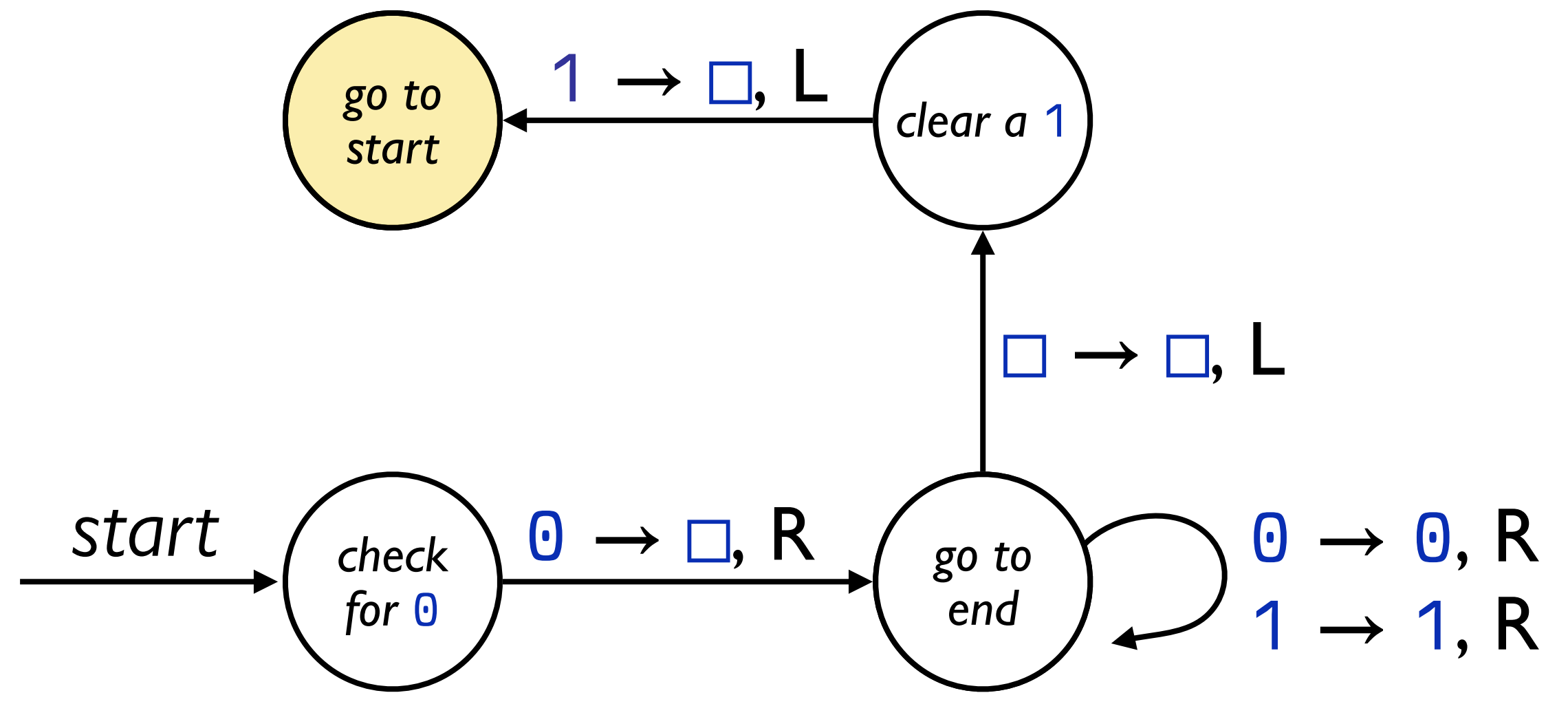
















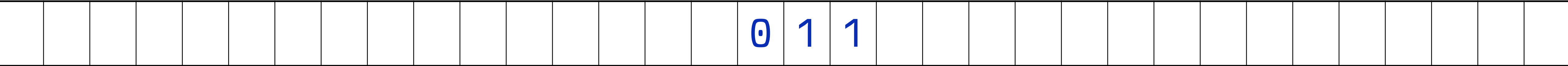
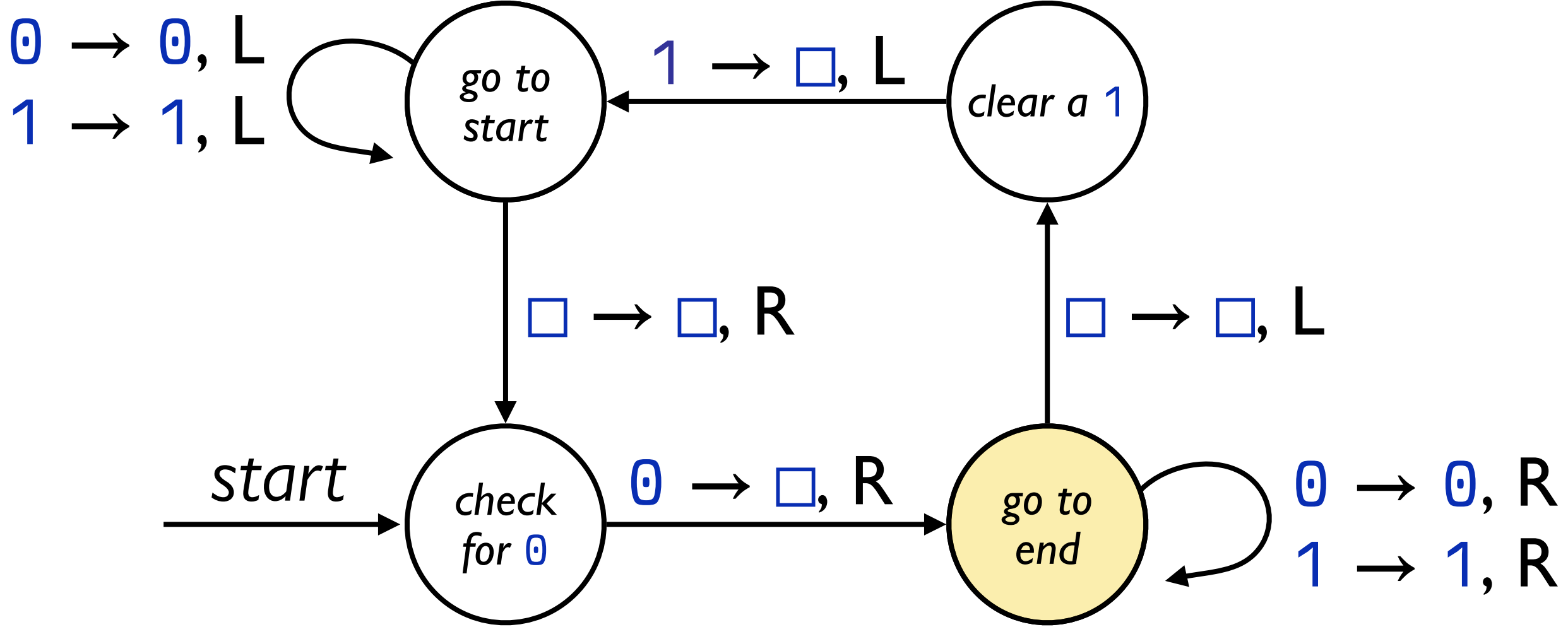










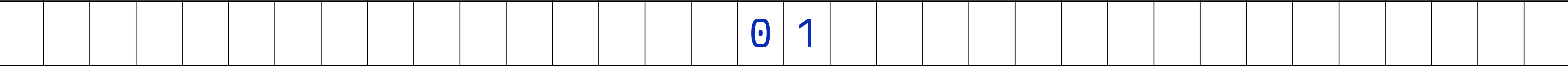
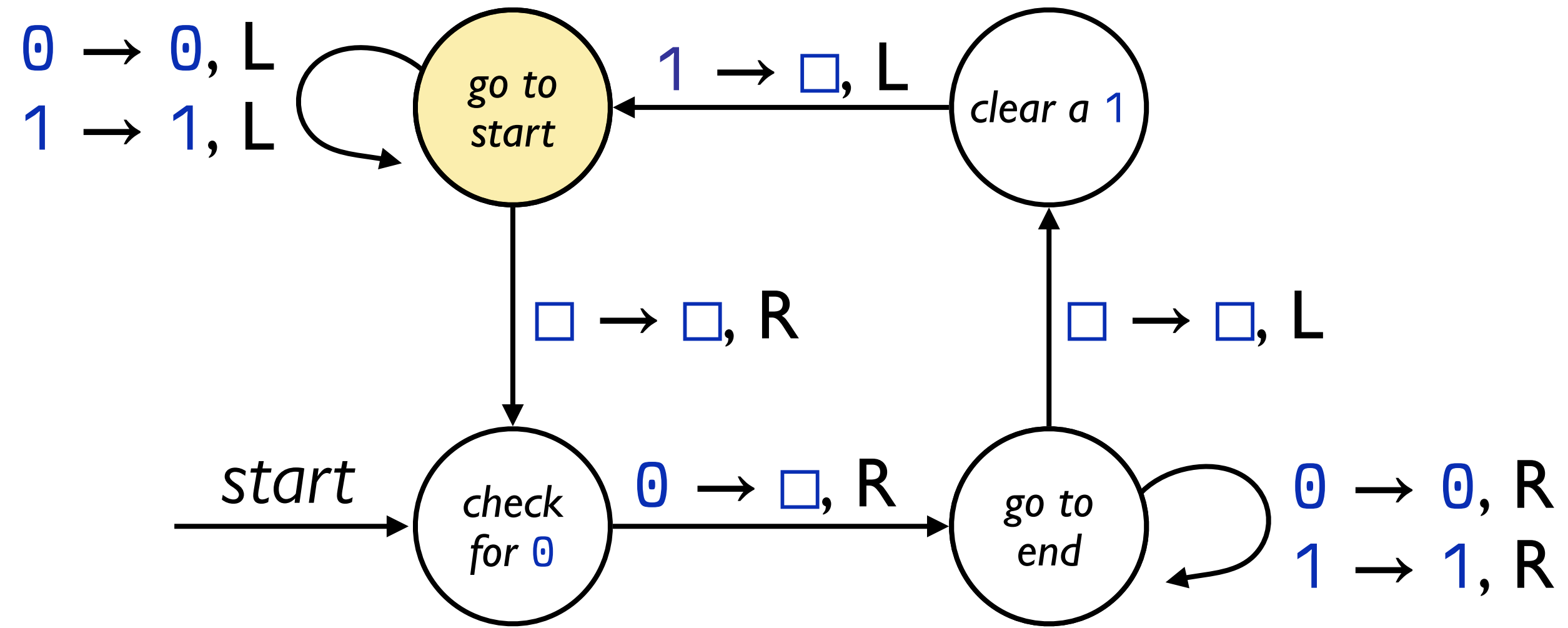




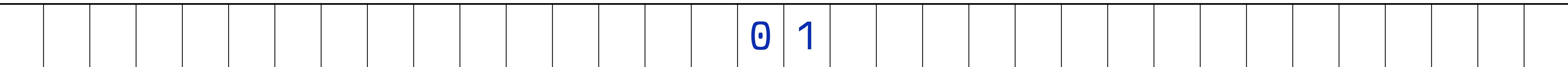
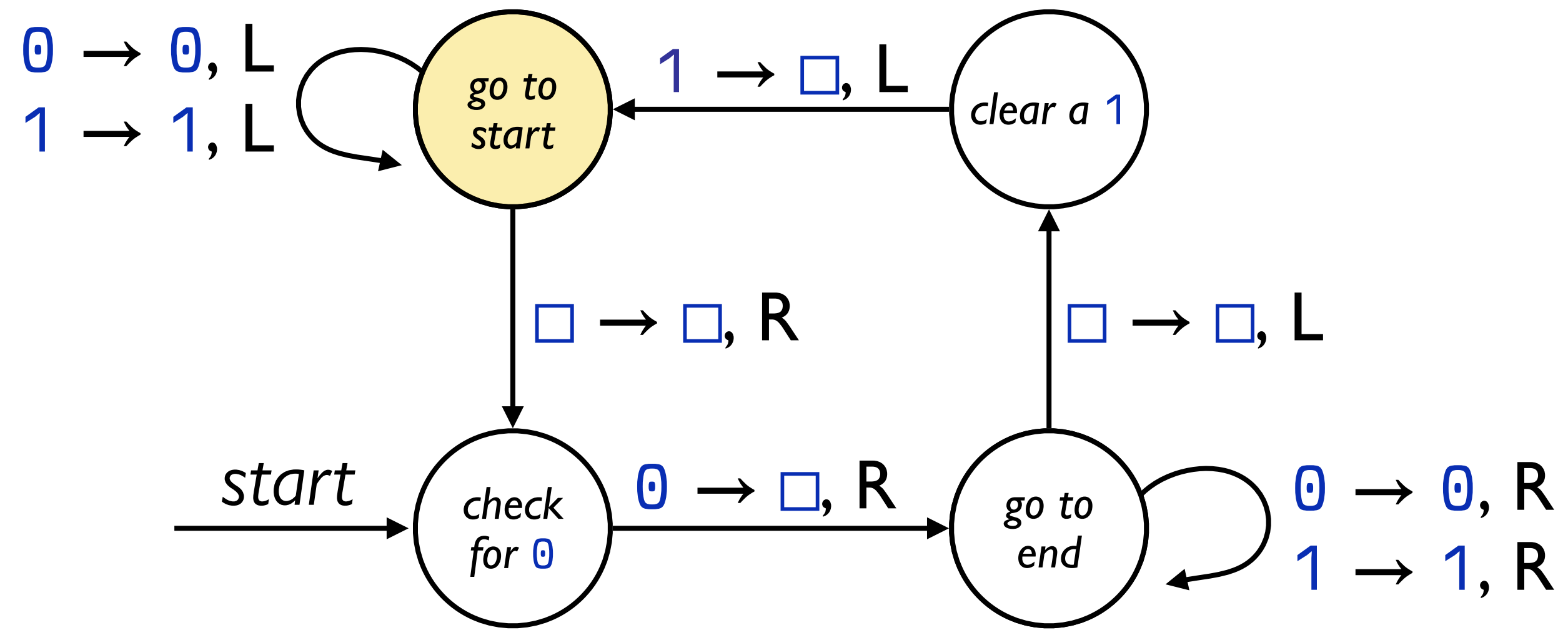




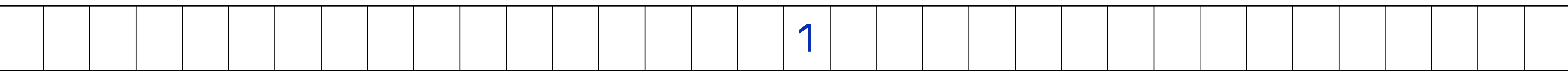
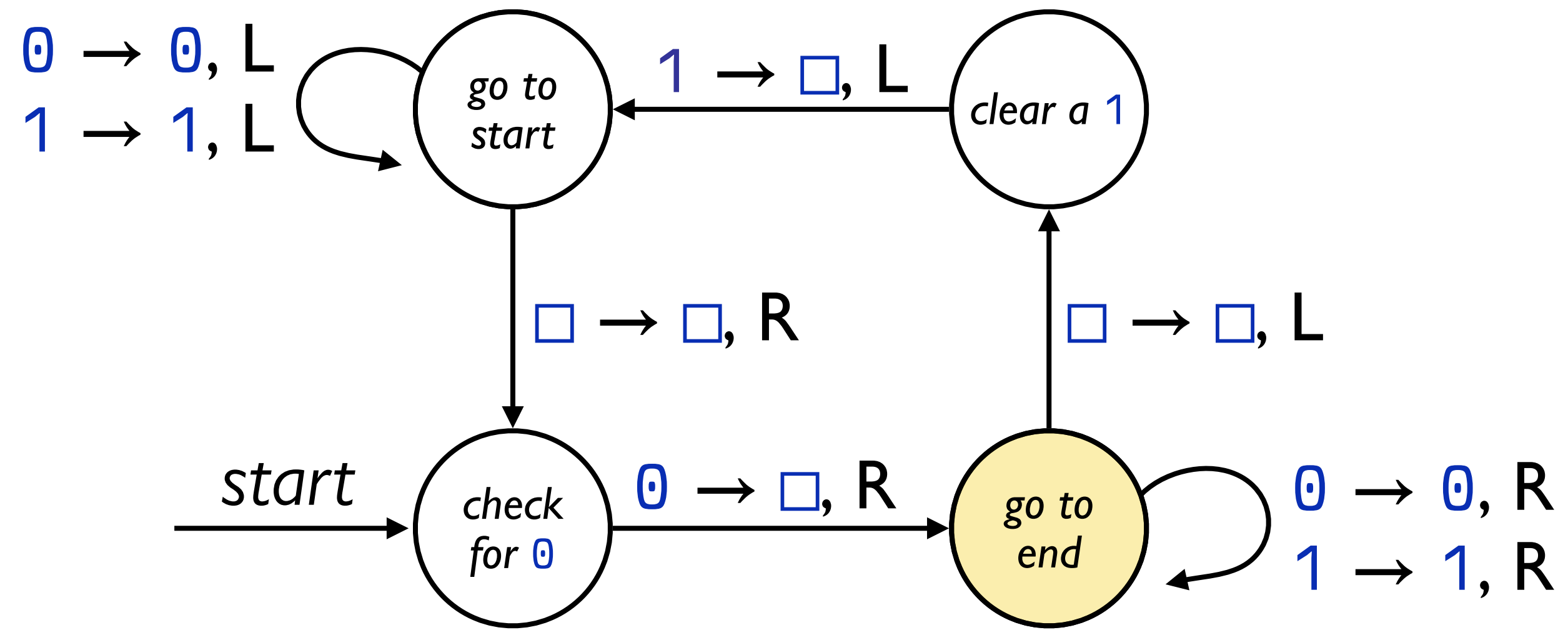


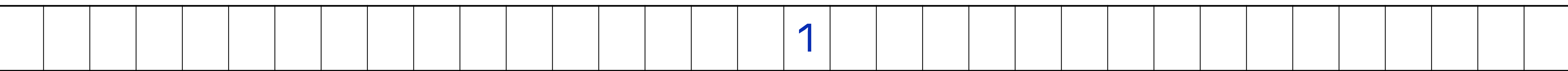
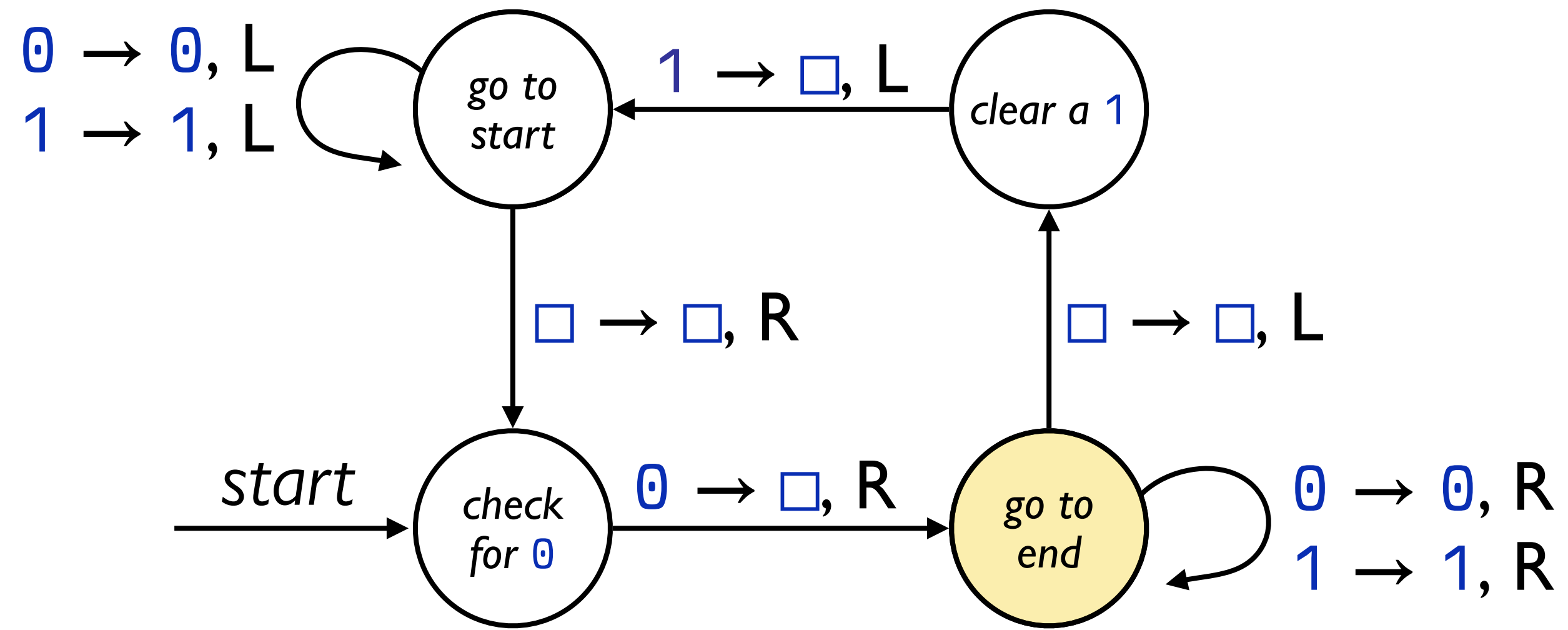


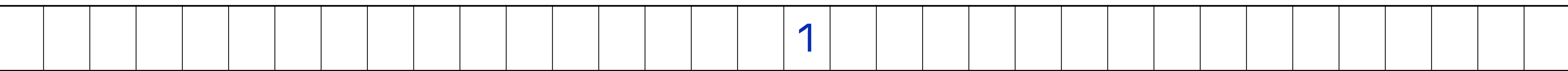
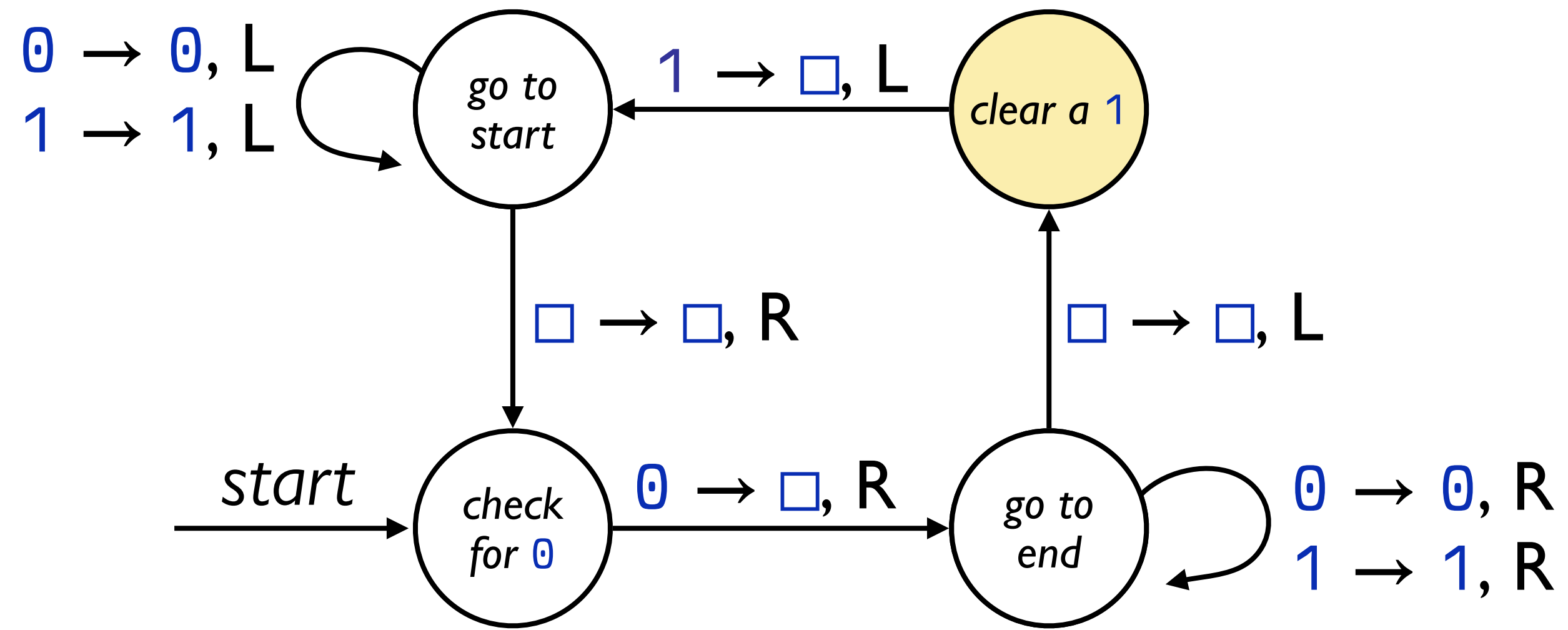


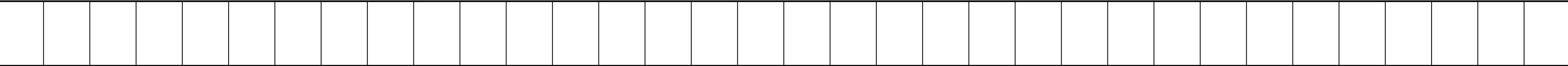
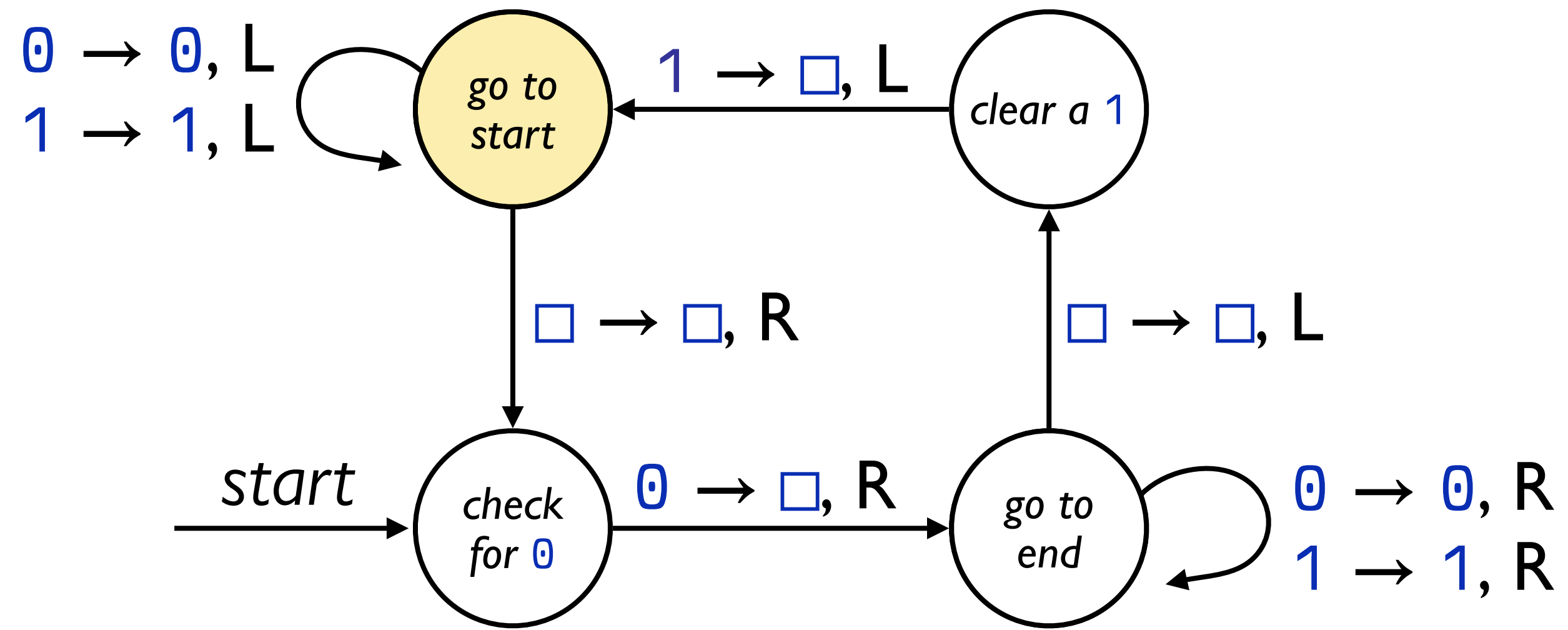


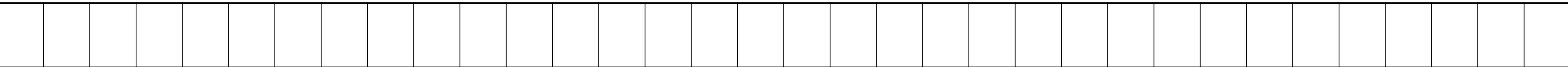
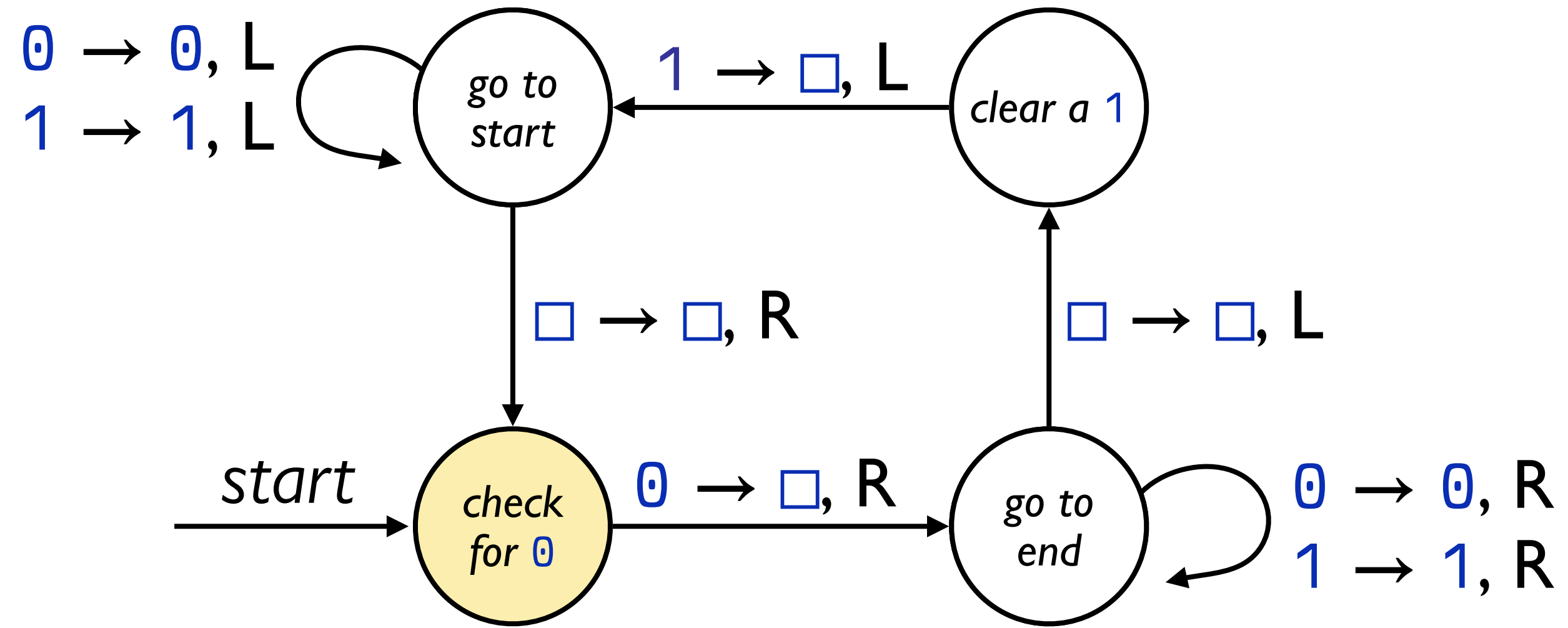


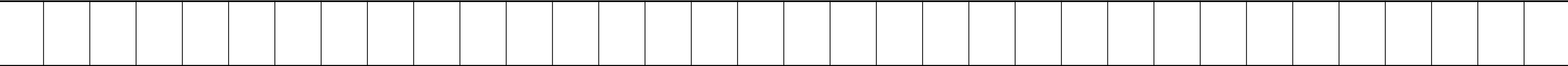
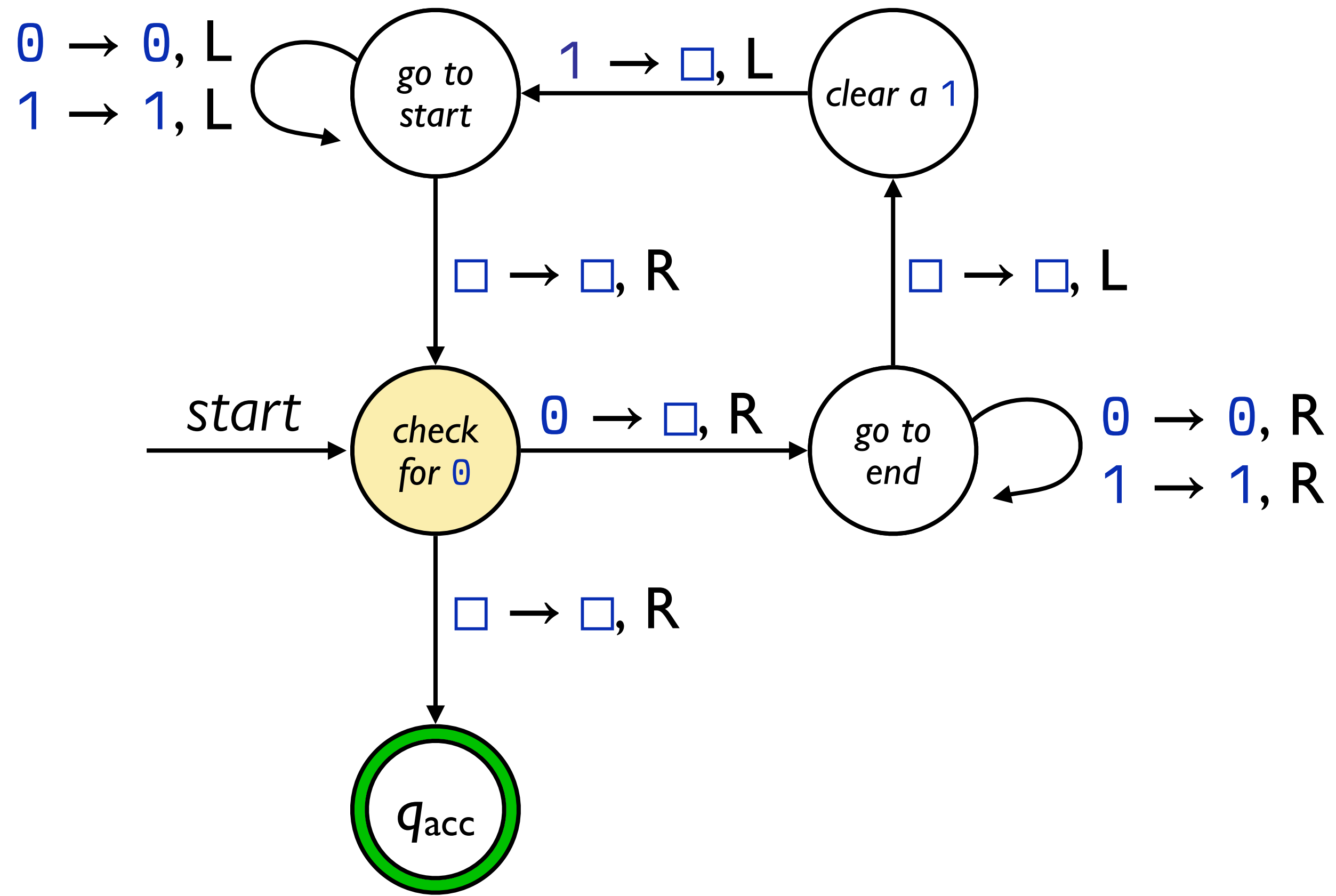


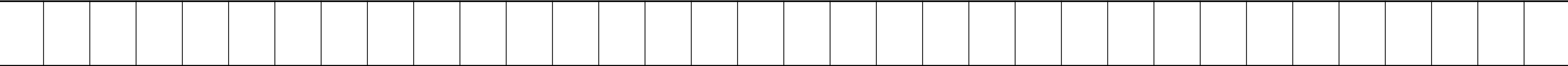
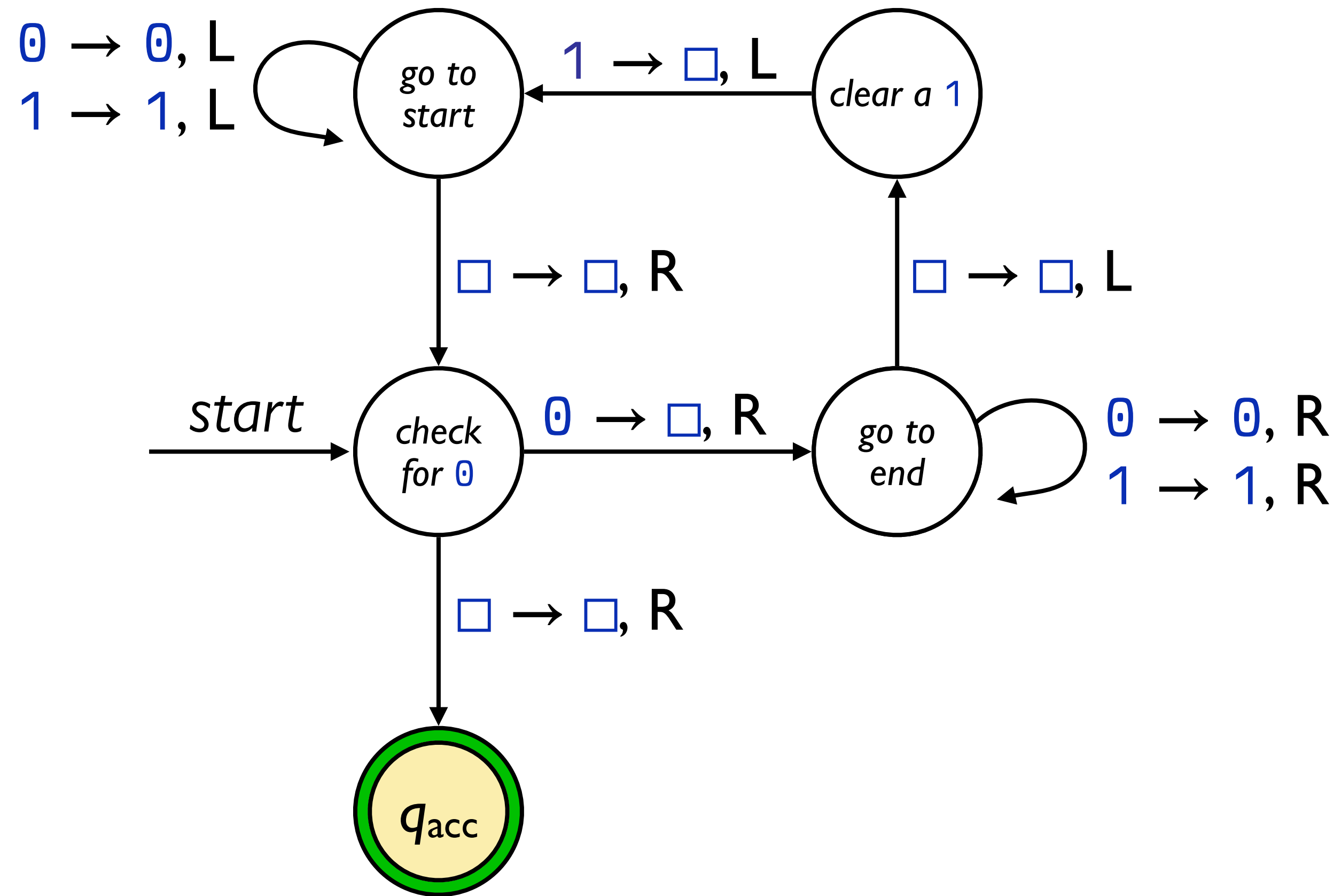


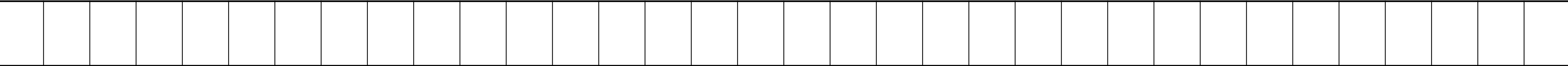
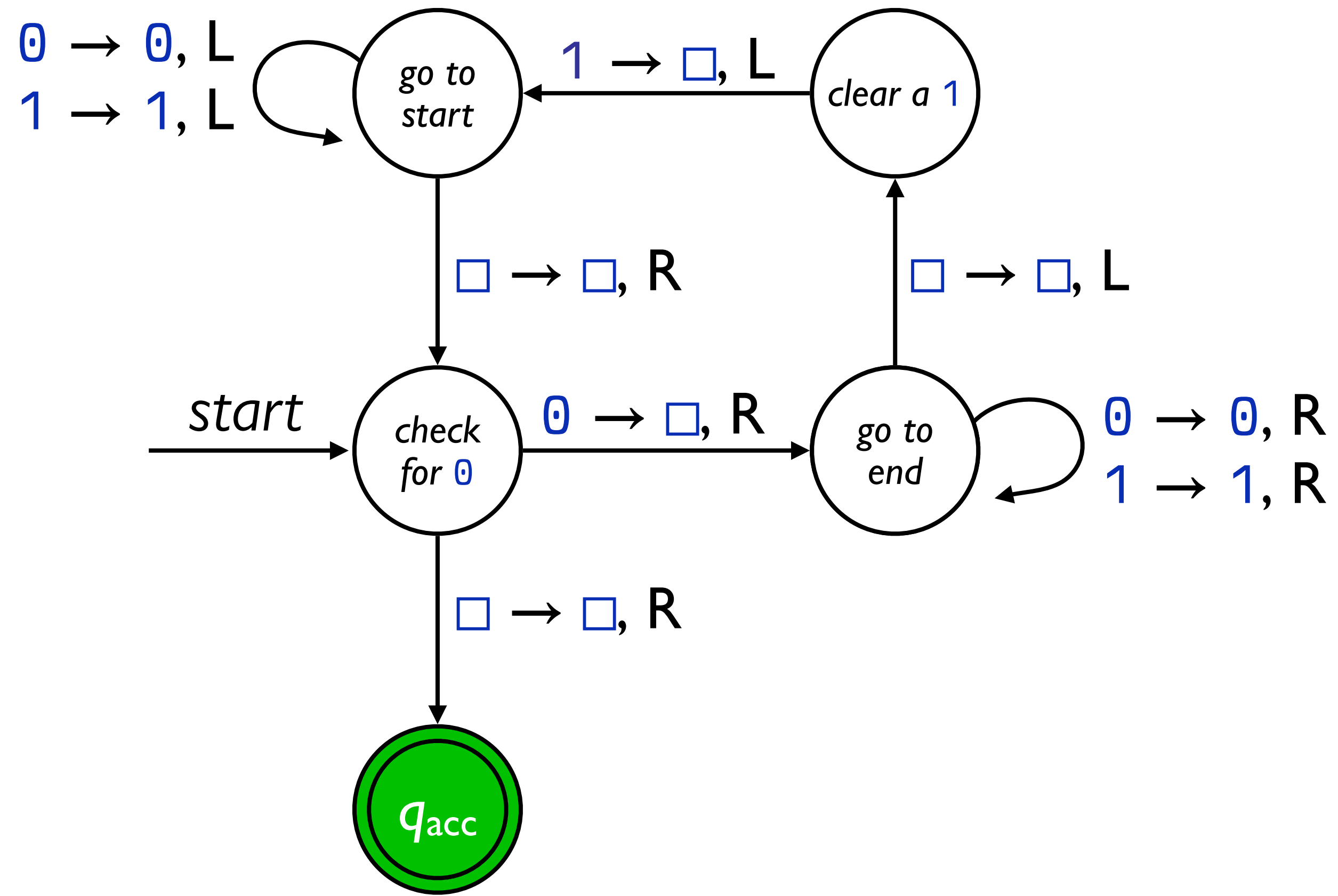


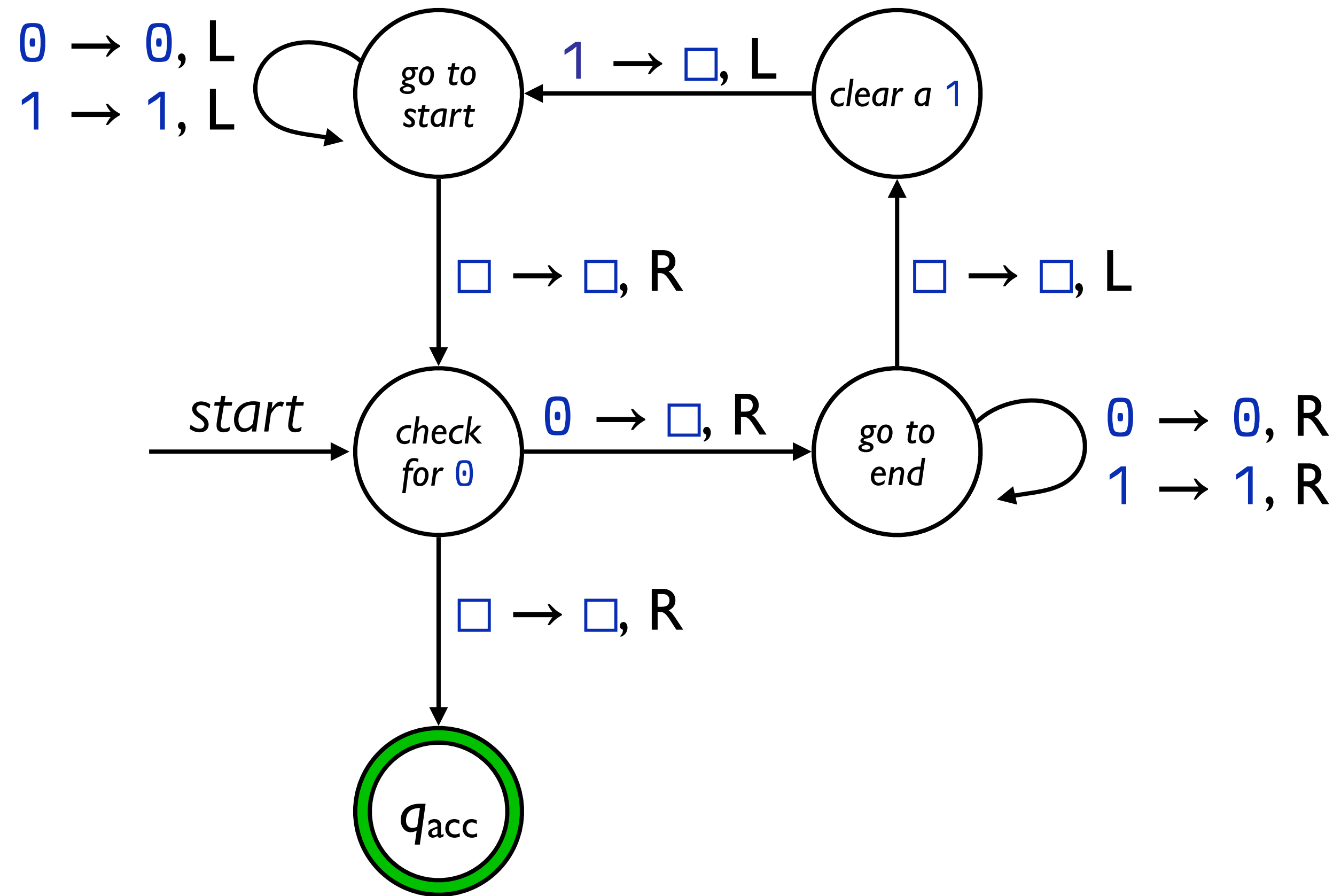


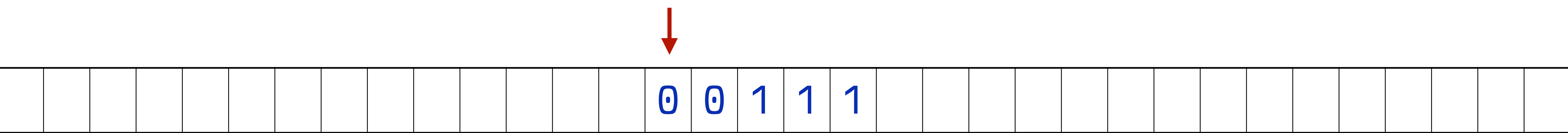
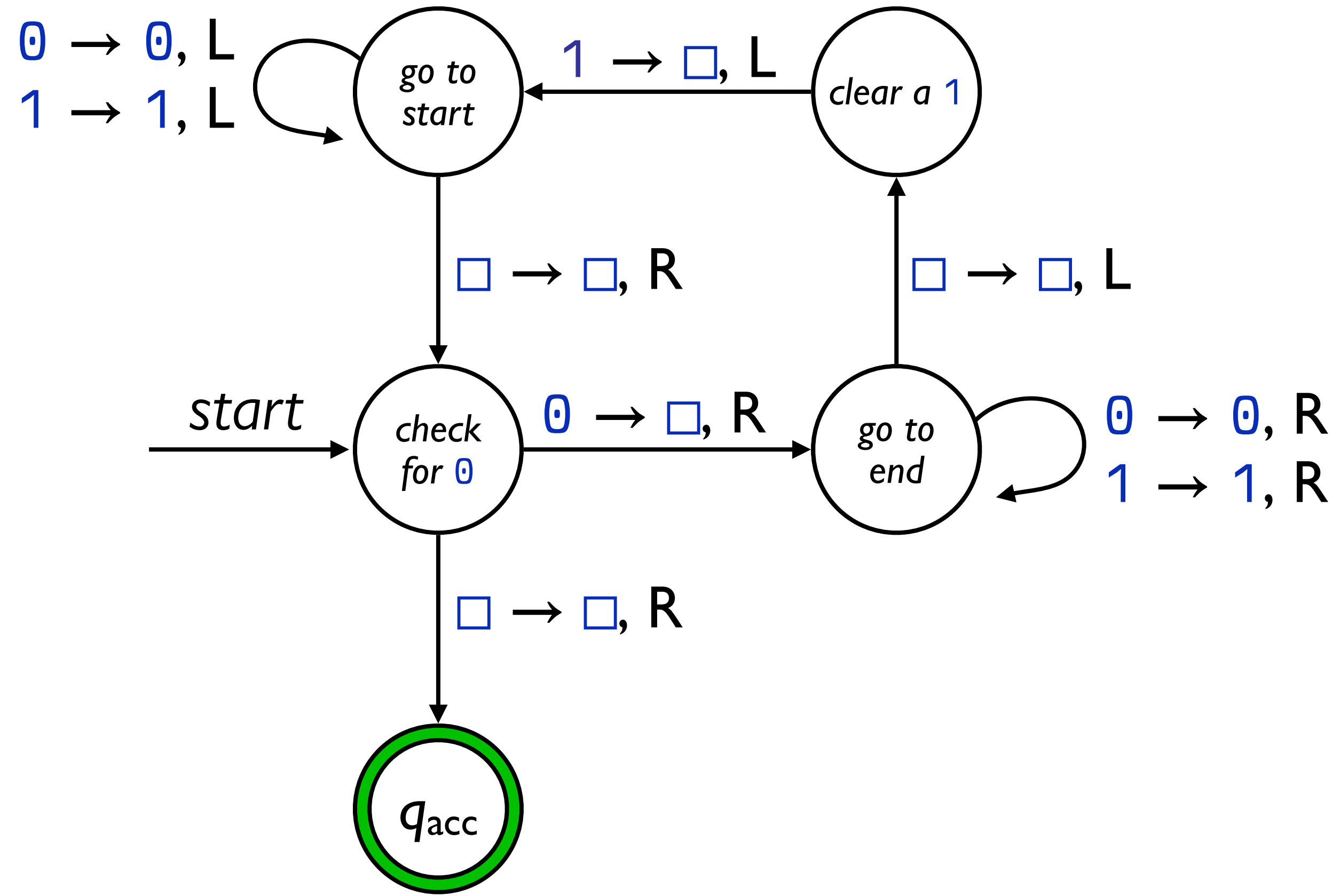




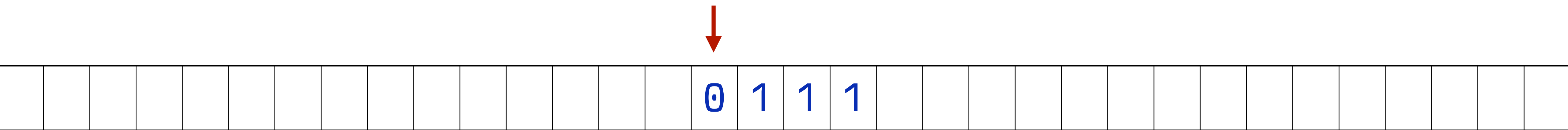
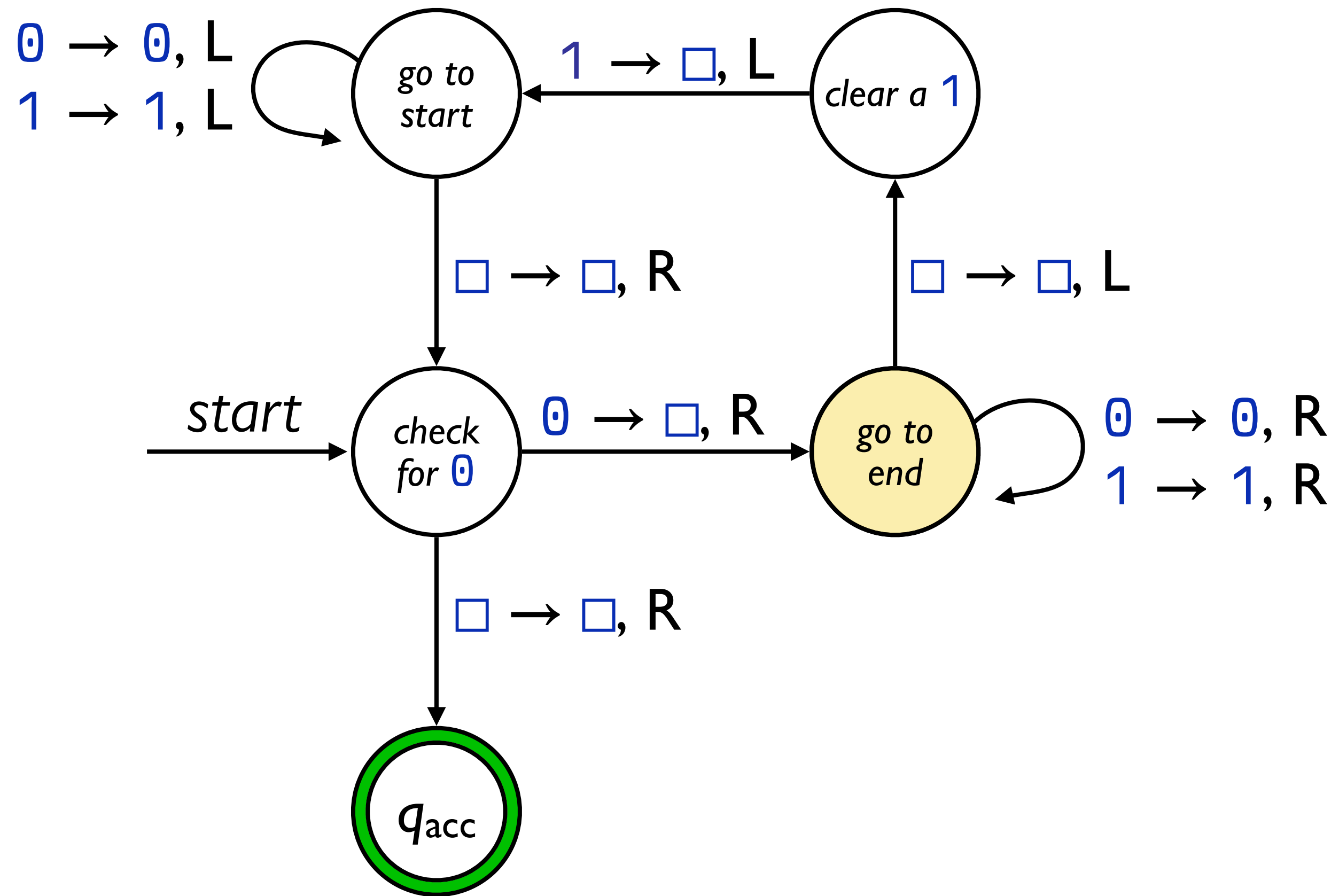


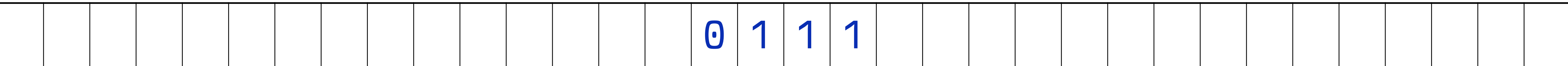
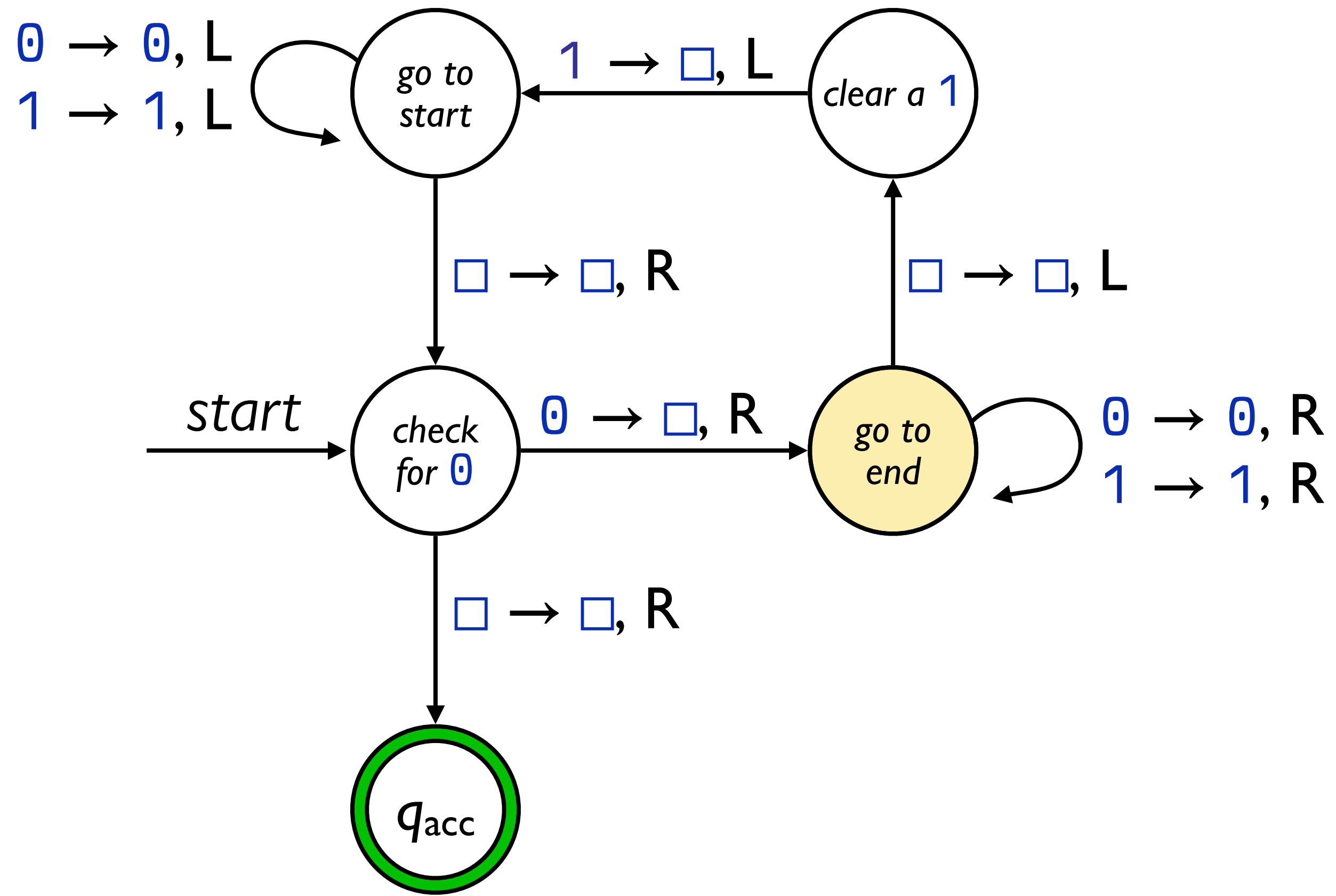


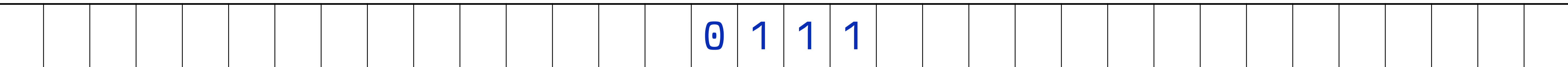
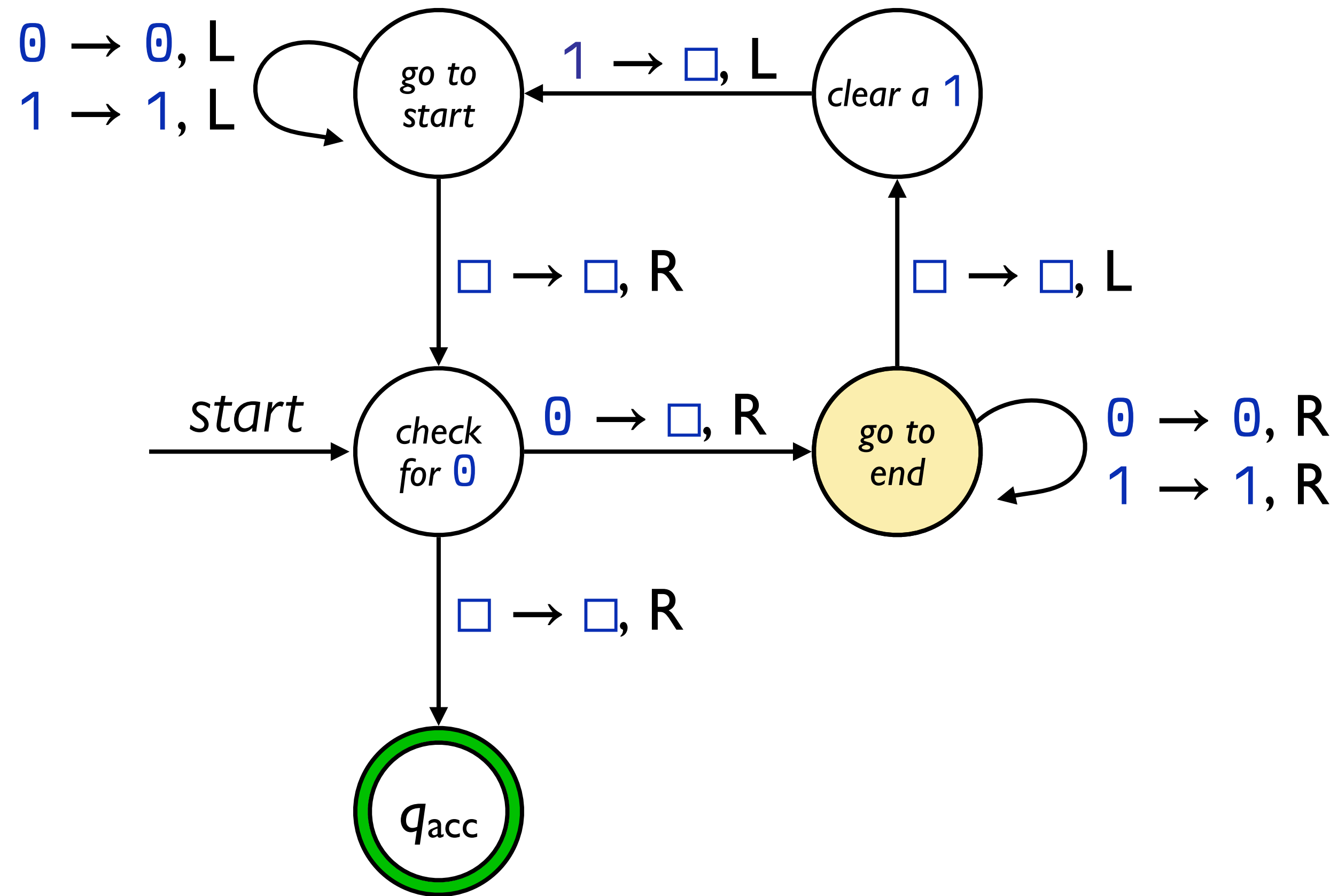


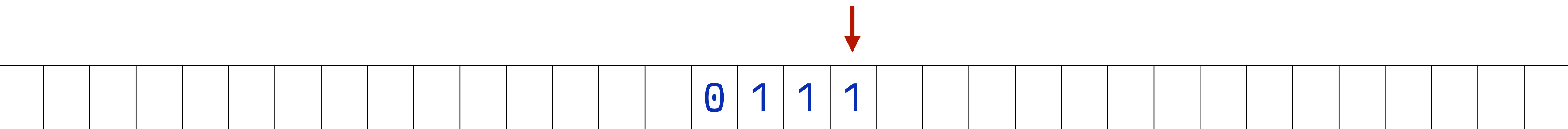
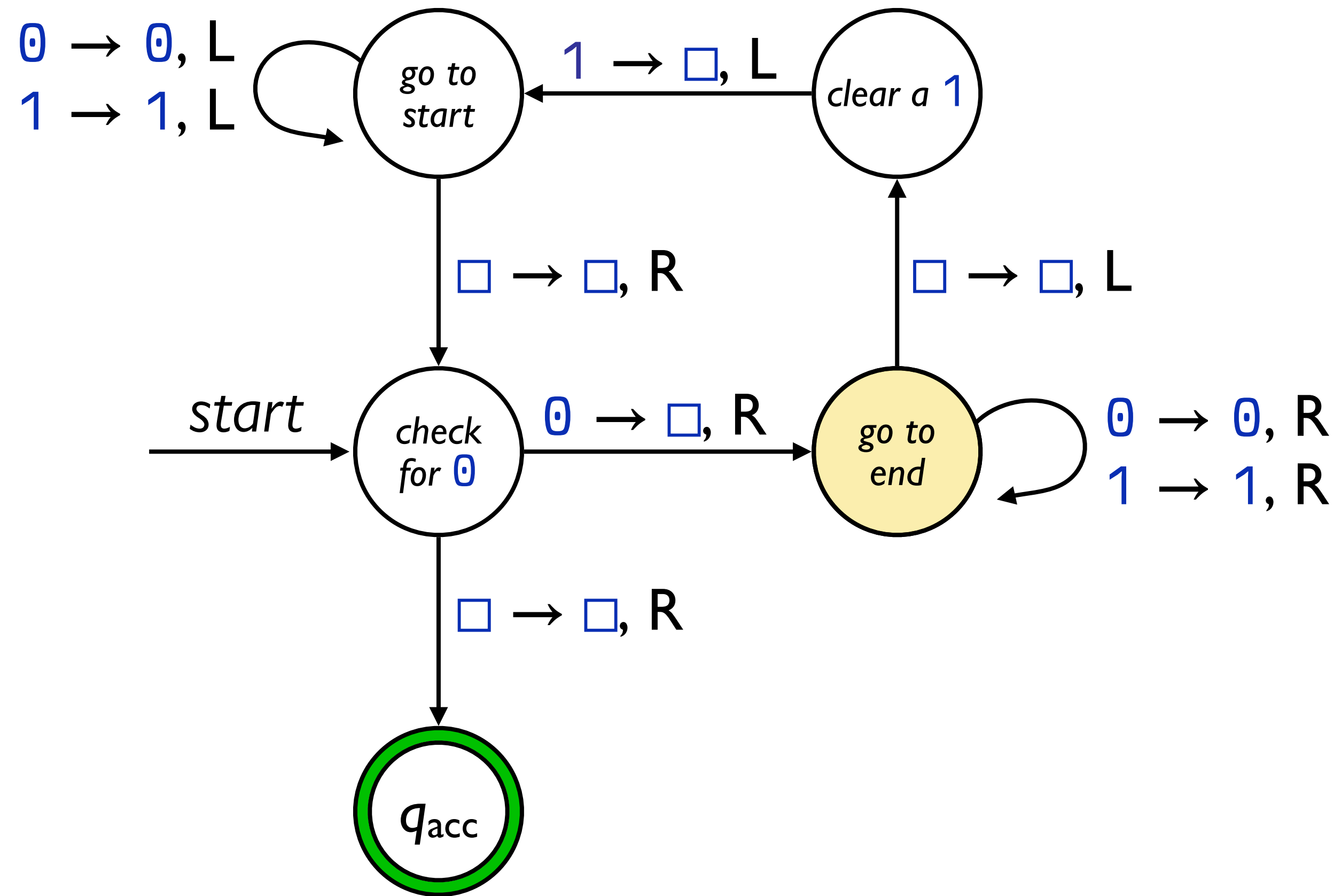


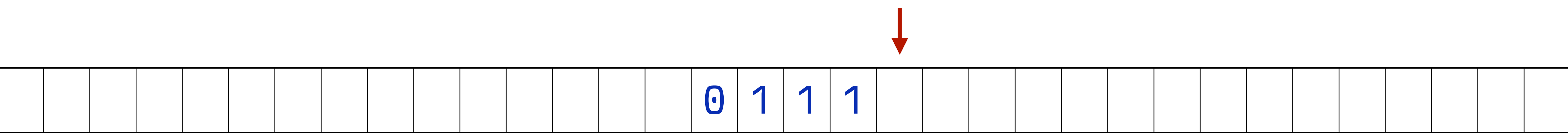
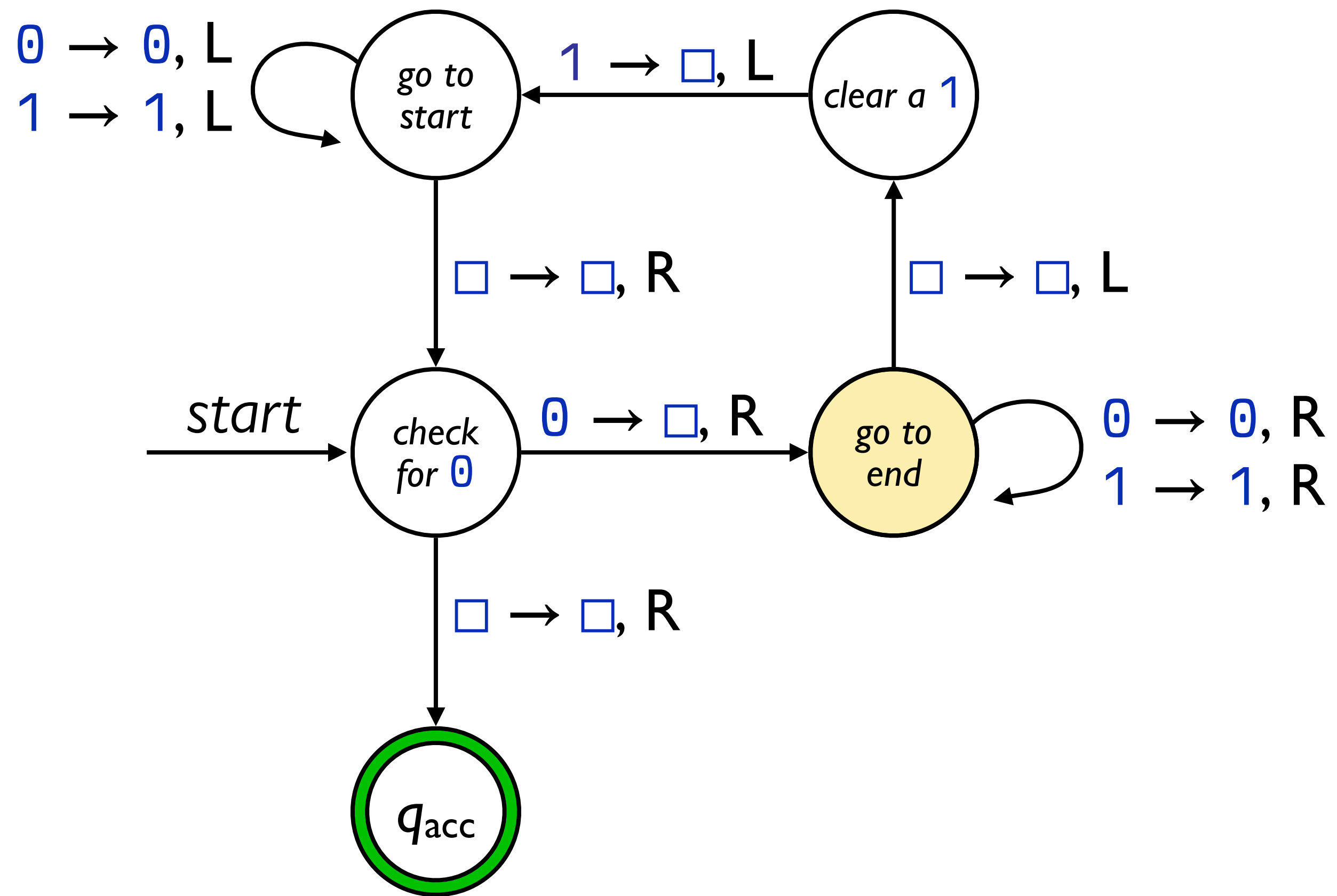


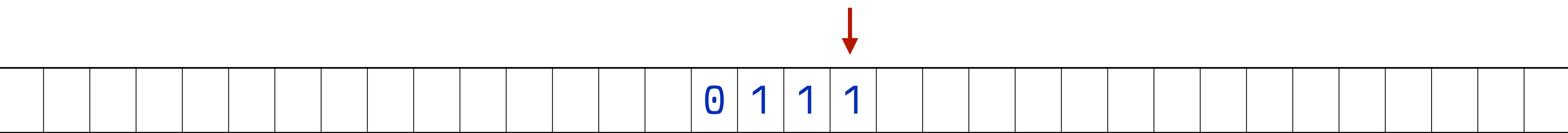
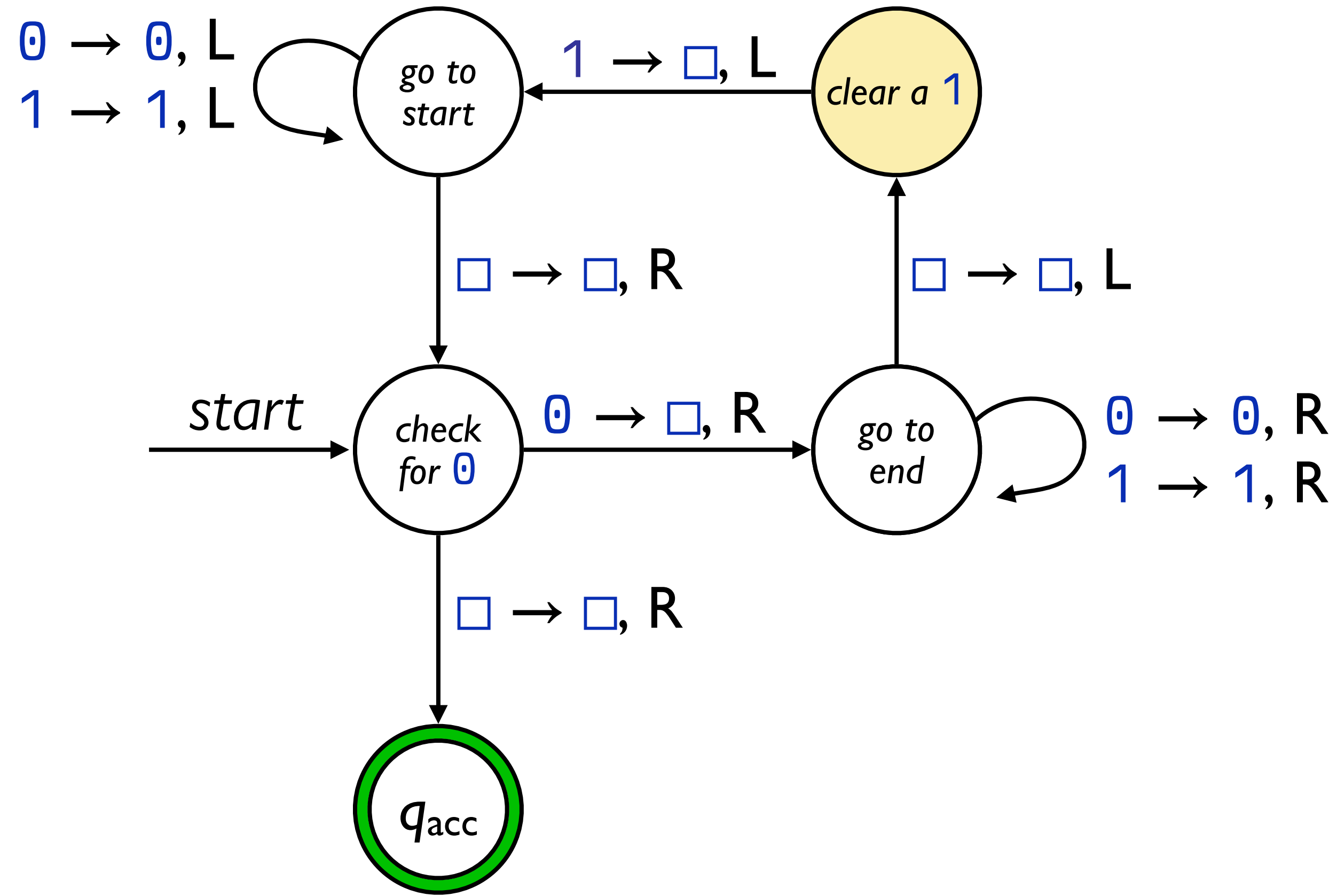






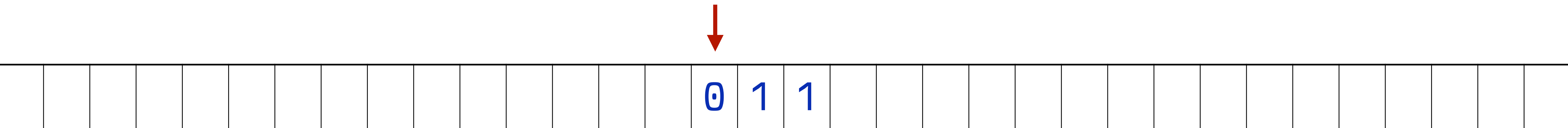
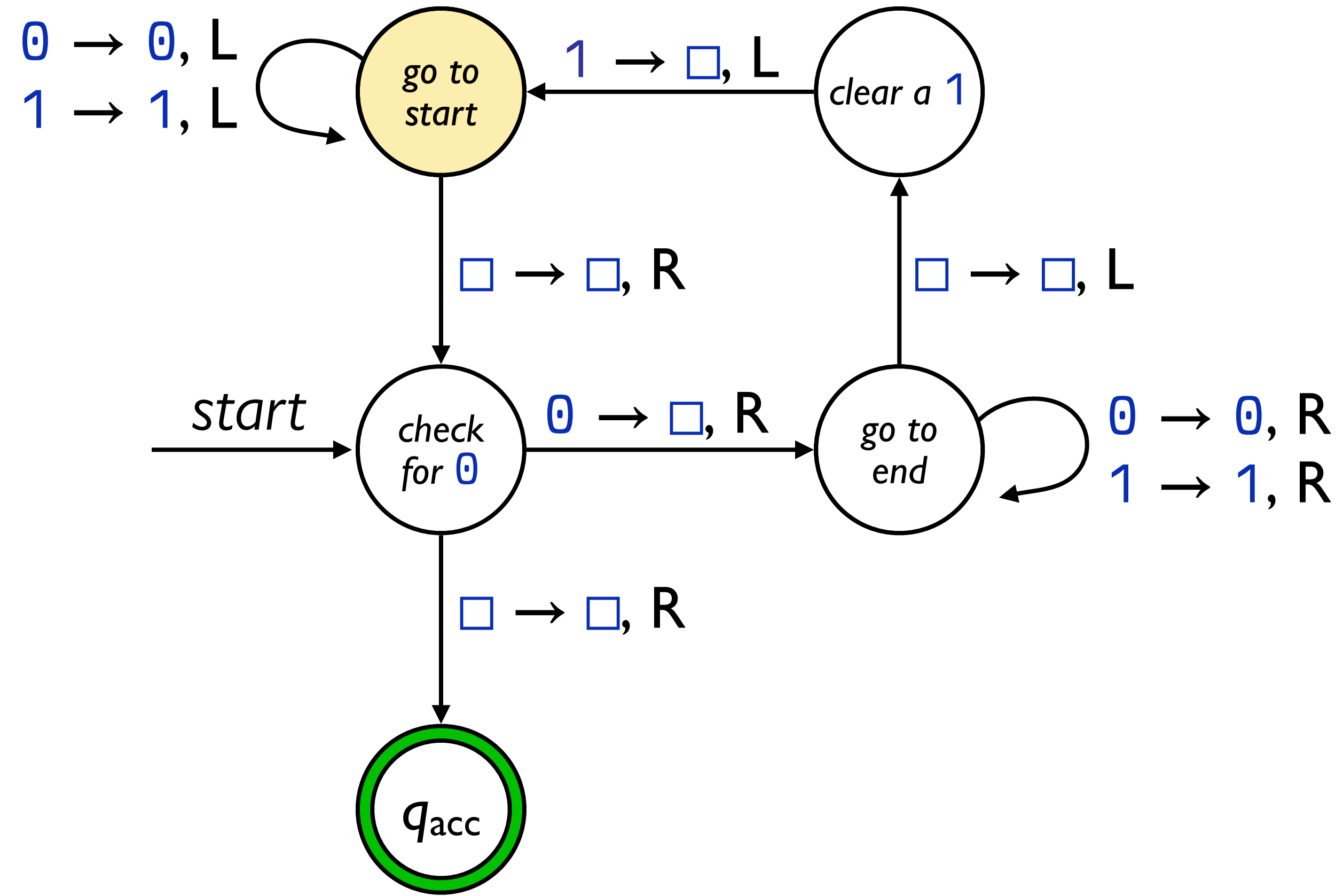


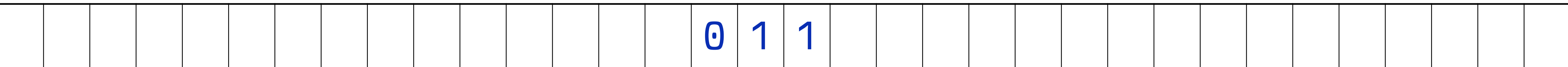
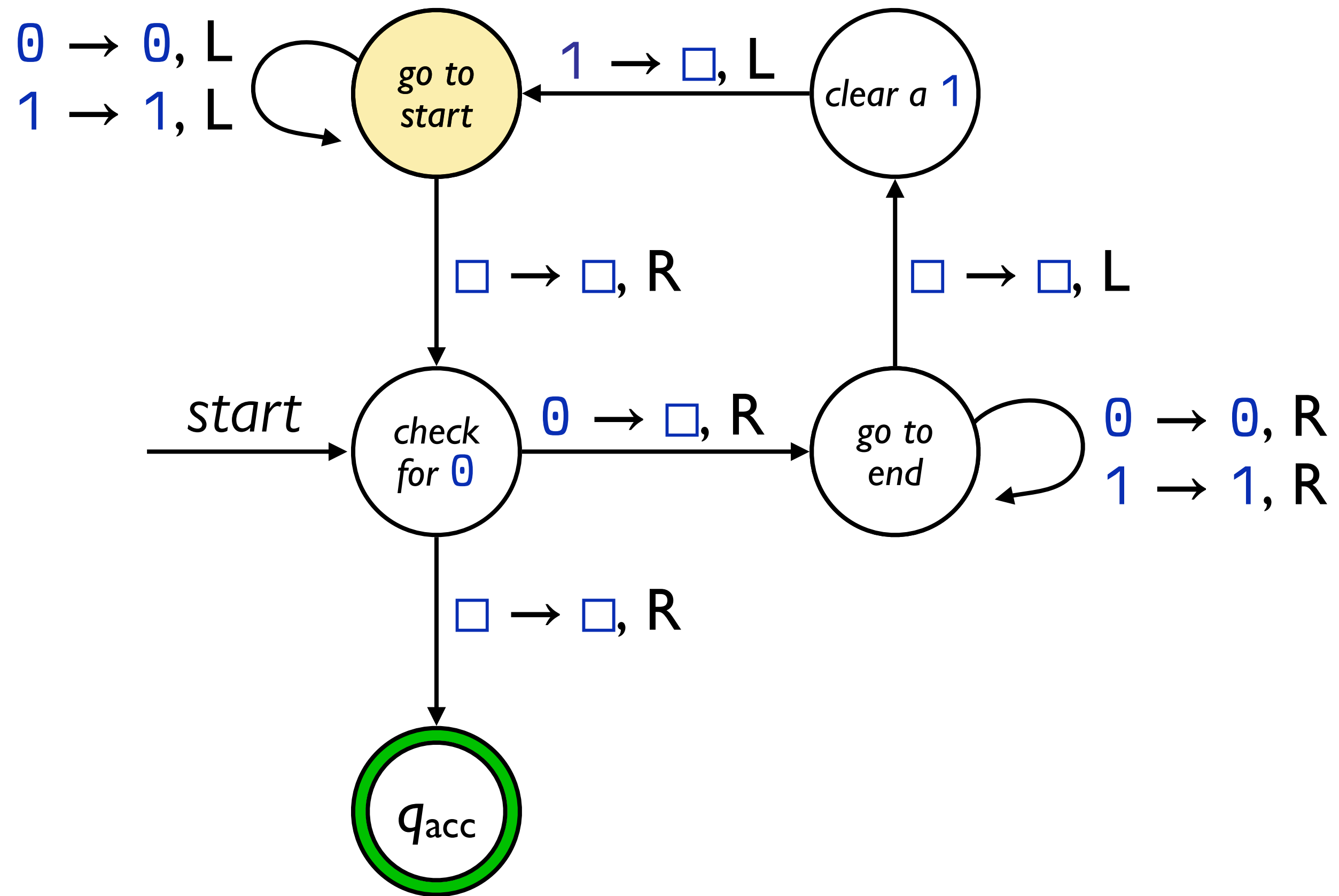


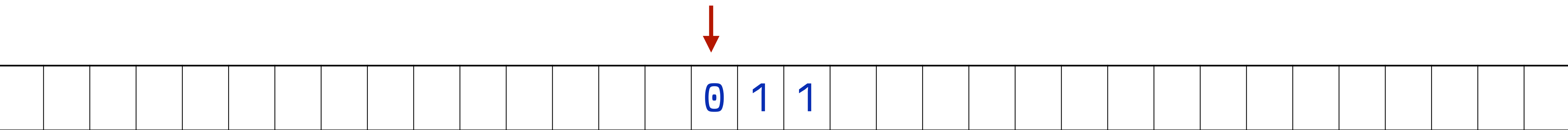
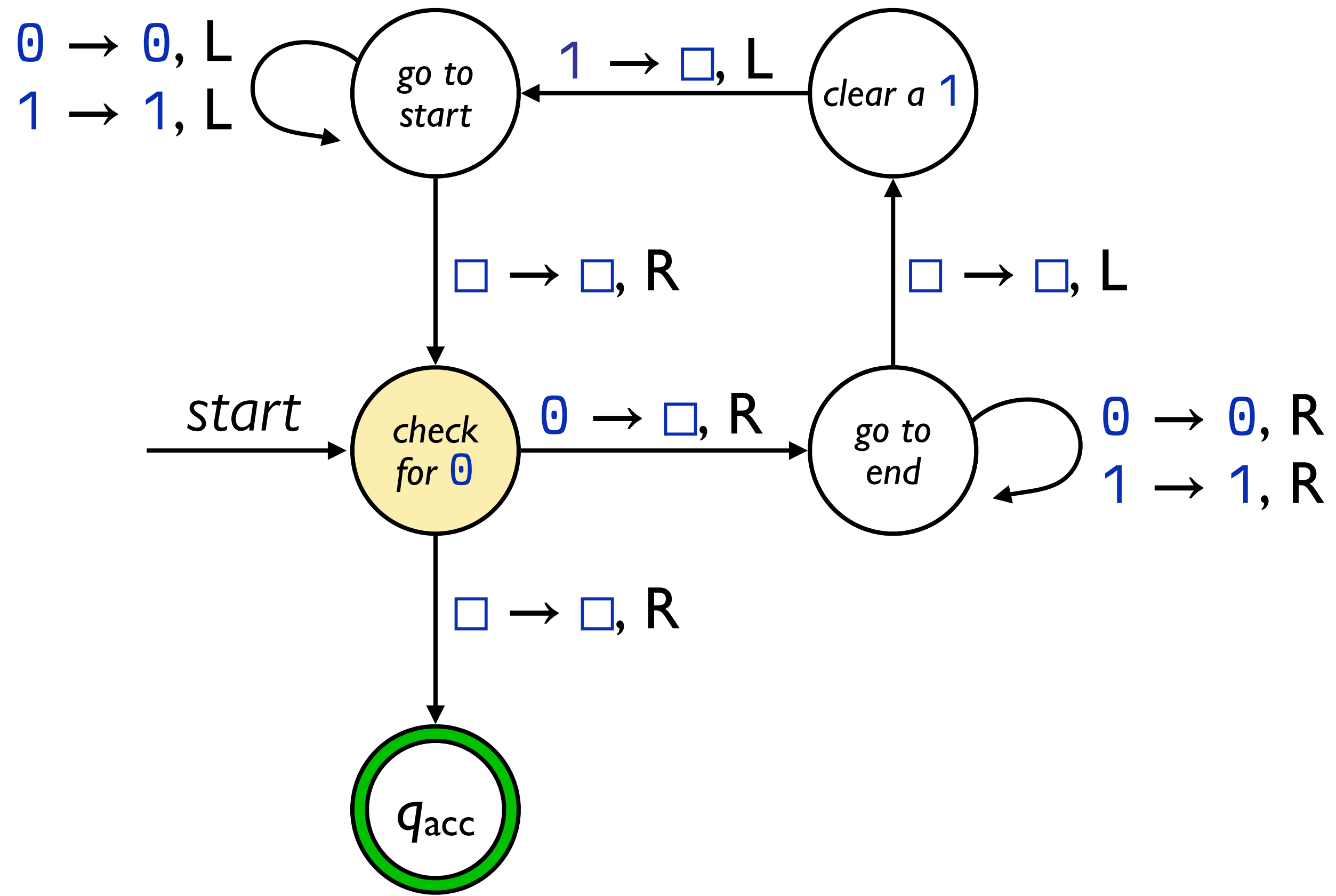


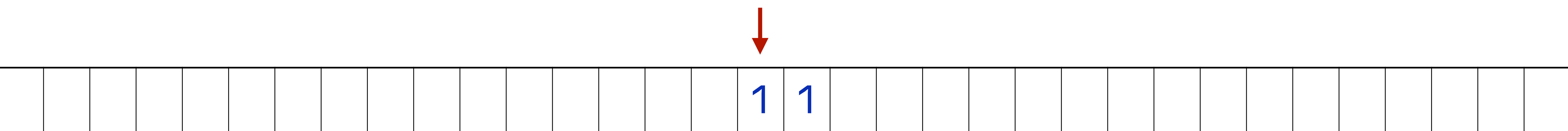
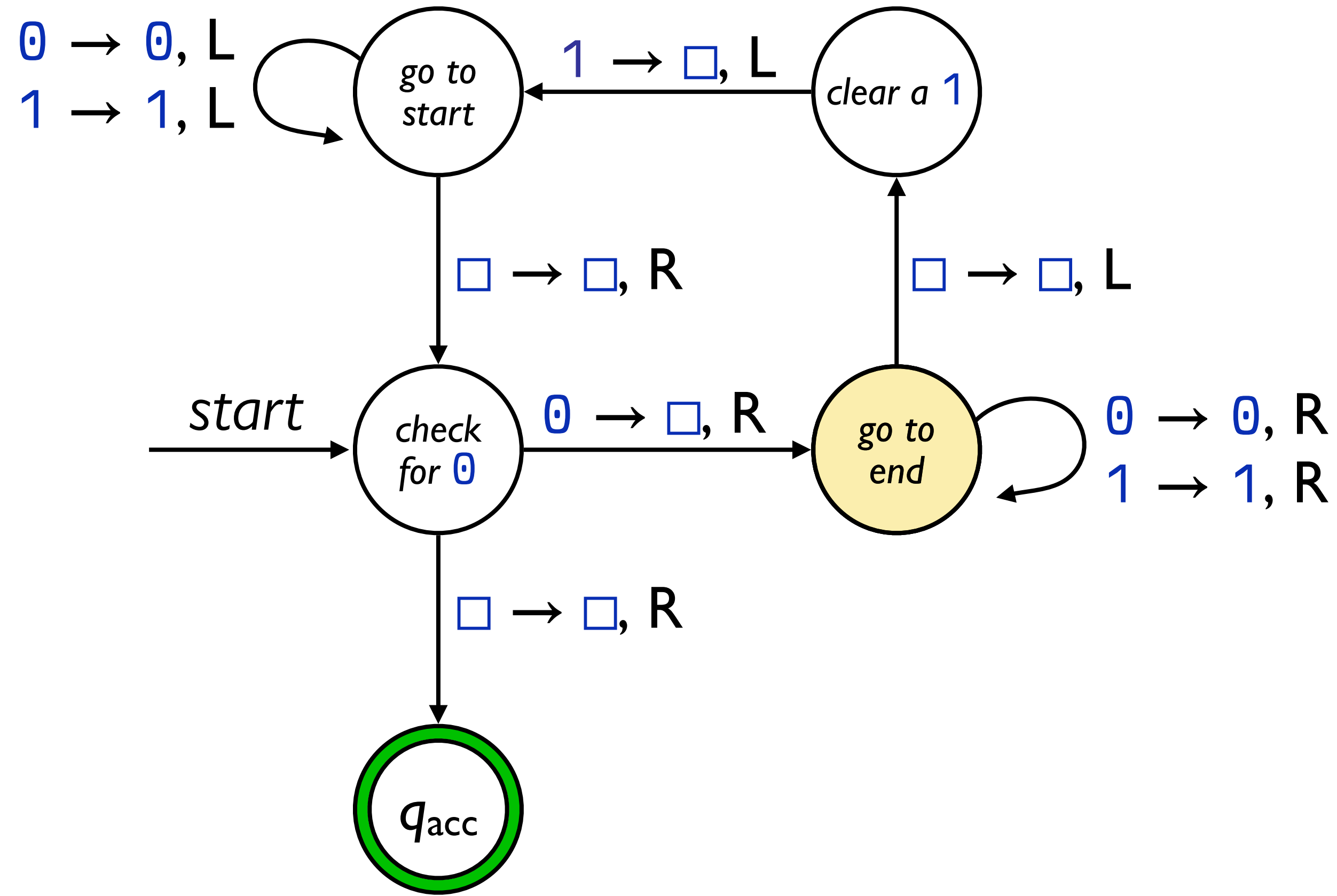


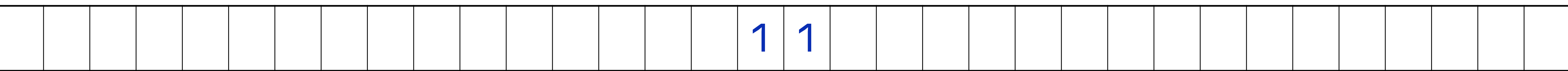
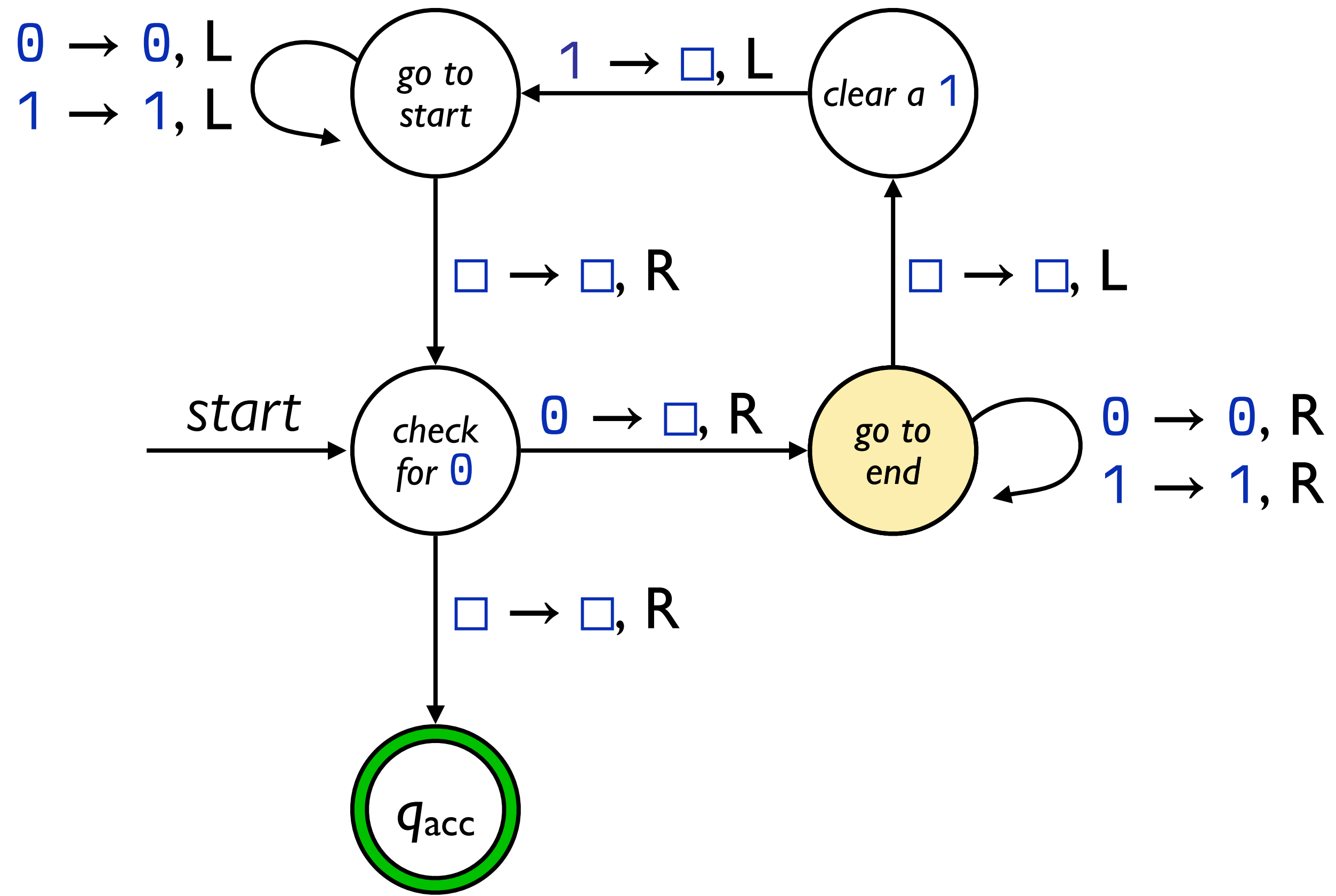




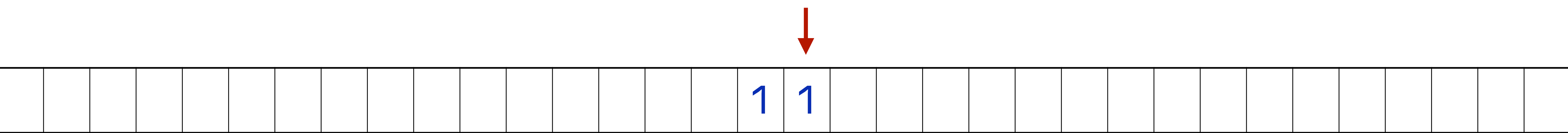
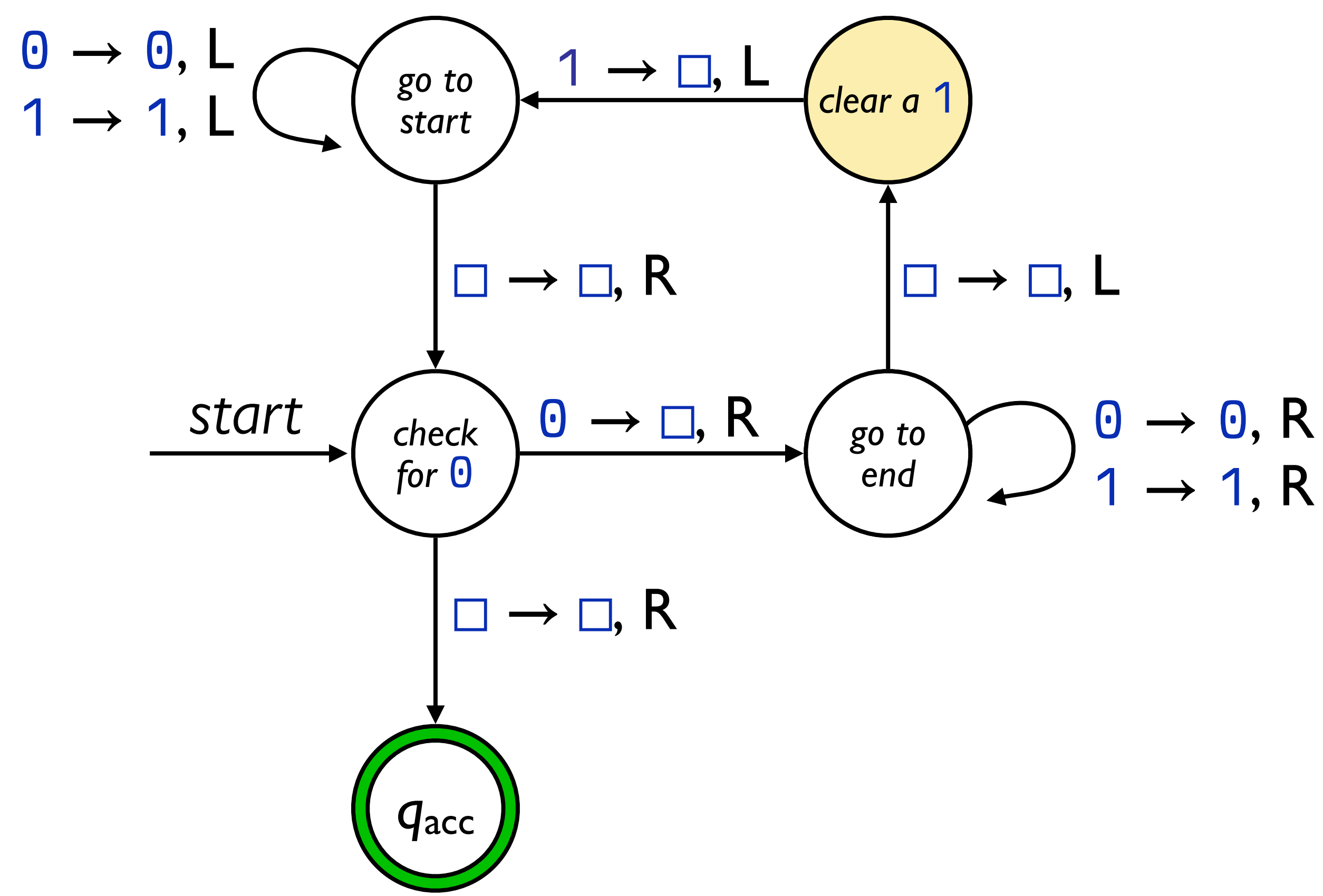




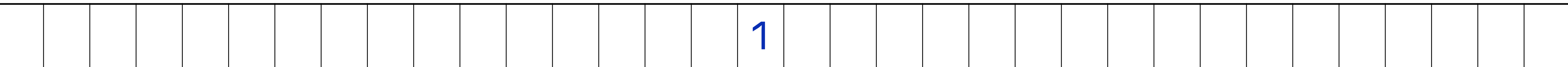
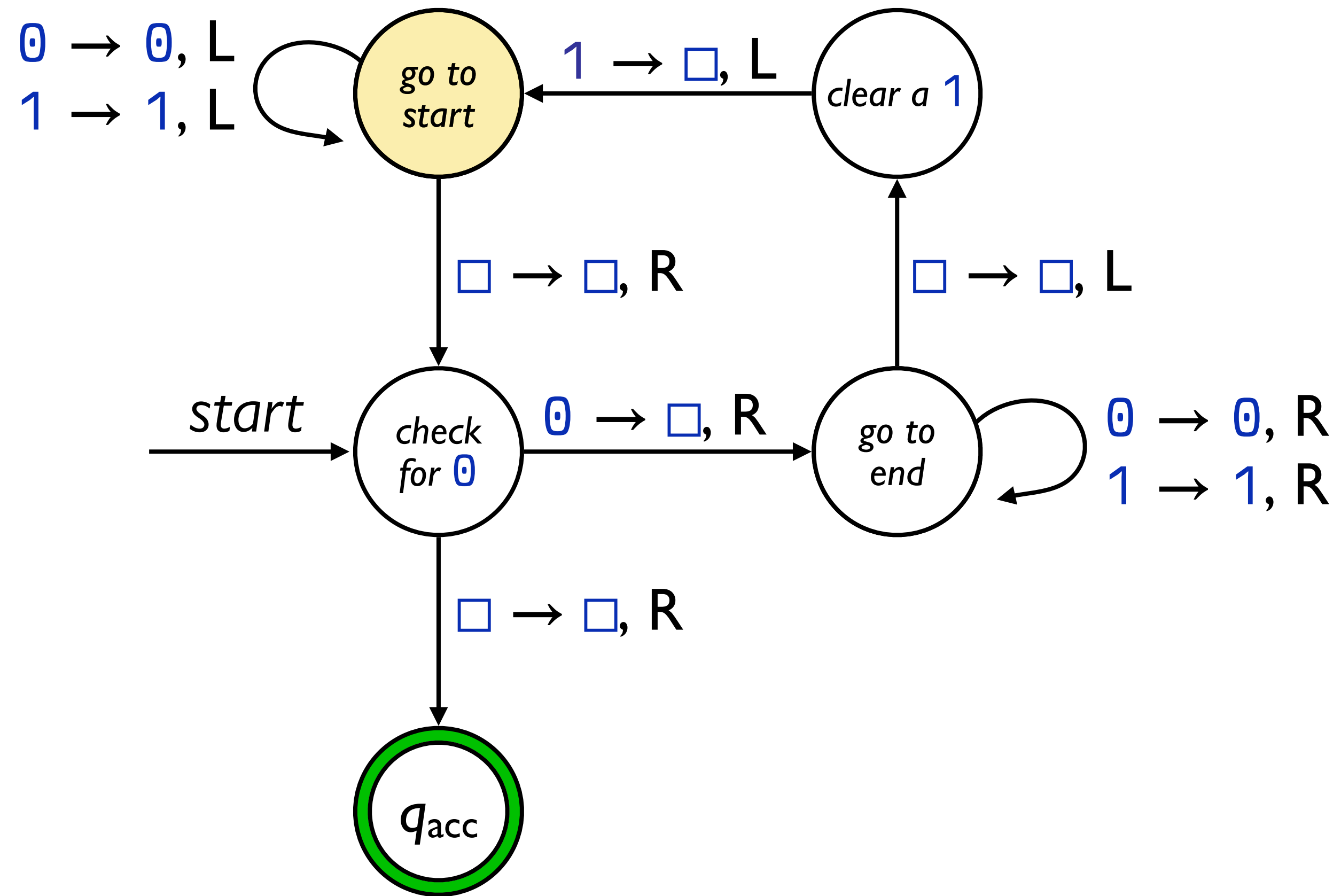


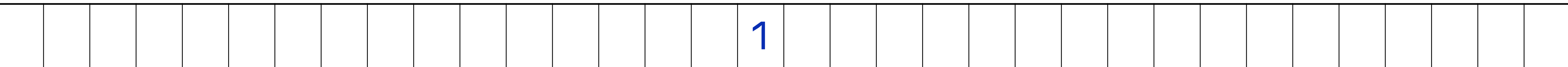
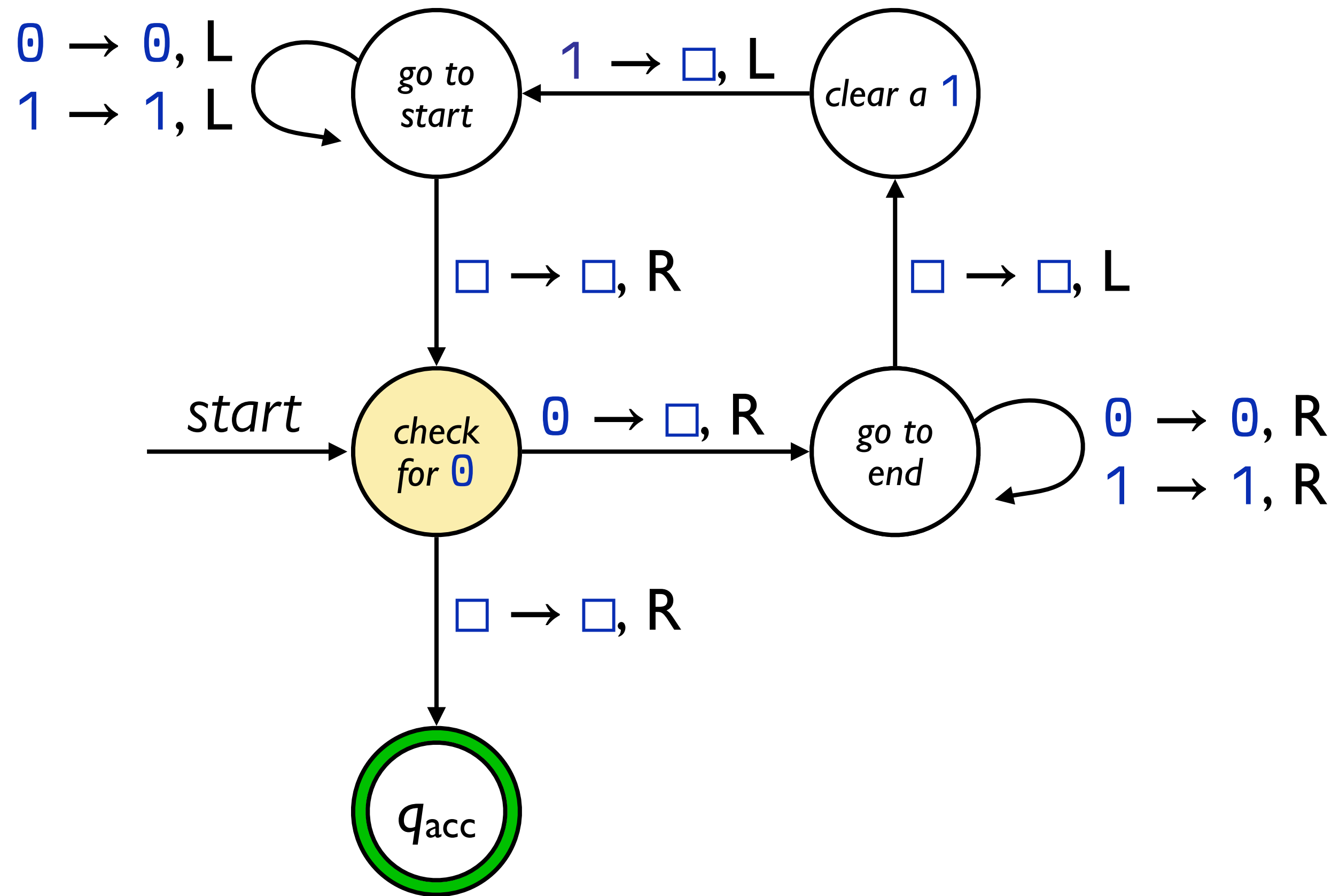


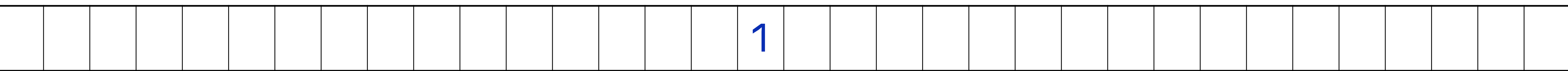
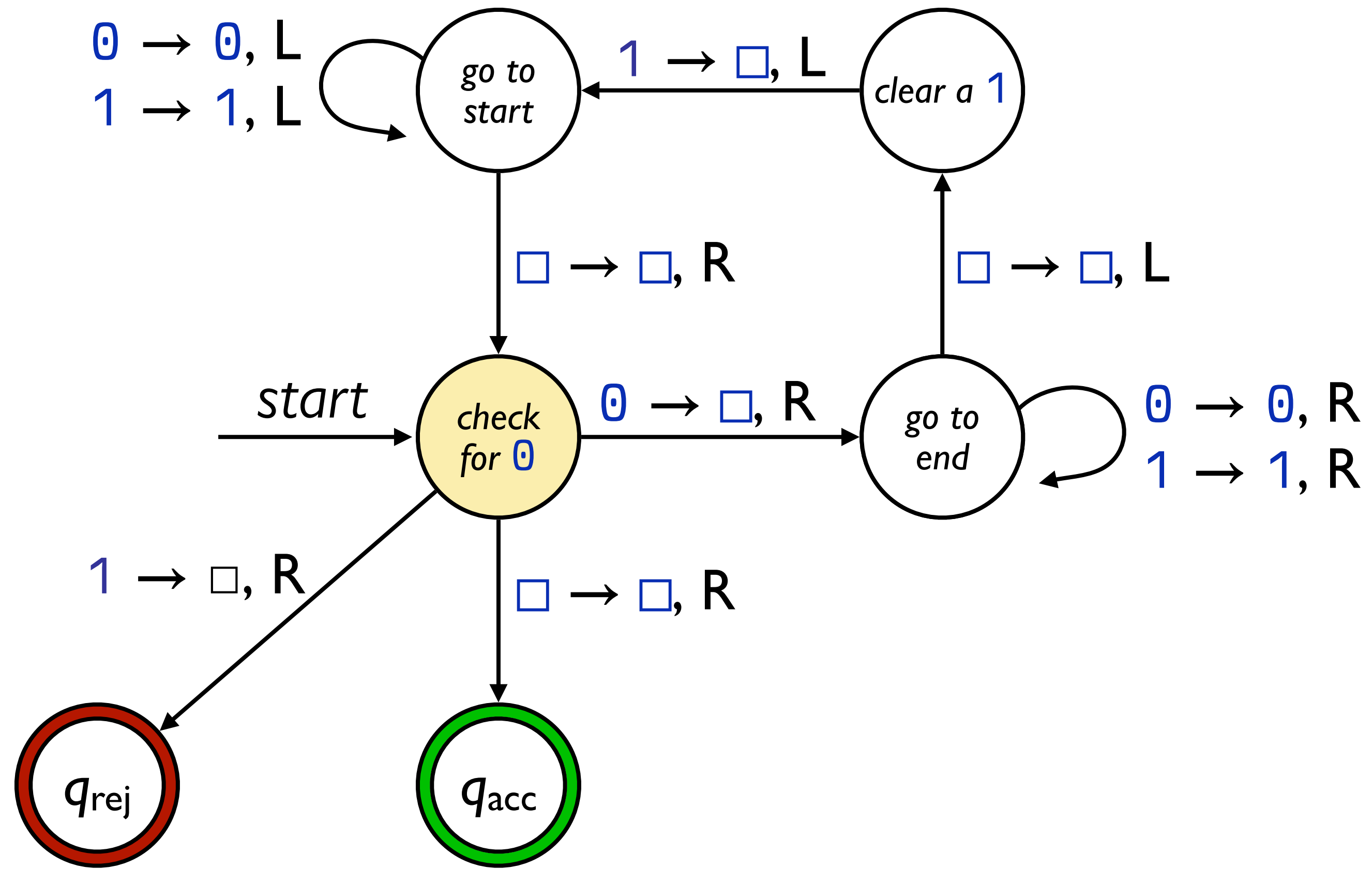


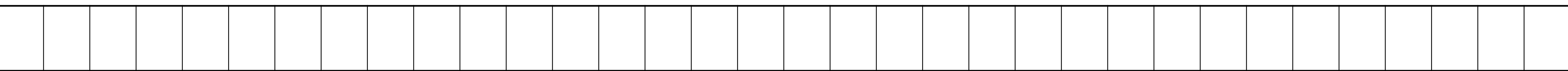
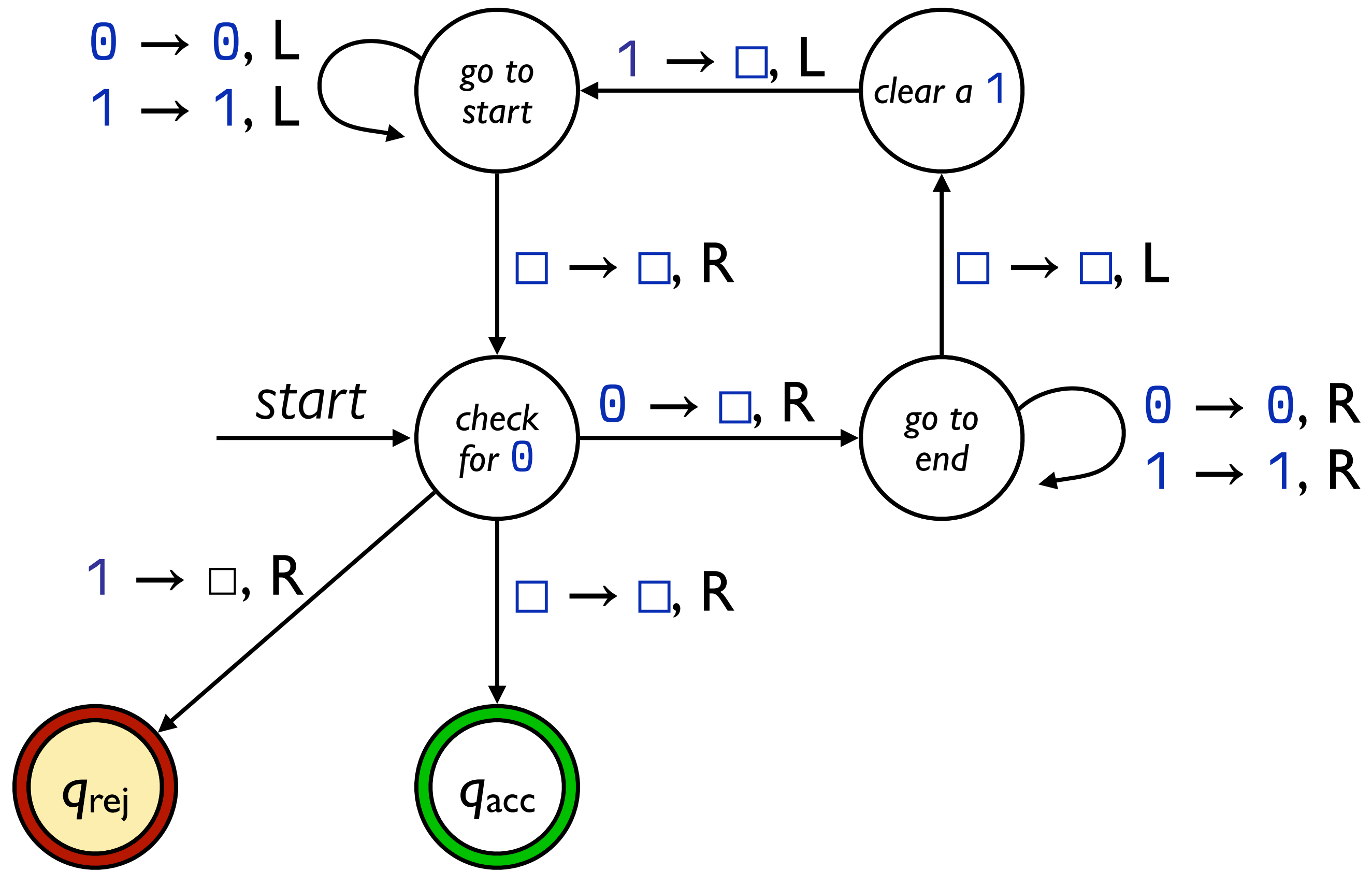


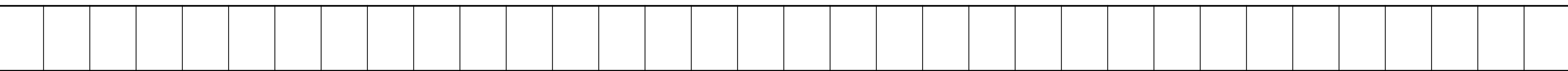
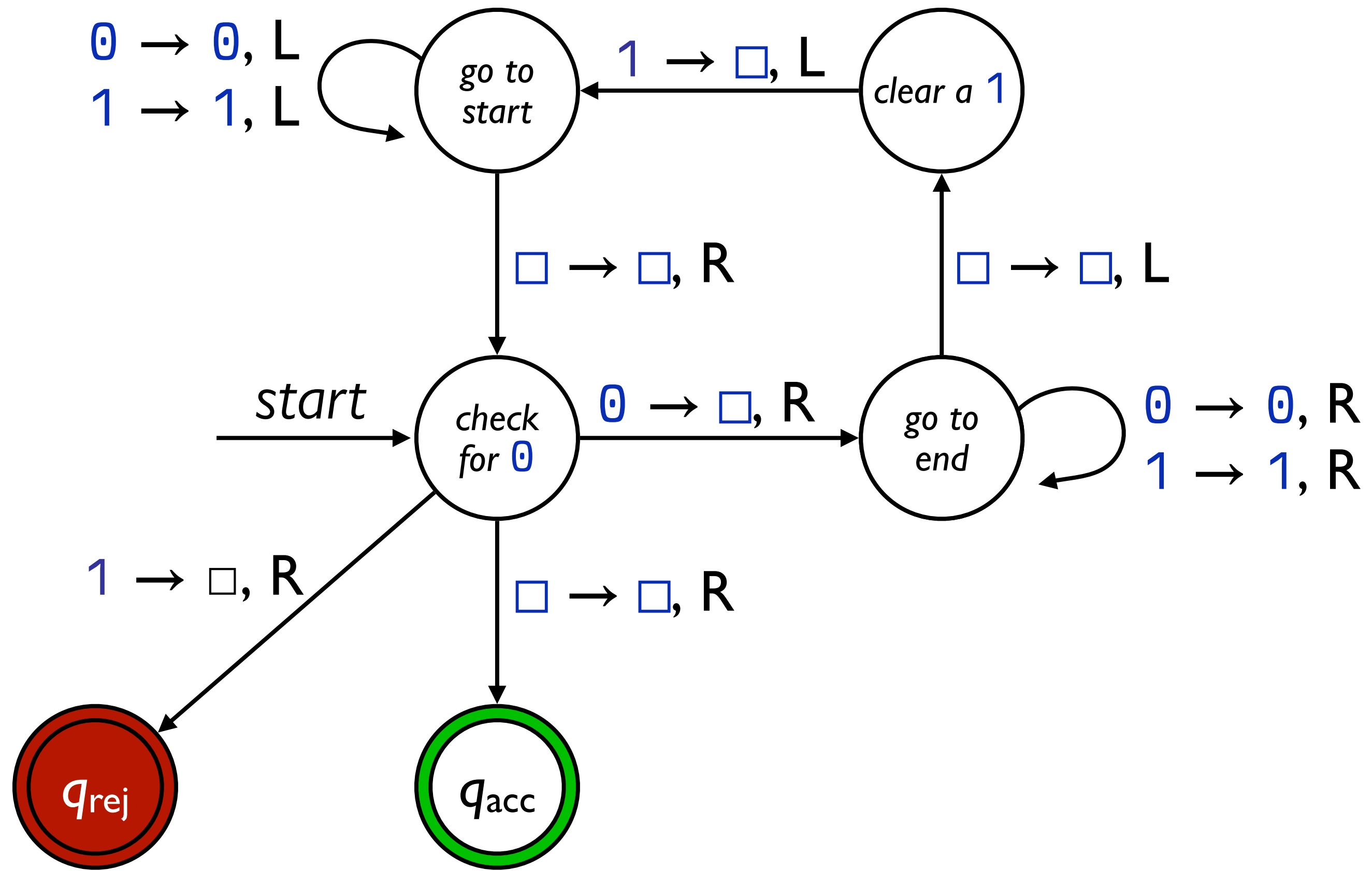


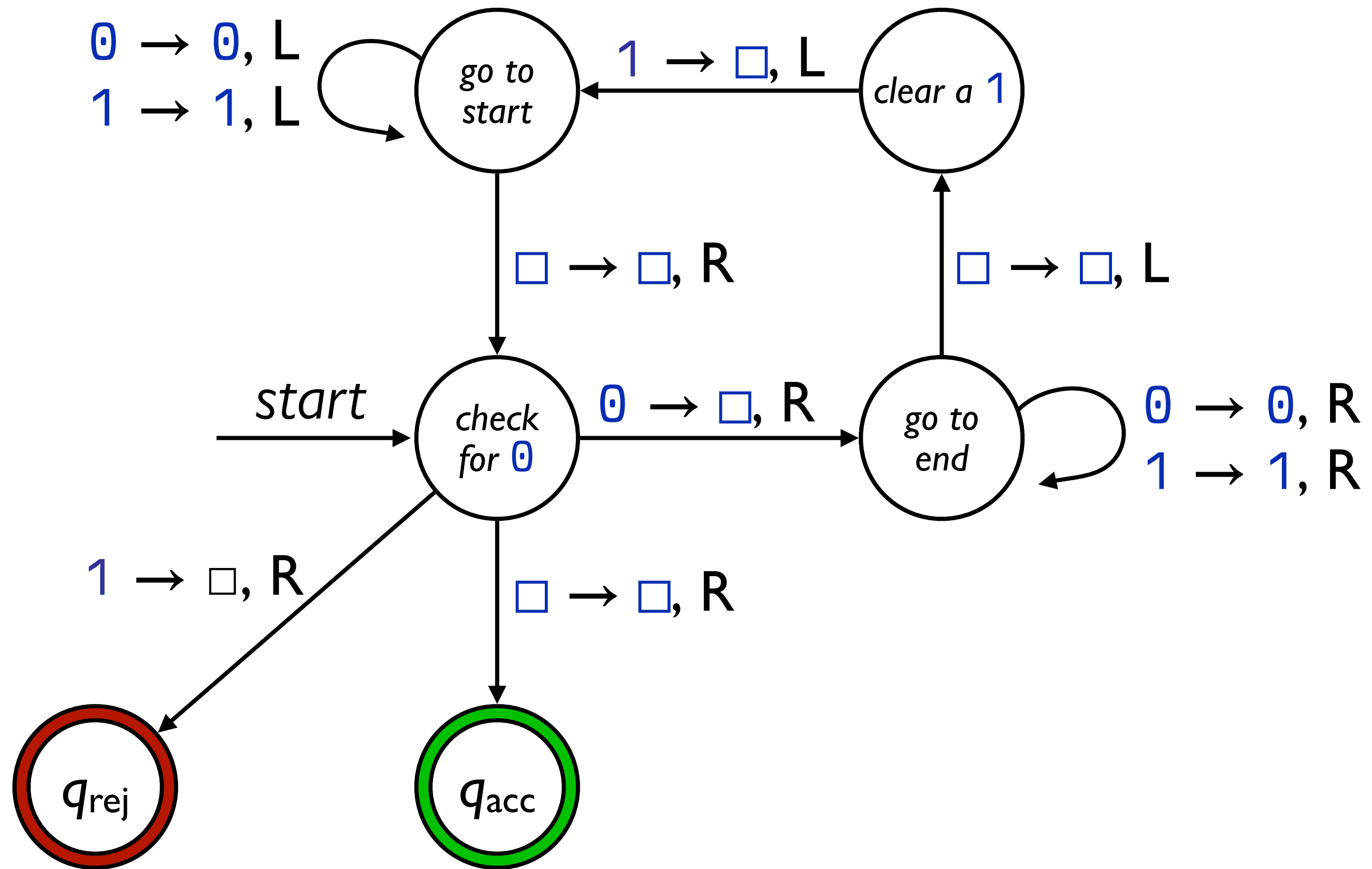


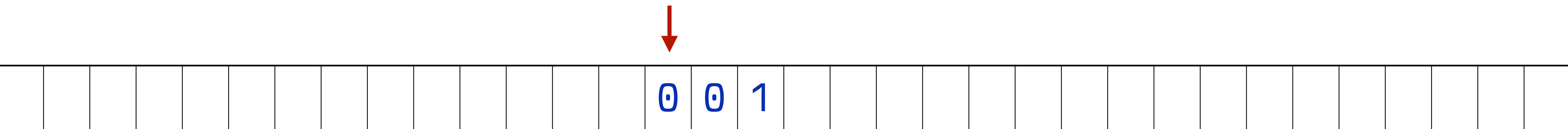
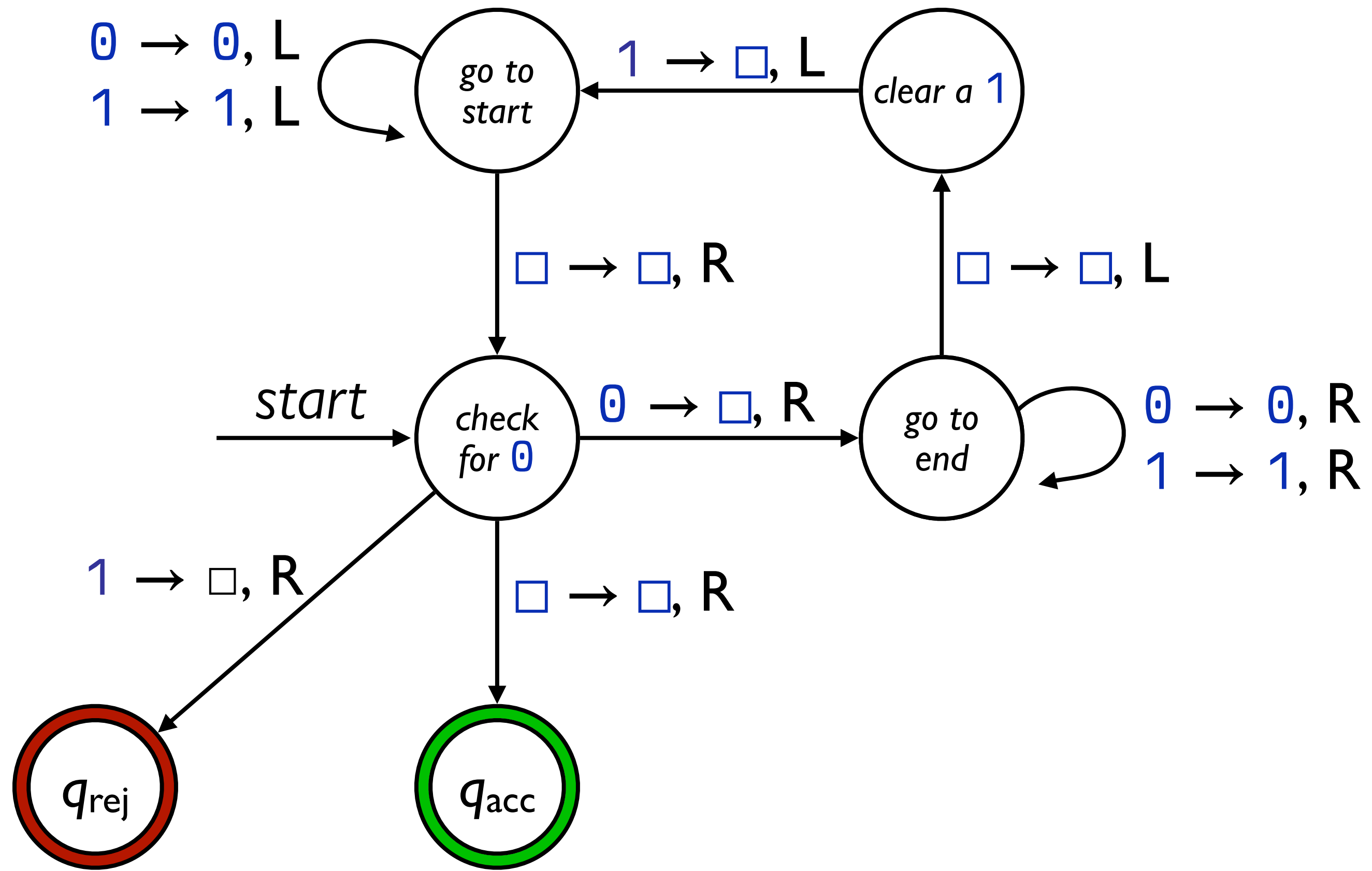




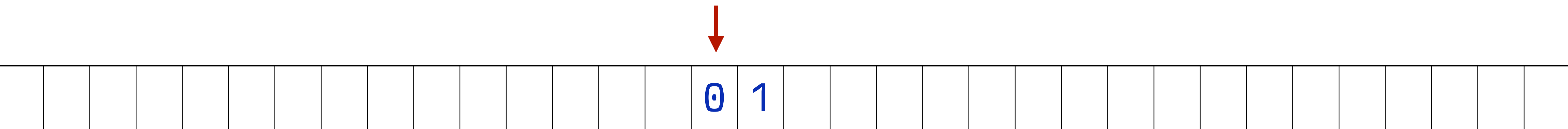
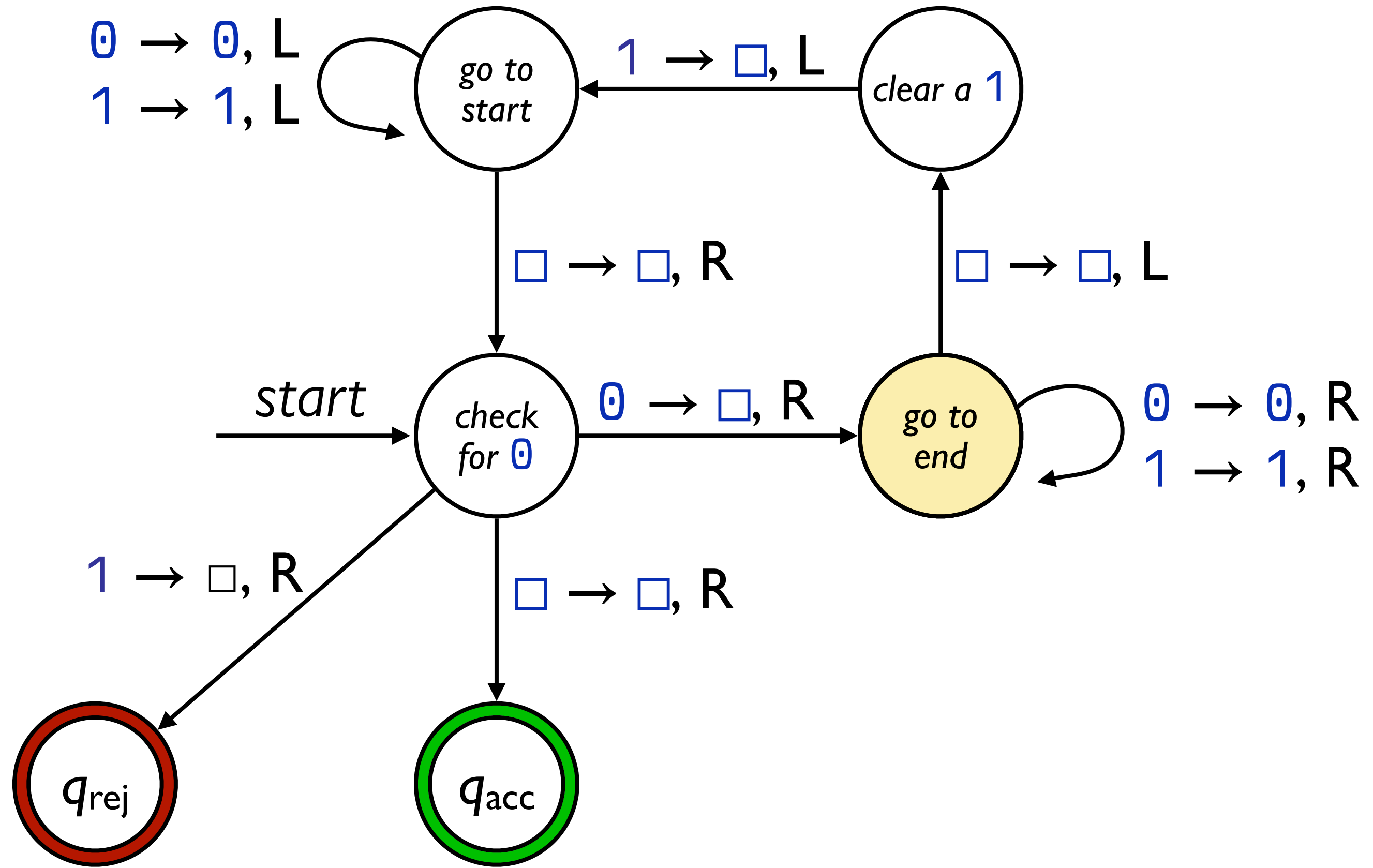




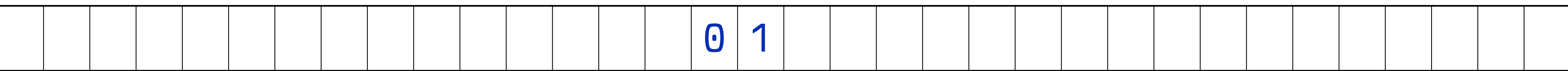
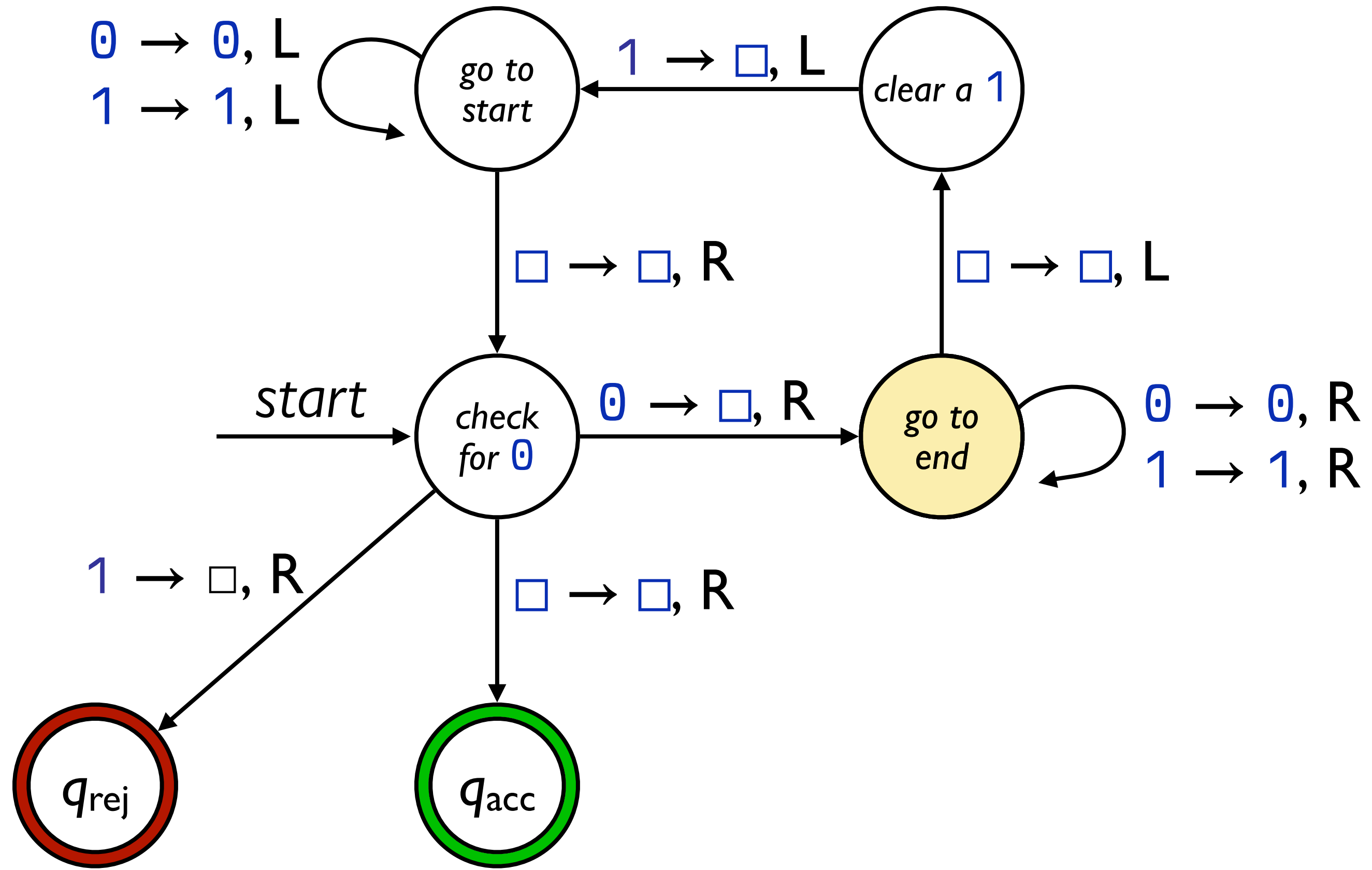


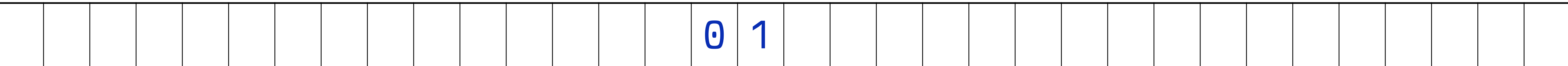
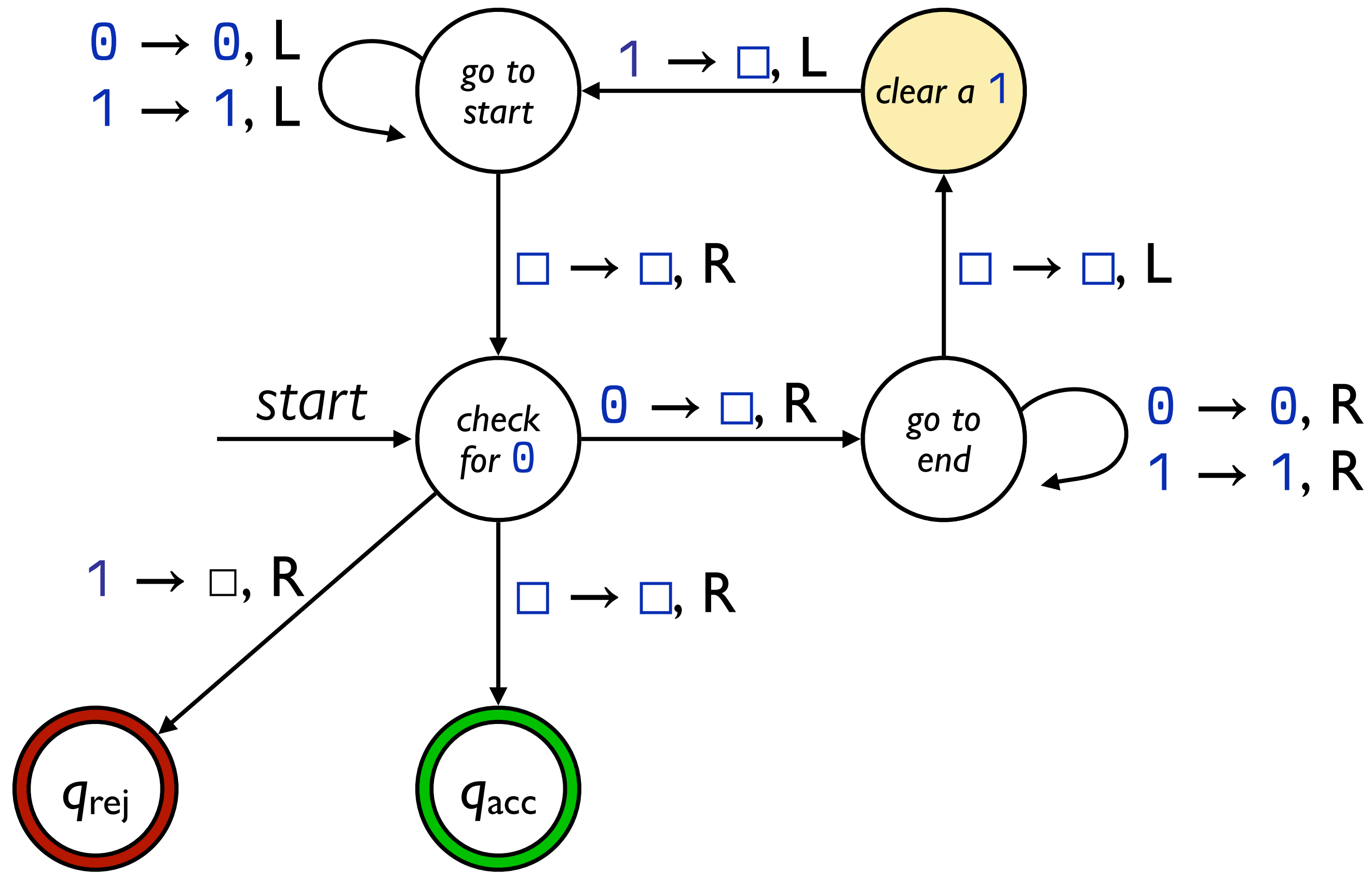


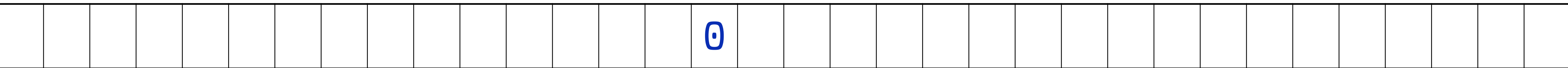
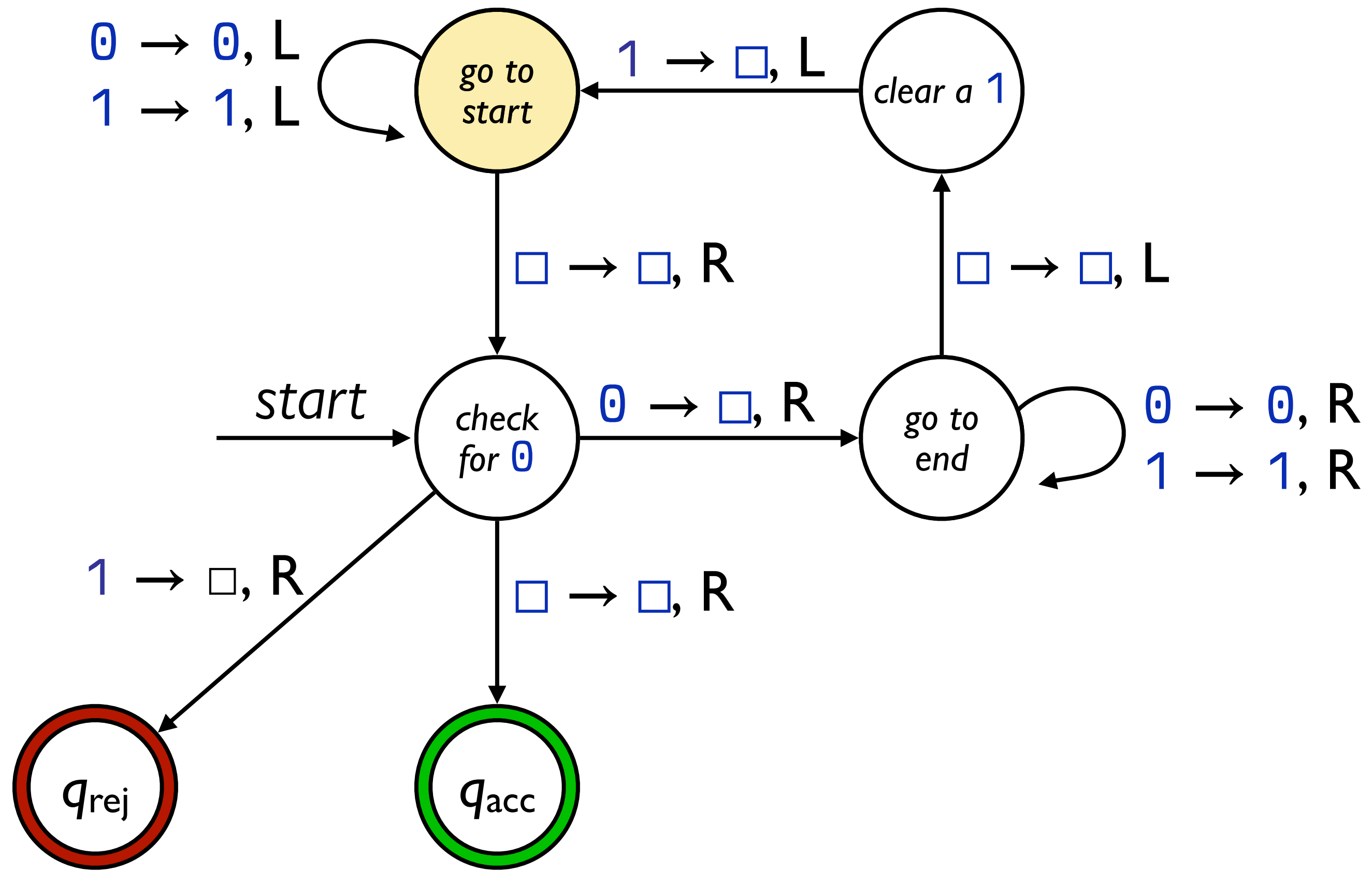


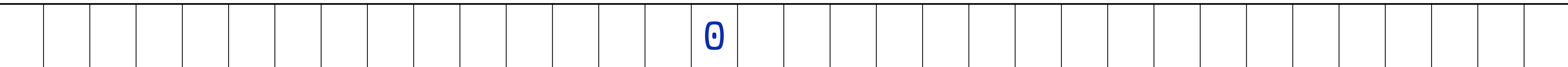
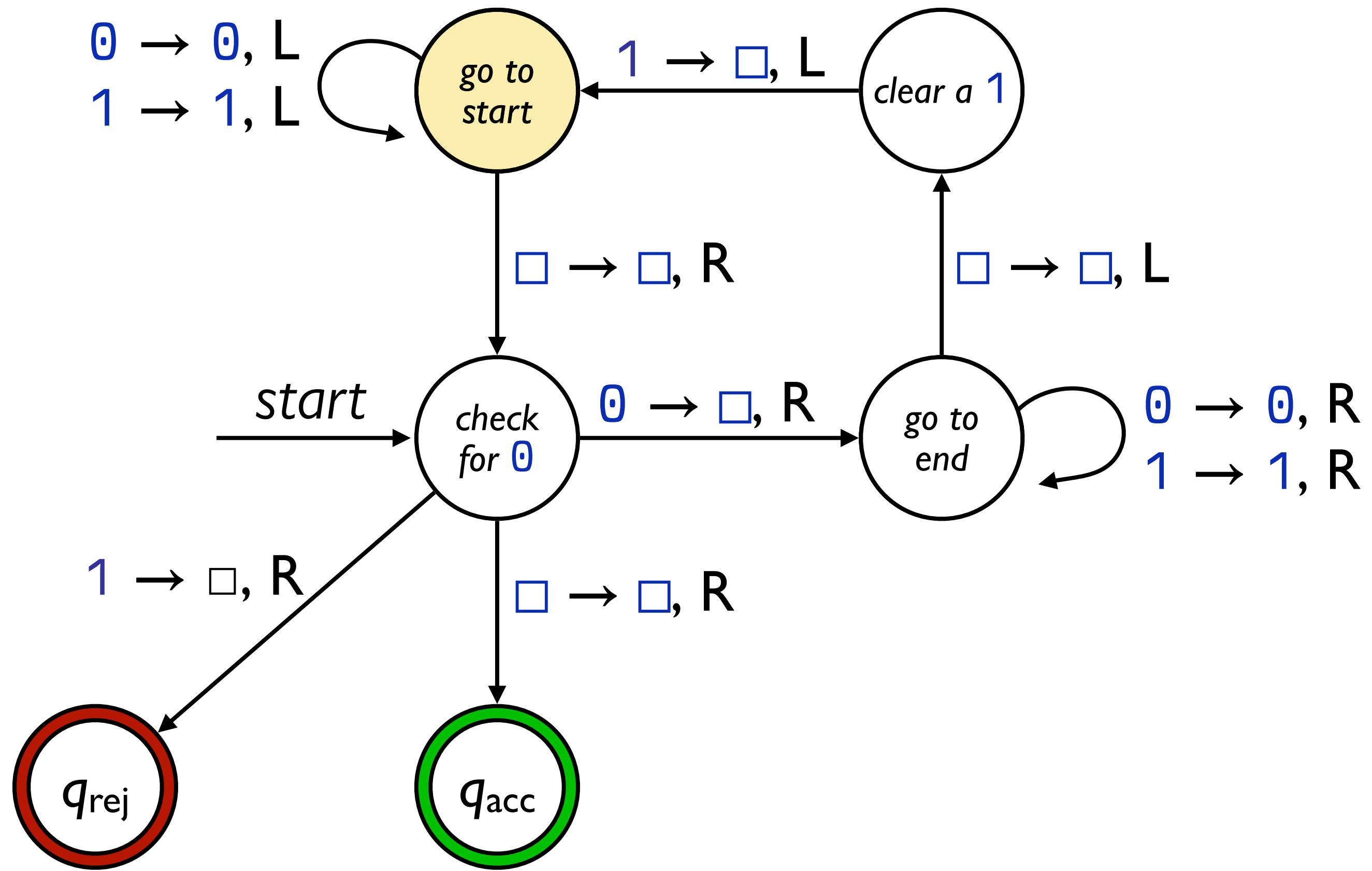




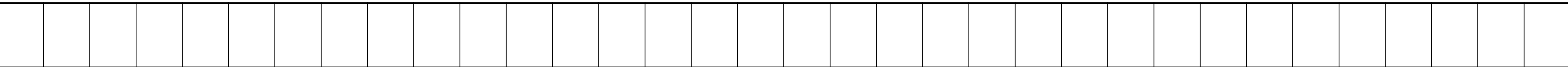
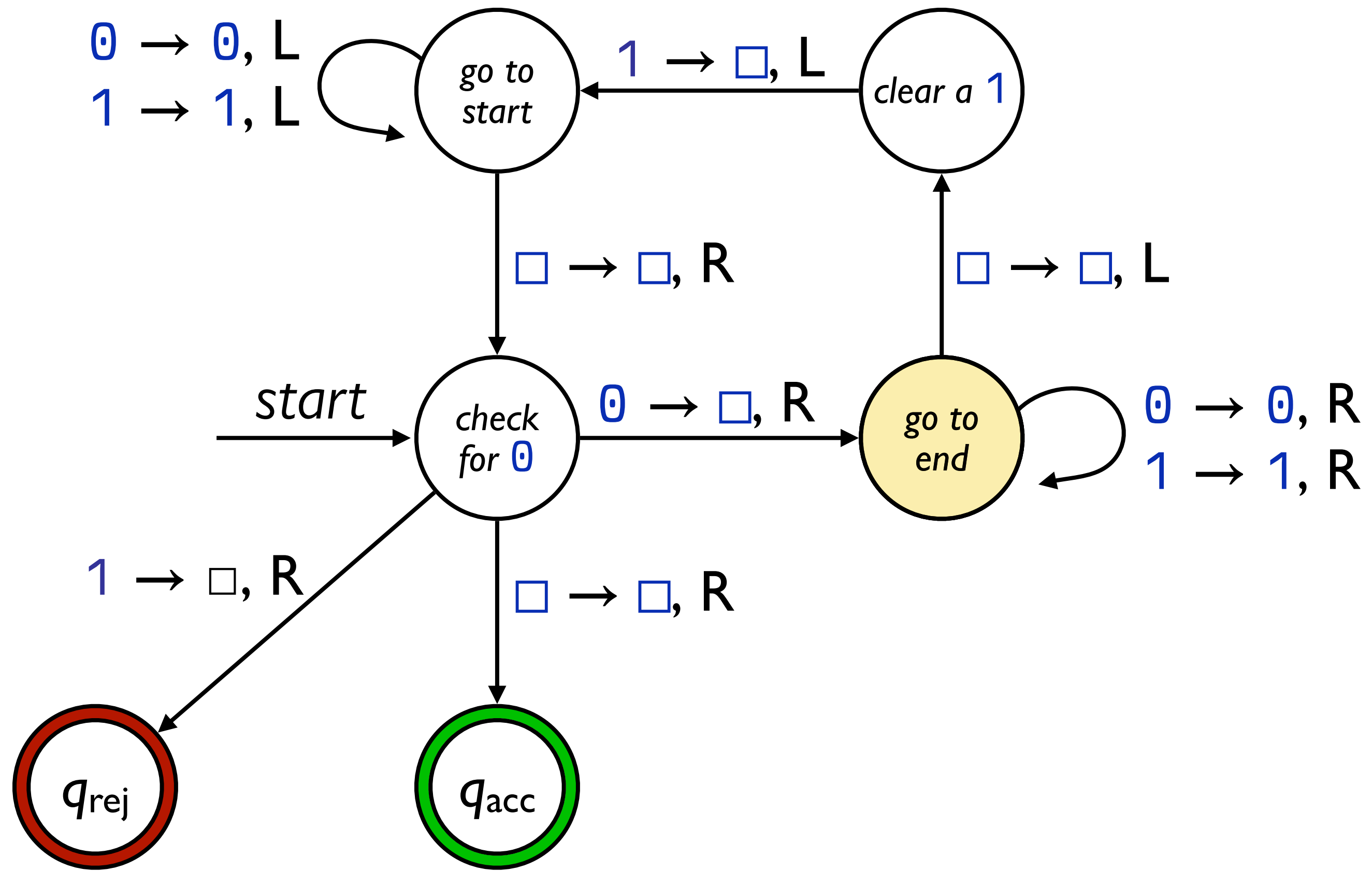


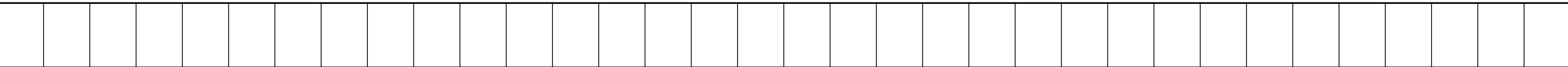
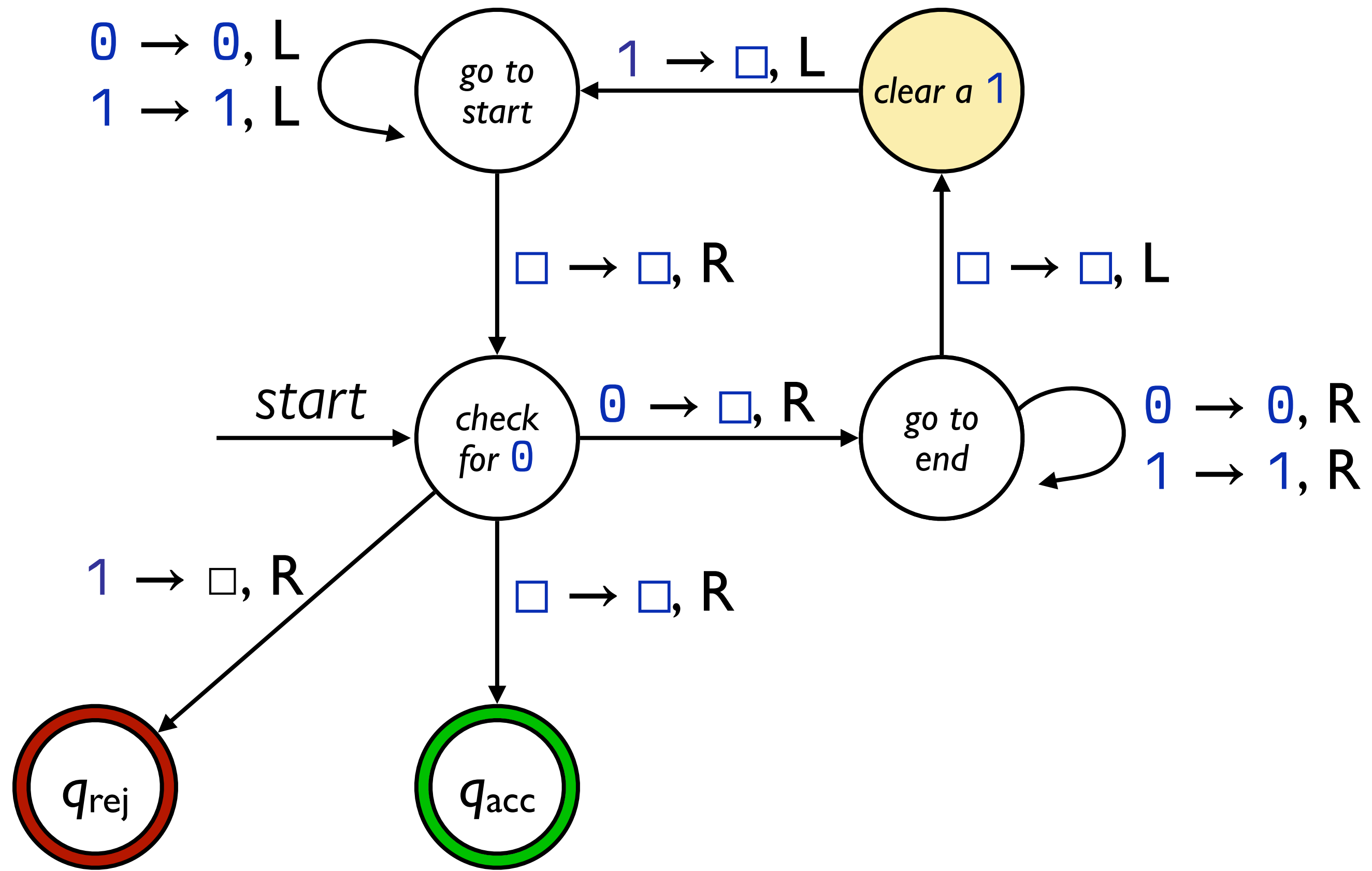


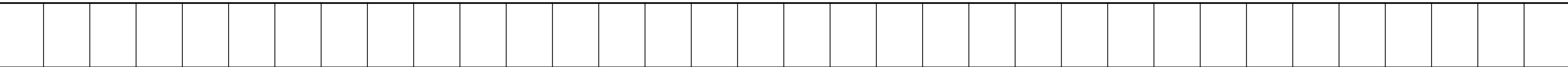
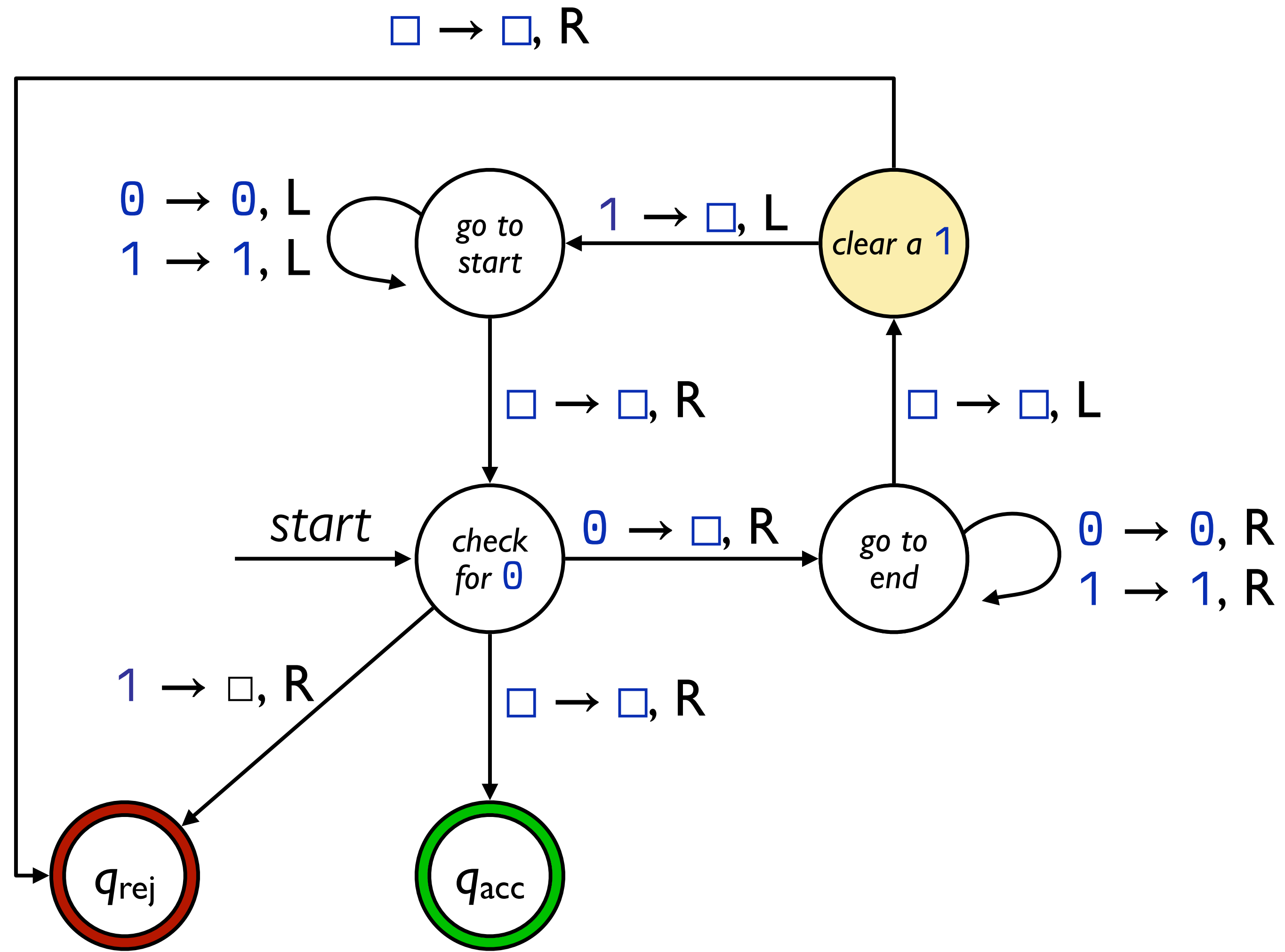


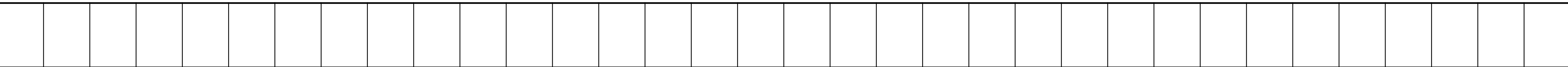
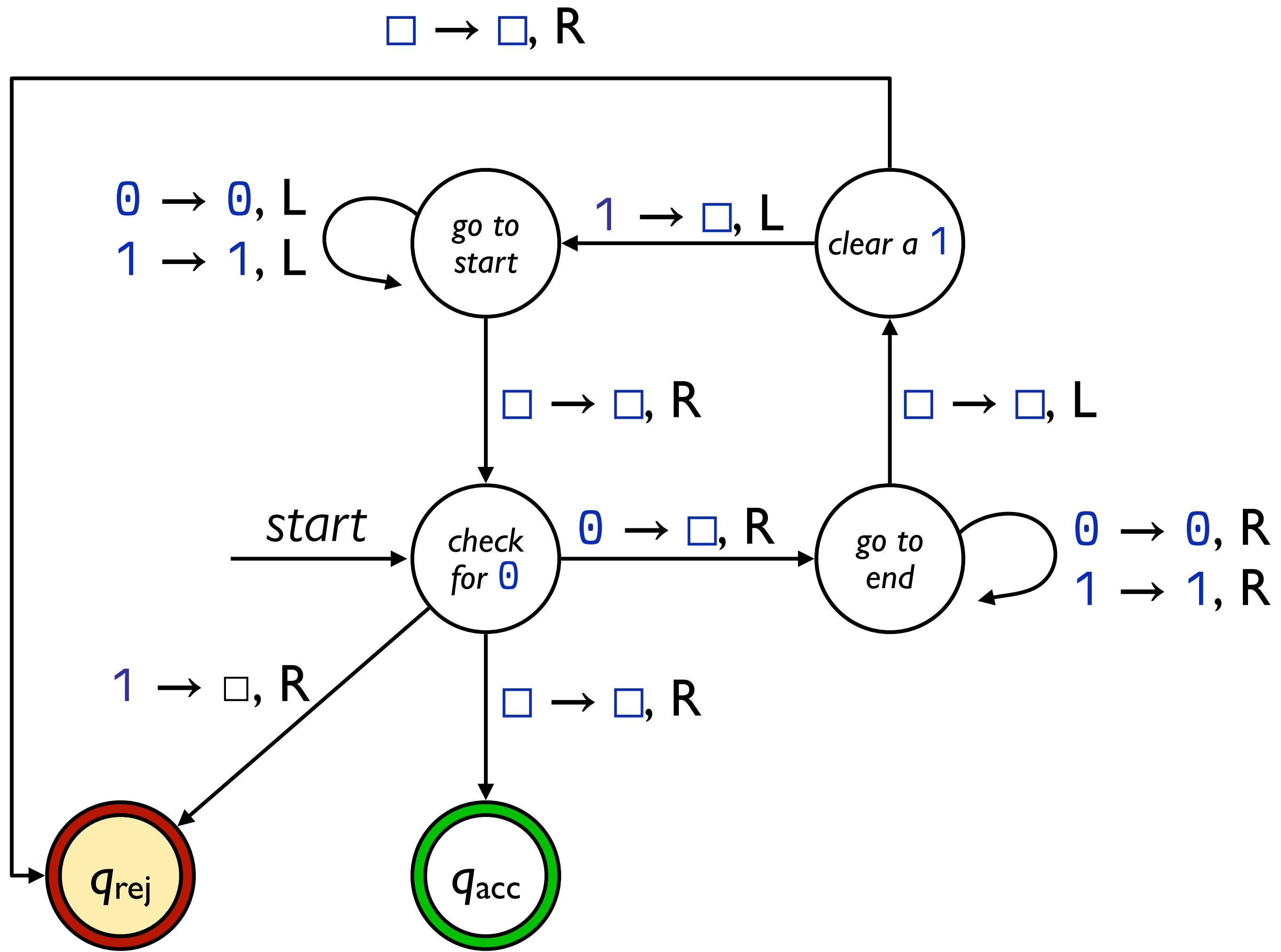


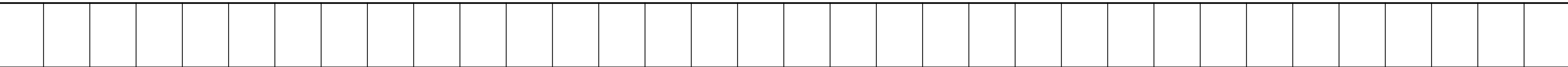
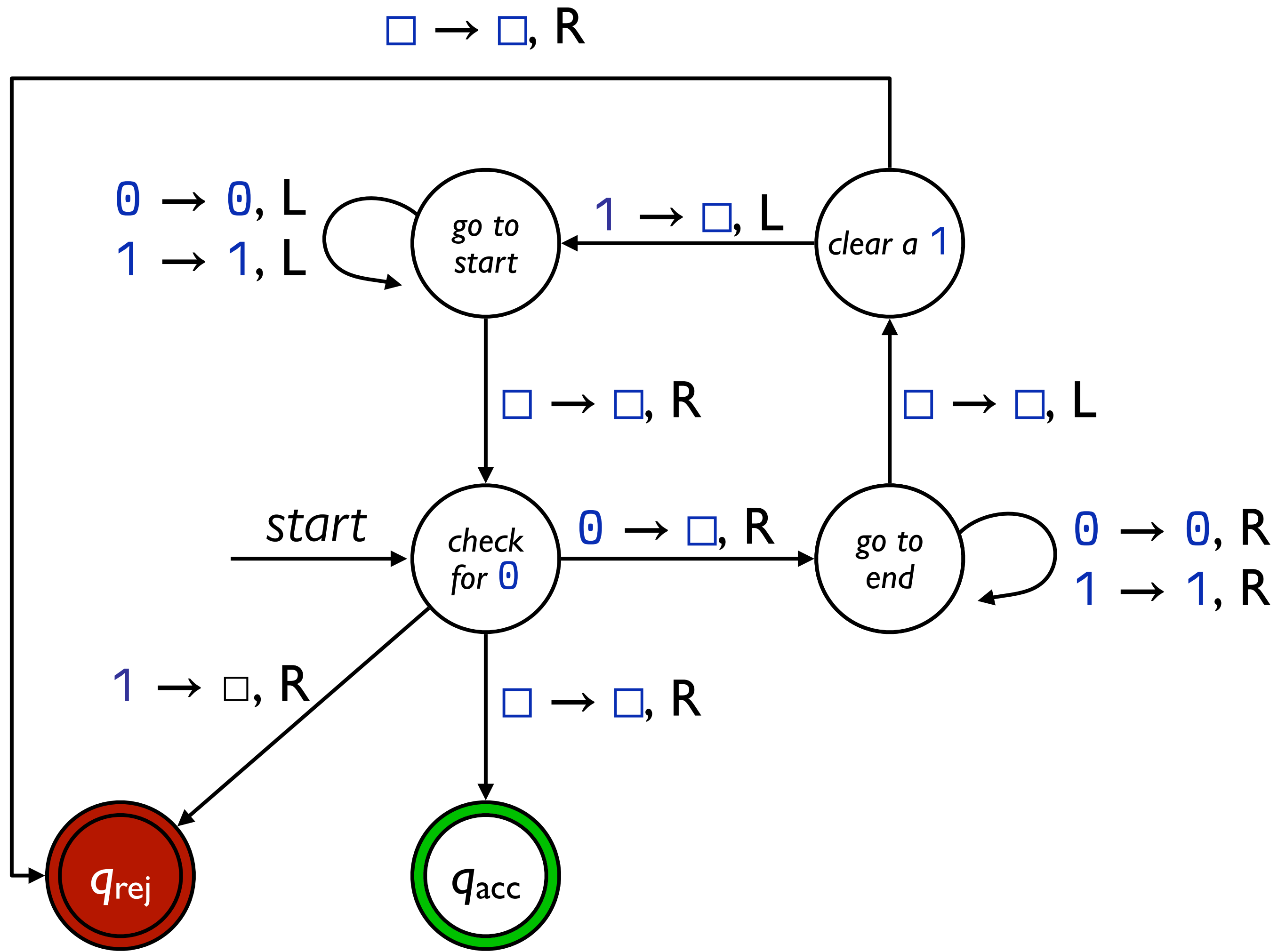


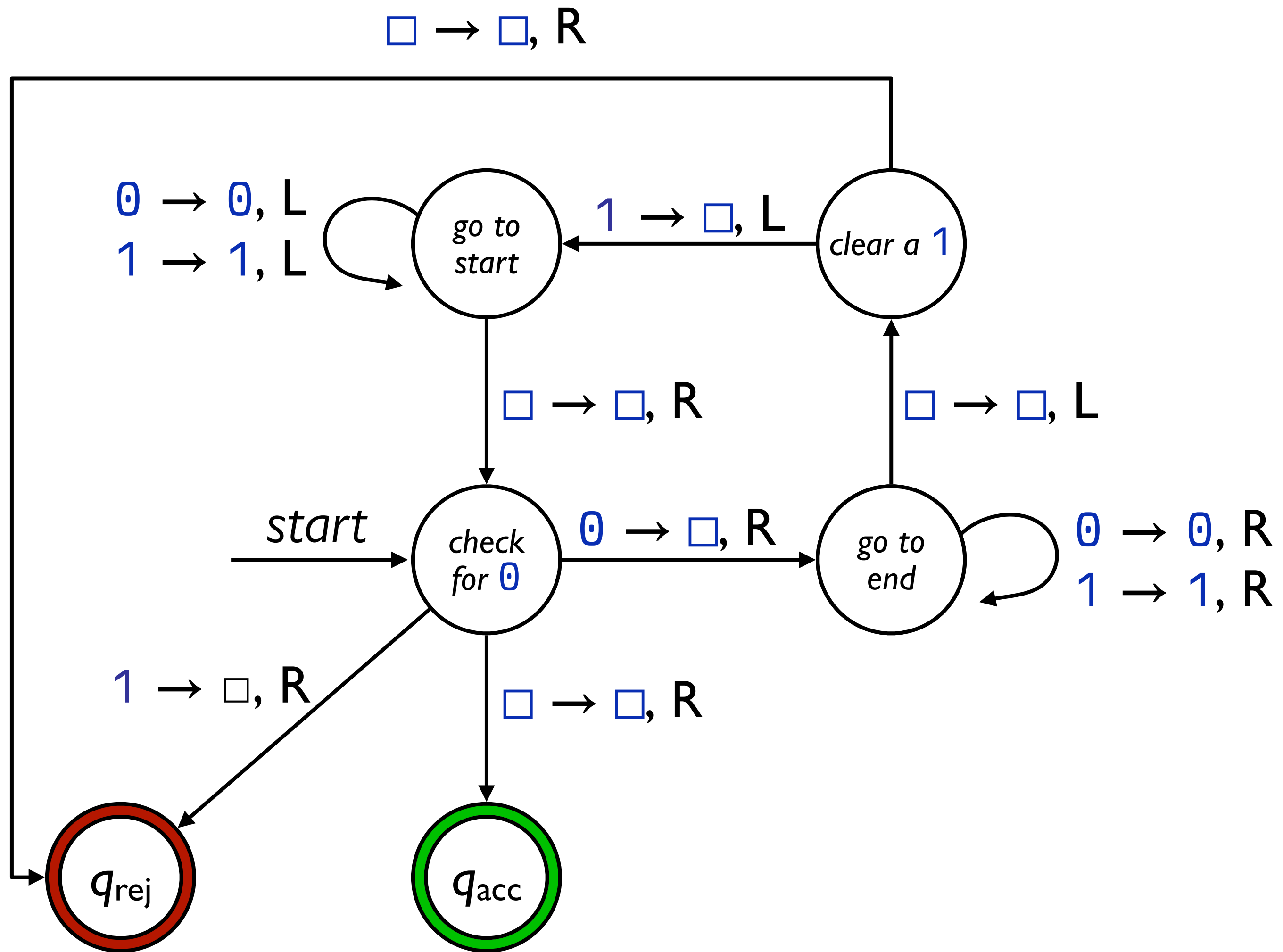






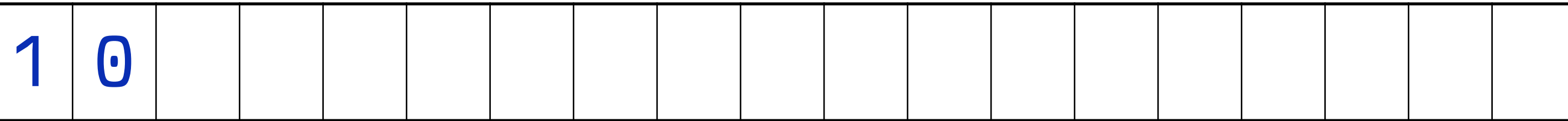
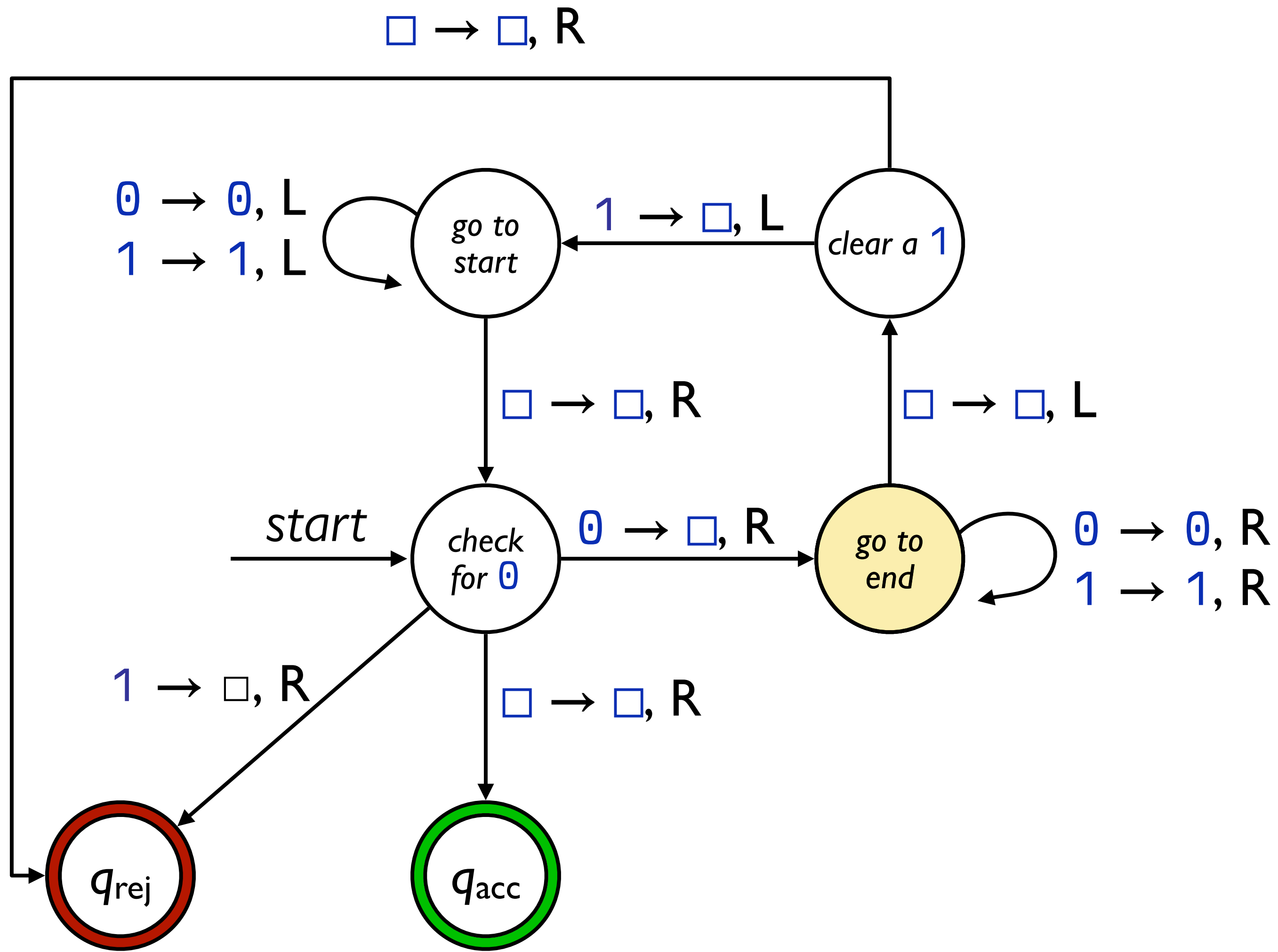


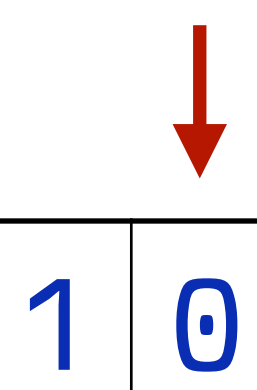
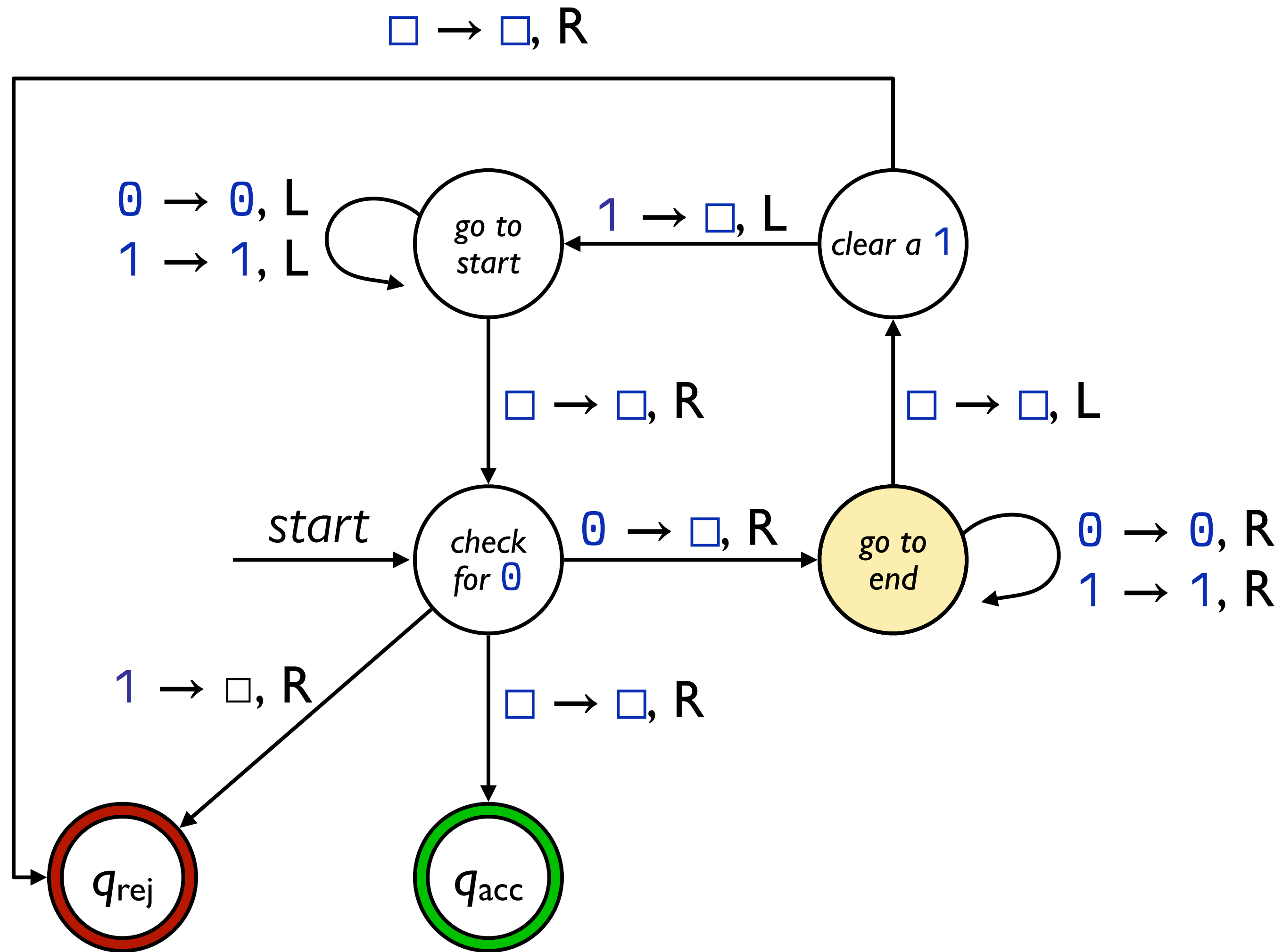


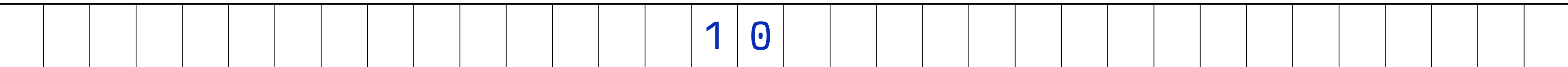
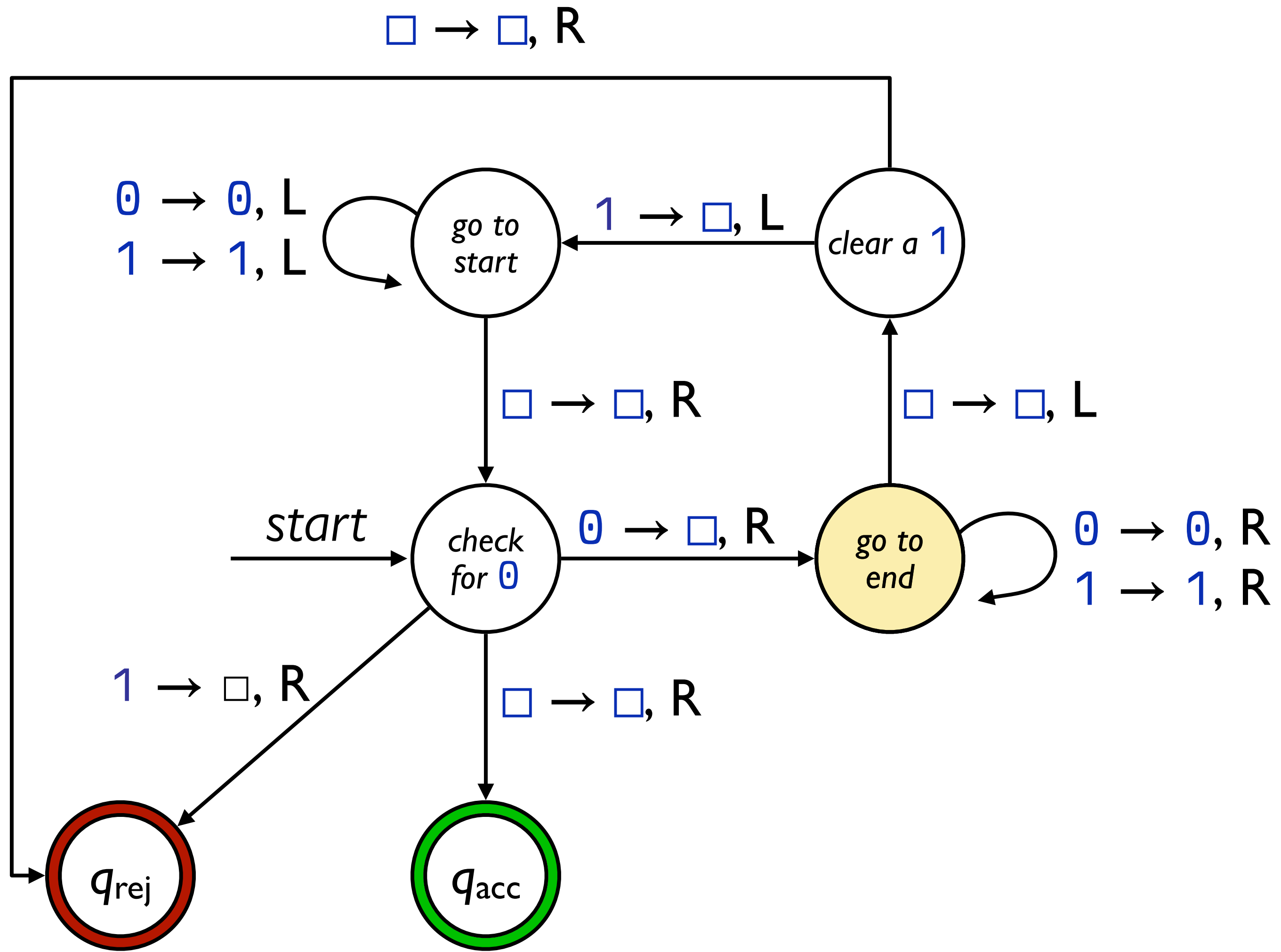








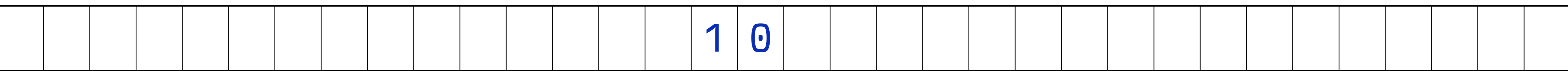
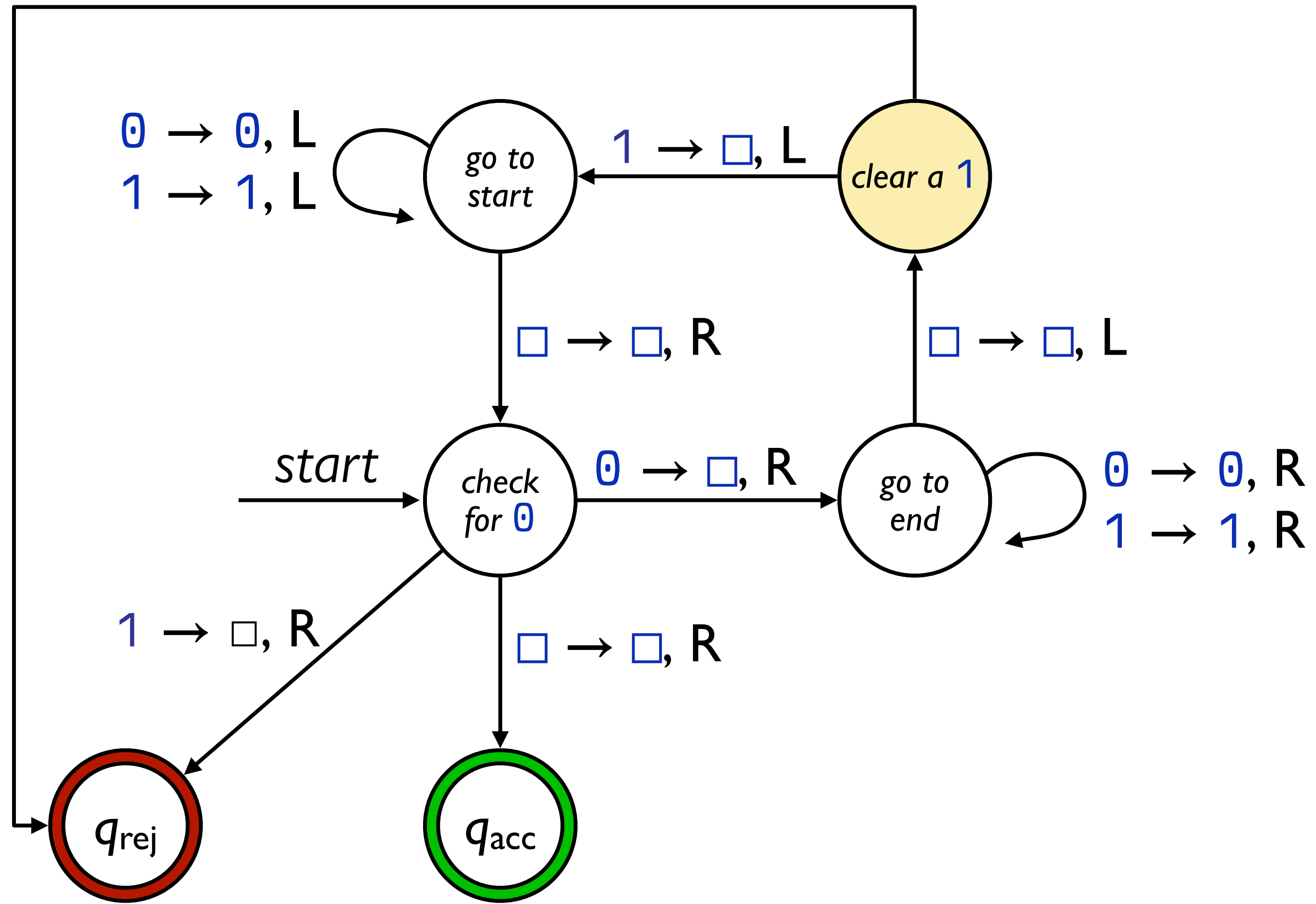






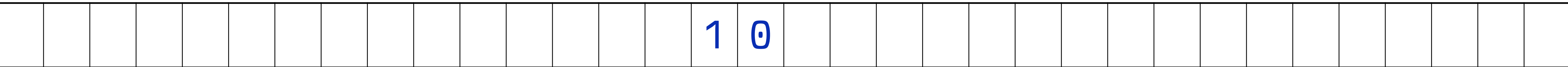
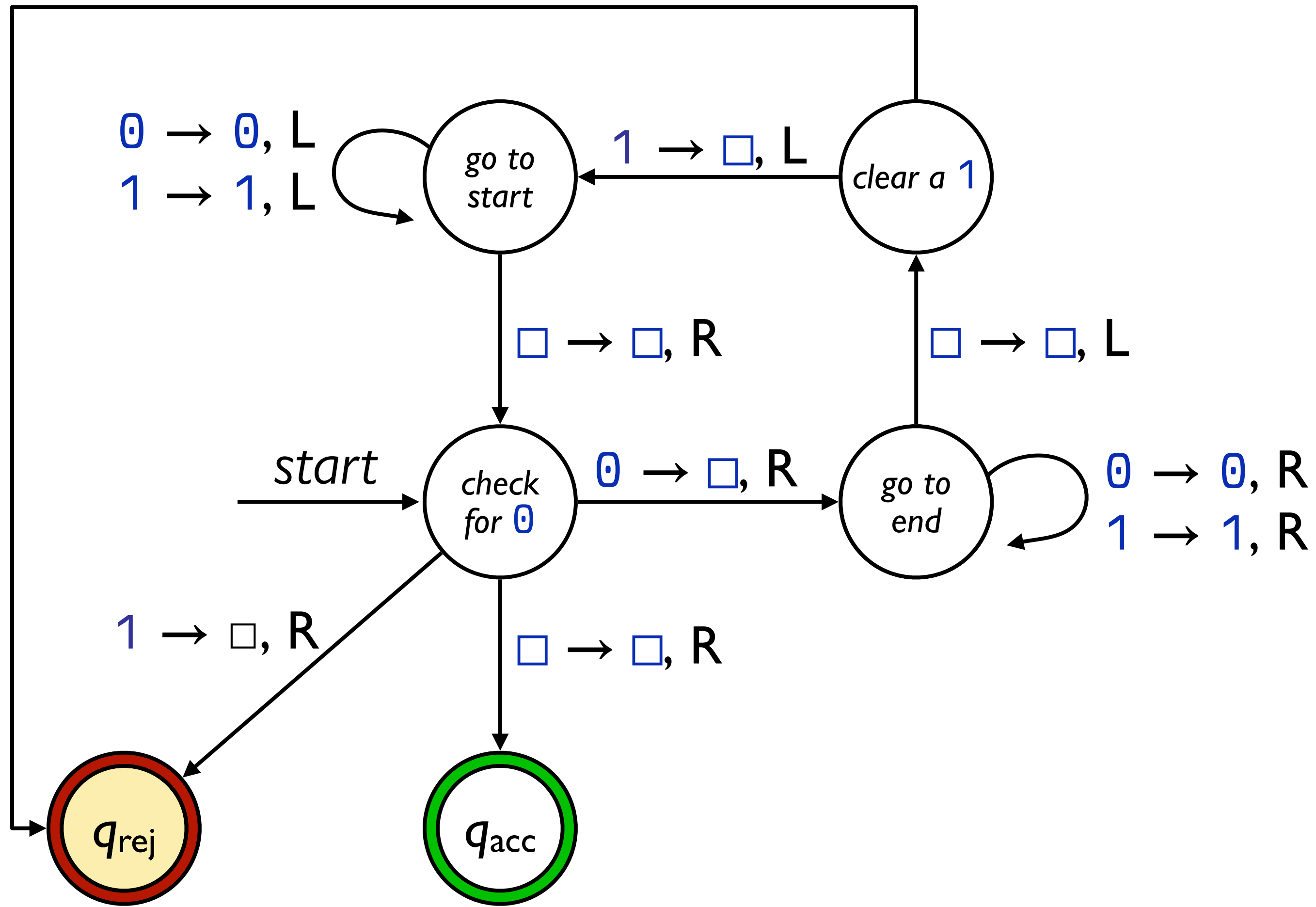
$\square \rightarrow \square, R$

$0 \rightarrow 0, R$



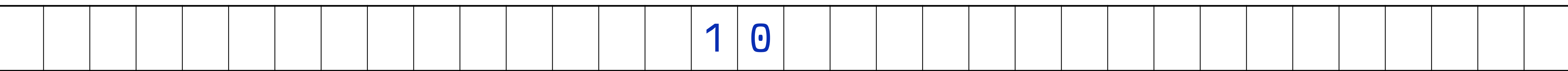
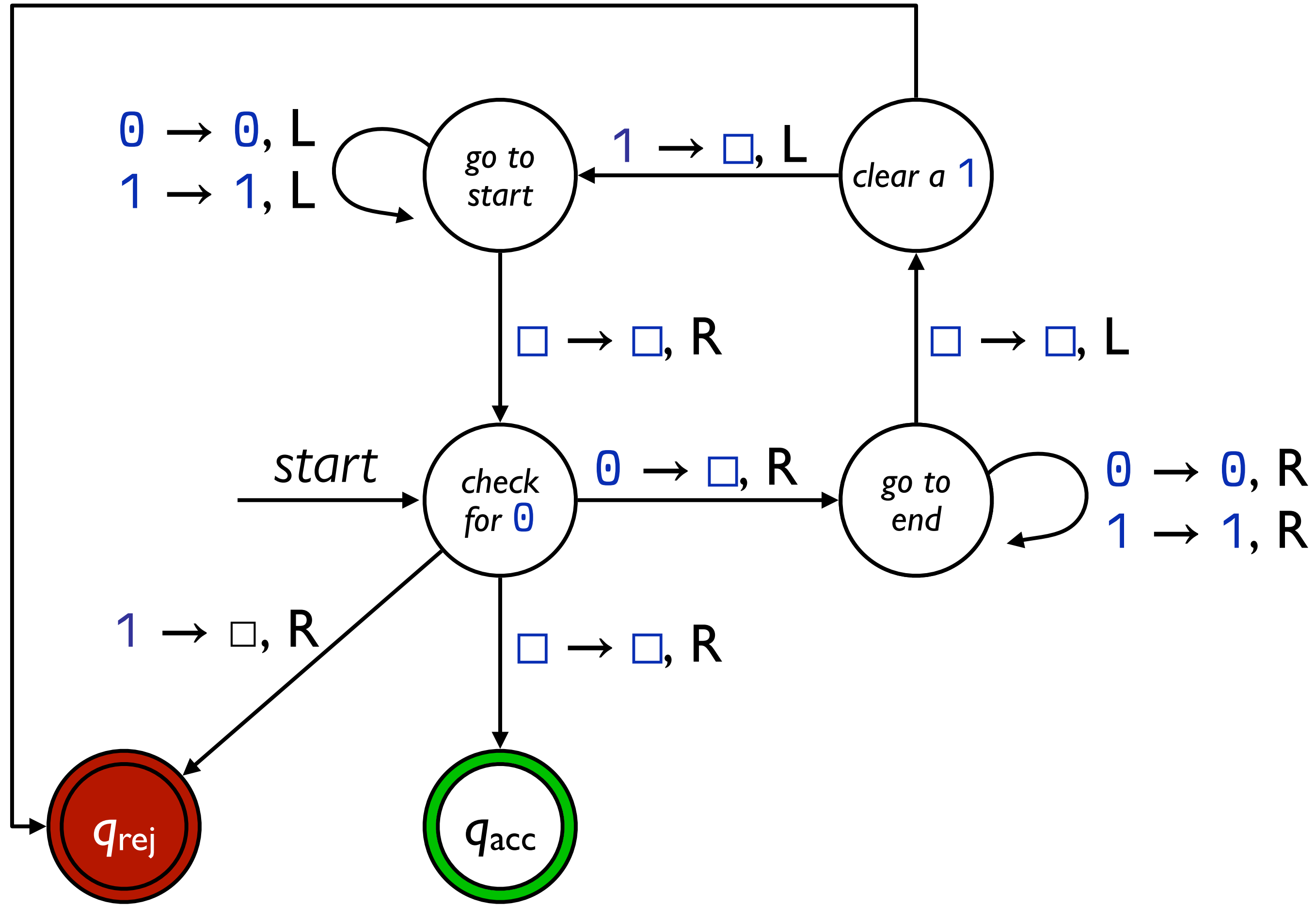
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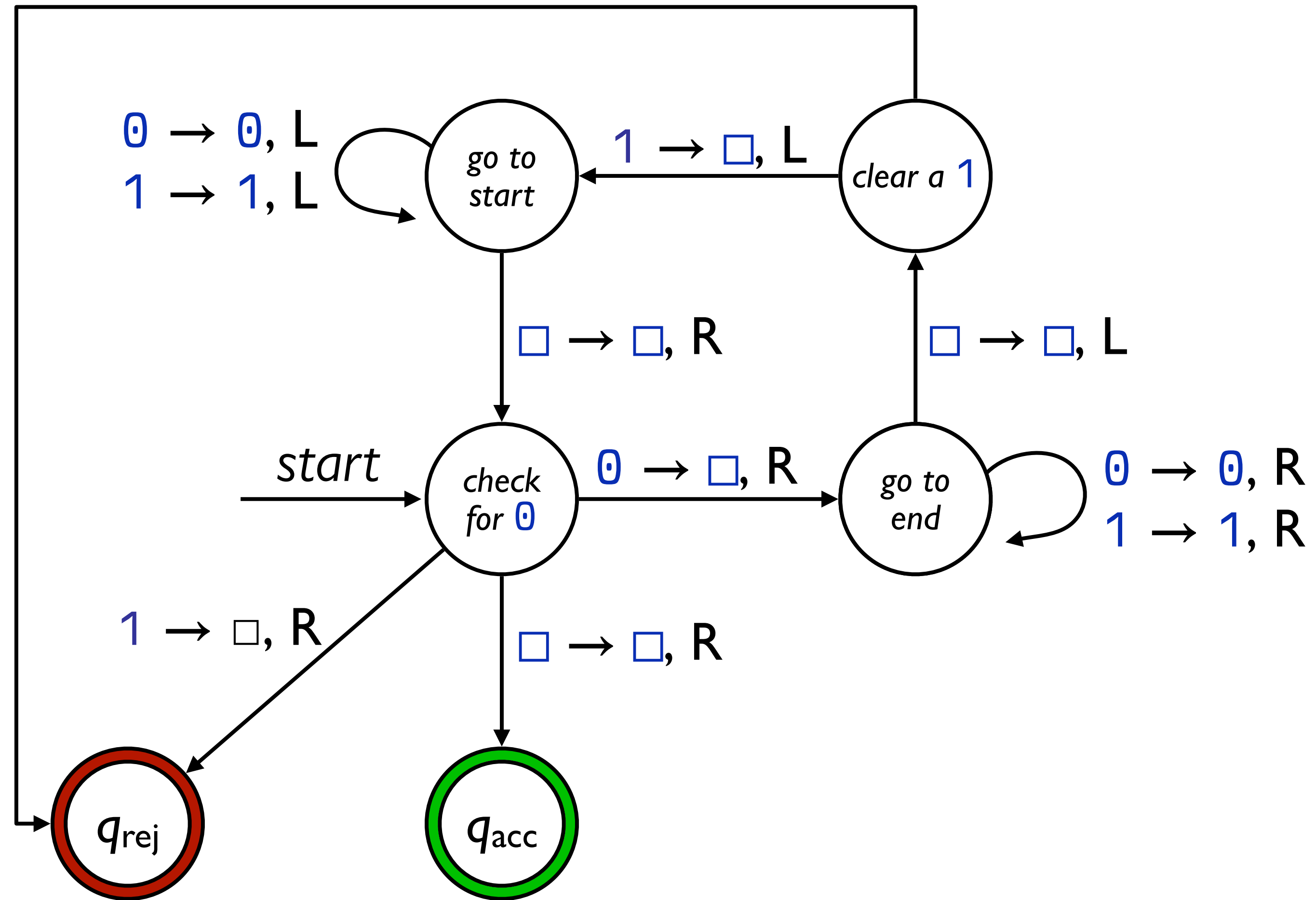
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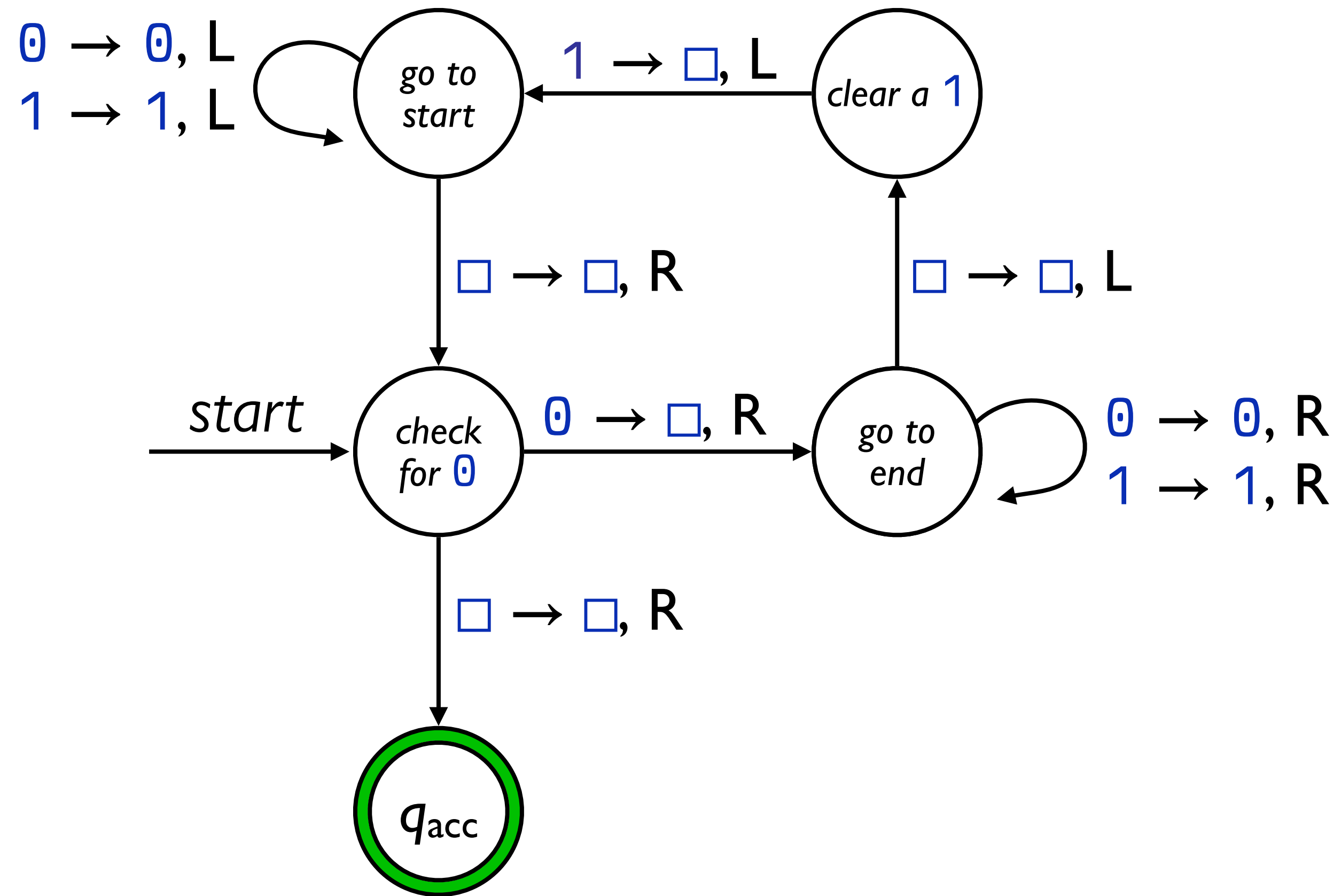
$0 \rightarrow 0, R$



$\square \rightarrow \square, R$

$0 \rightarrow 0, R$





# Another Turing machine design

We designed a Turing machine for  $\{0^n 1^n \mid n \in \mathbb{N}_0\}$ .

Let's consider how we could design a Turing machine for a related context-free language over

$\Sigma = \{0, 1\}$ :

$$L = \{w \in \Sigma^* \mid n_0(w) = n_1(w)\}$$











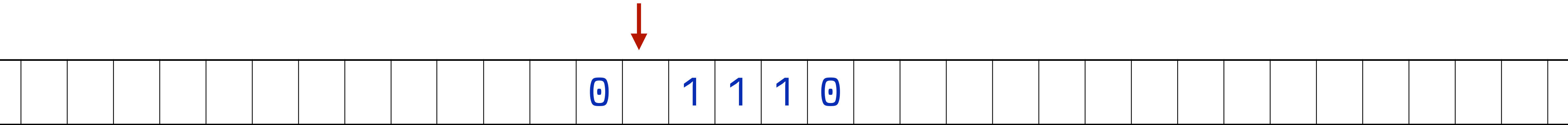








# A caveat



*How do we know that this blank isn't one of the infinitely many blanks after our input string?*

































































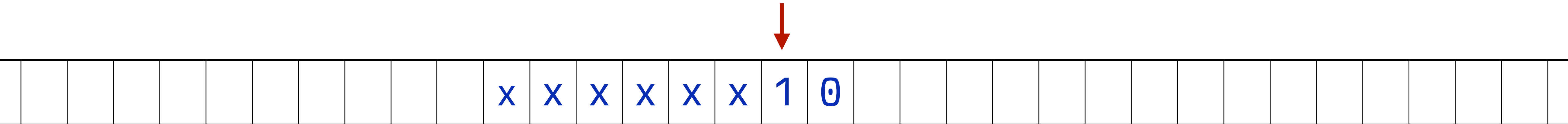








# One solution

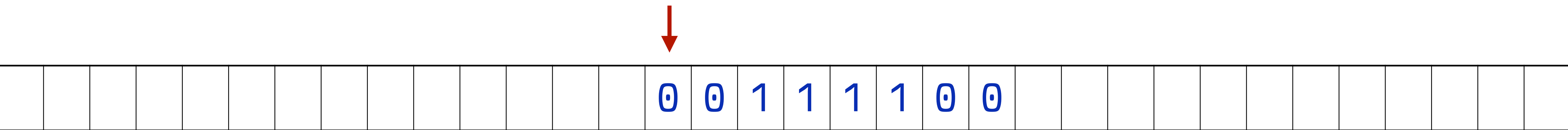
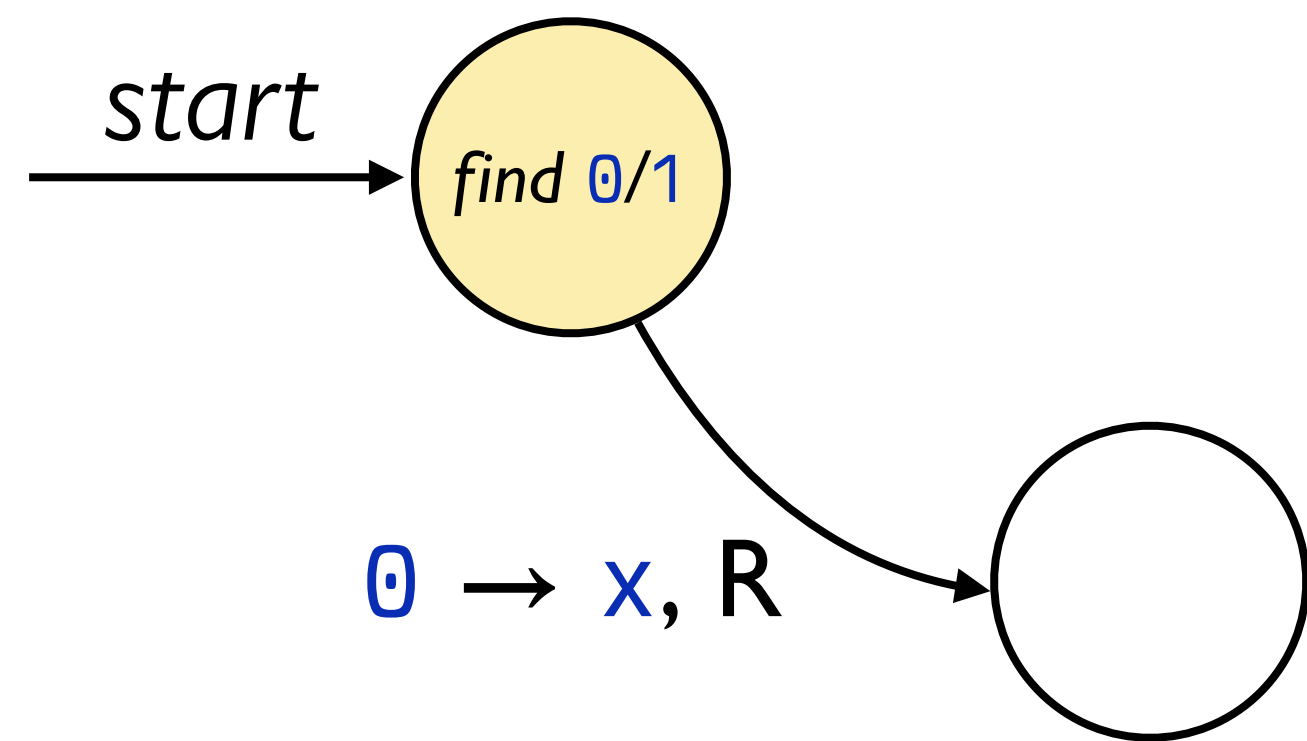


*Now the first non-x character we encounter is a 1, so we want to cross it off but remember that we're looking for a matching 0*





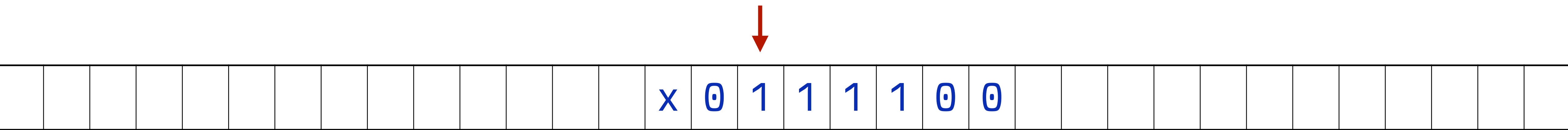
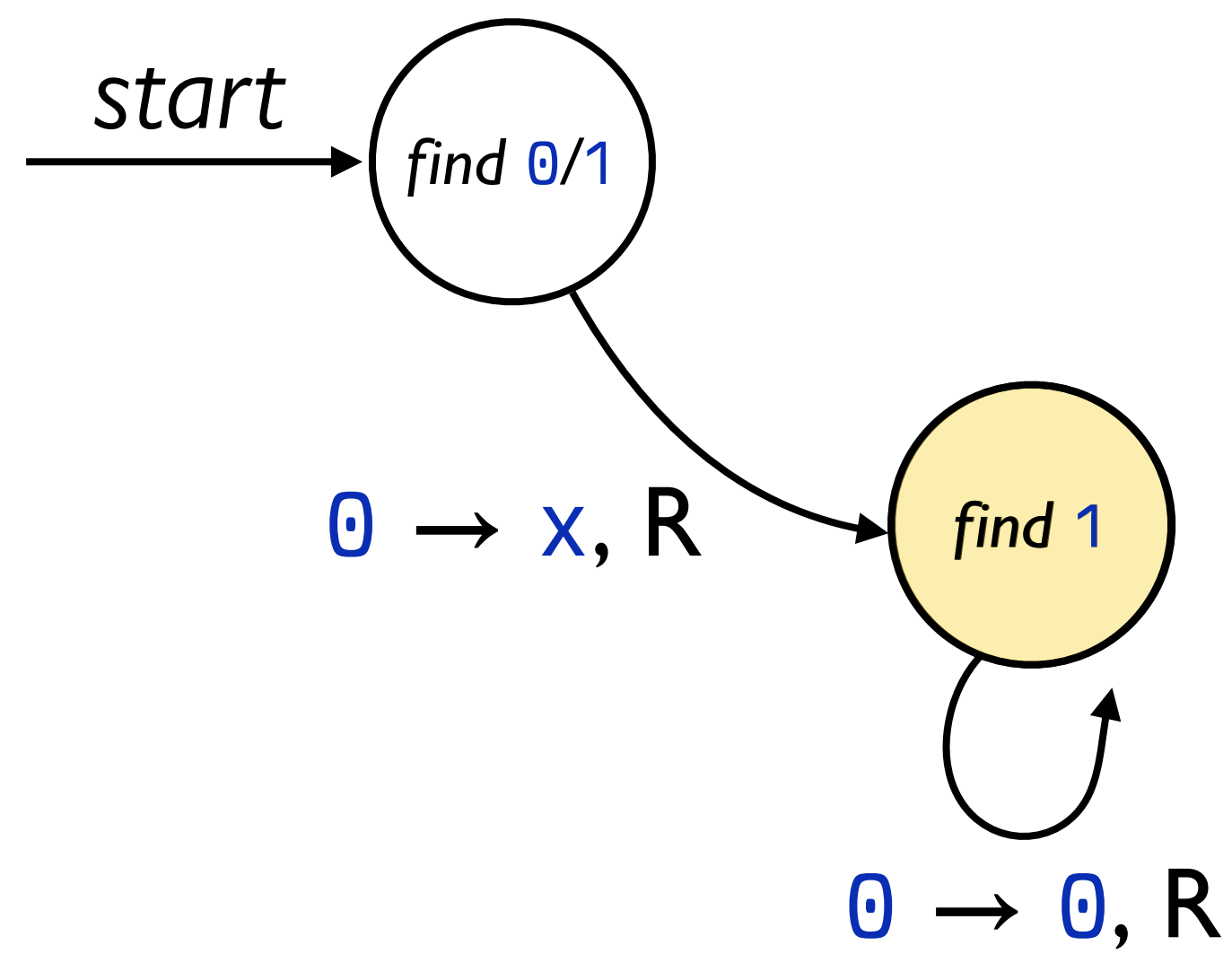


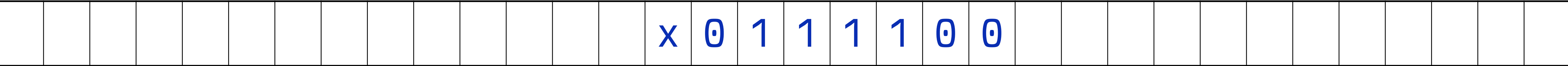
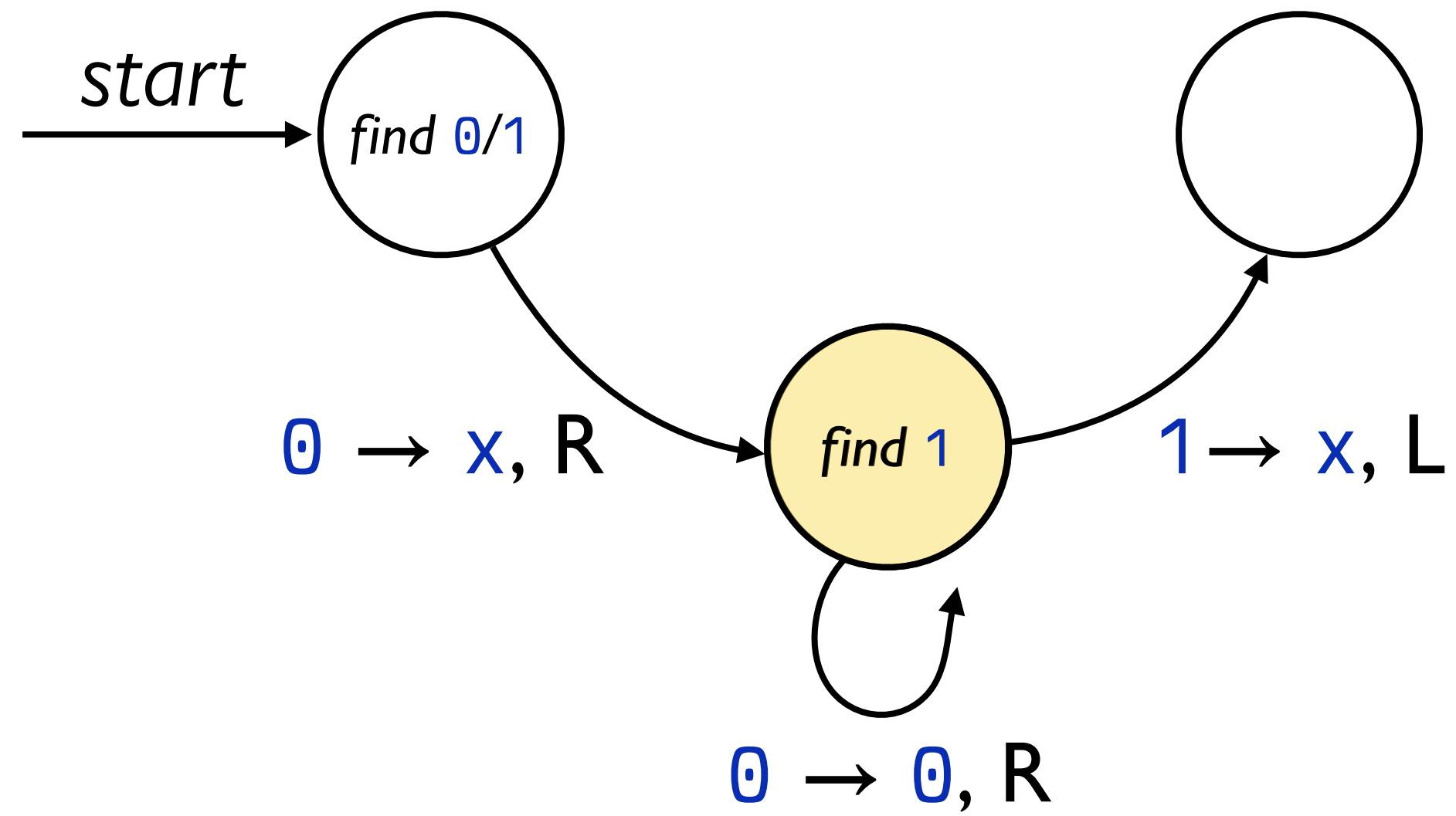




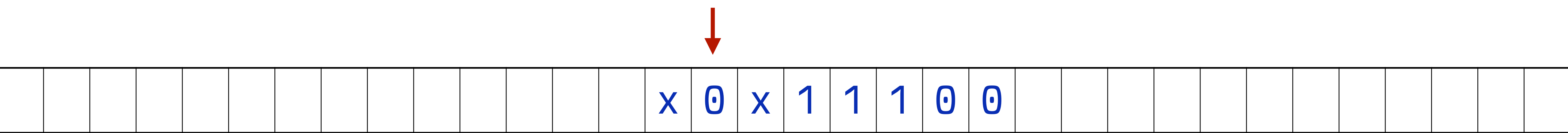
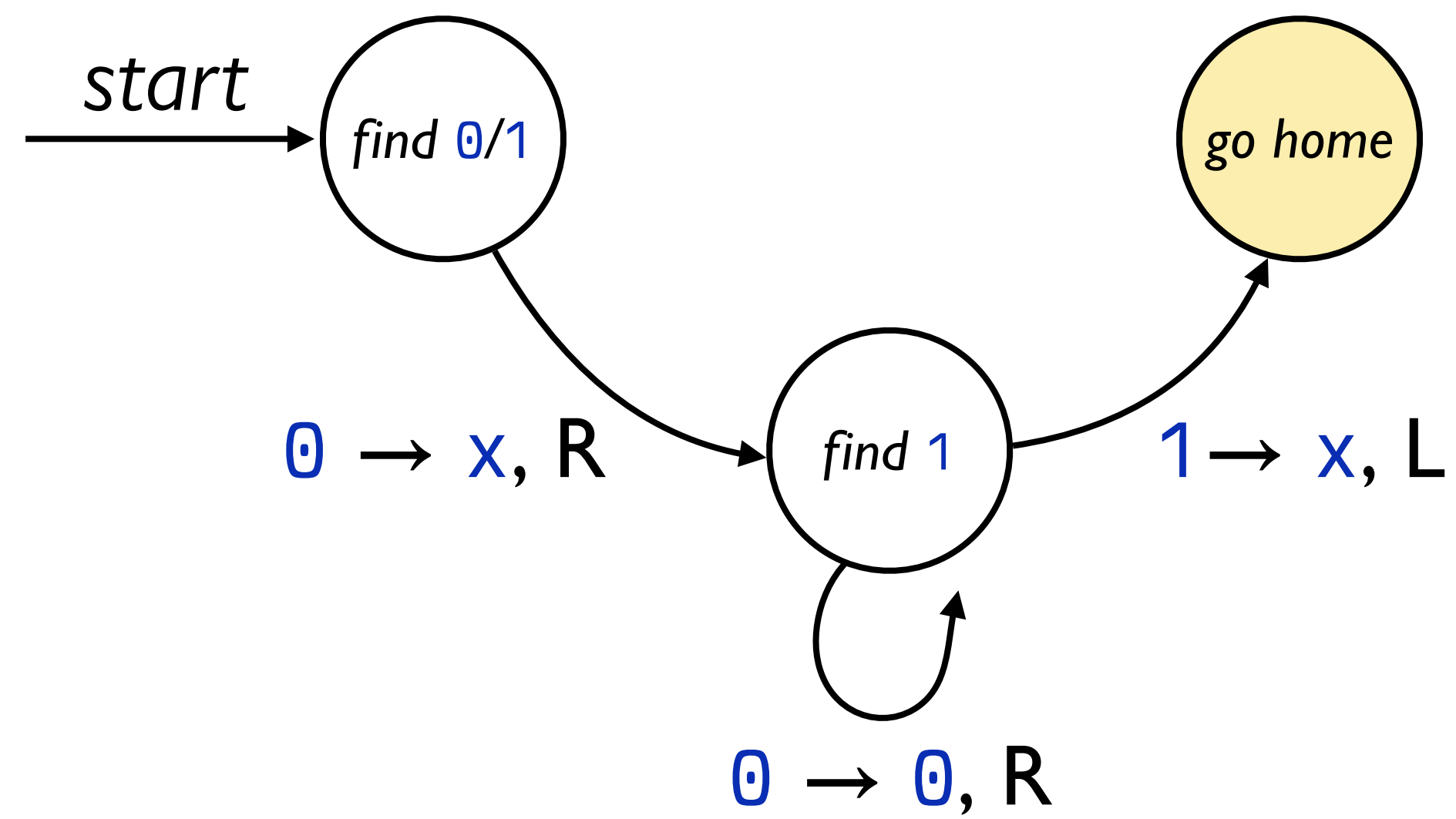






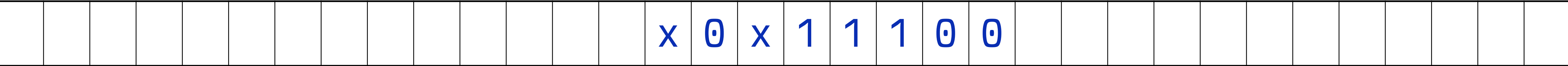
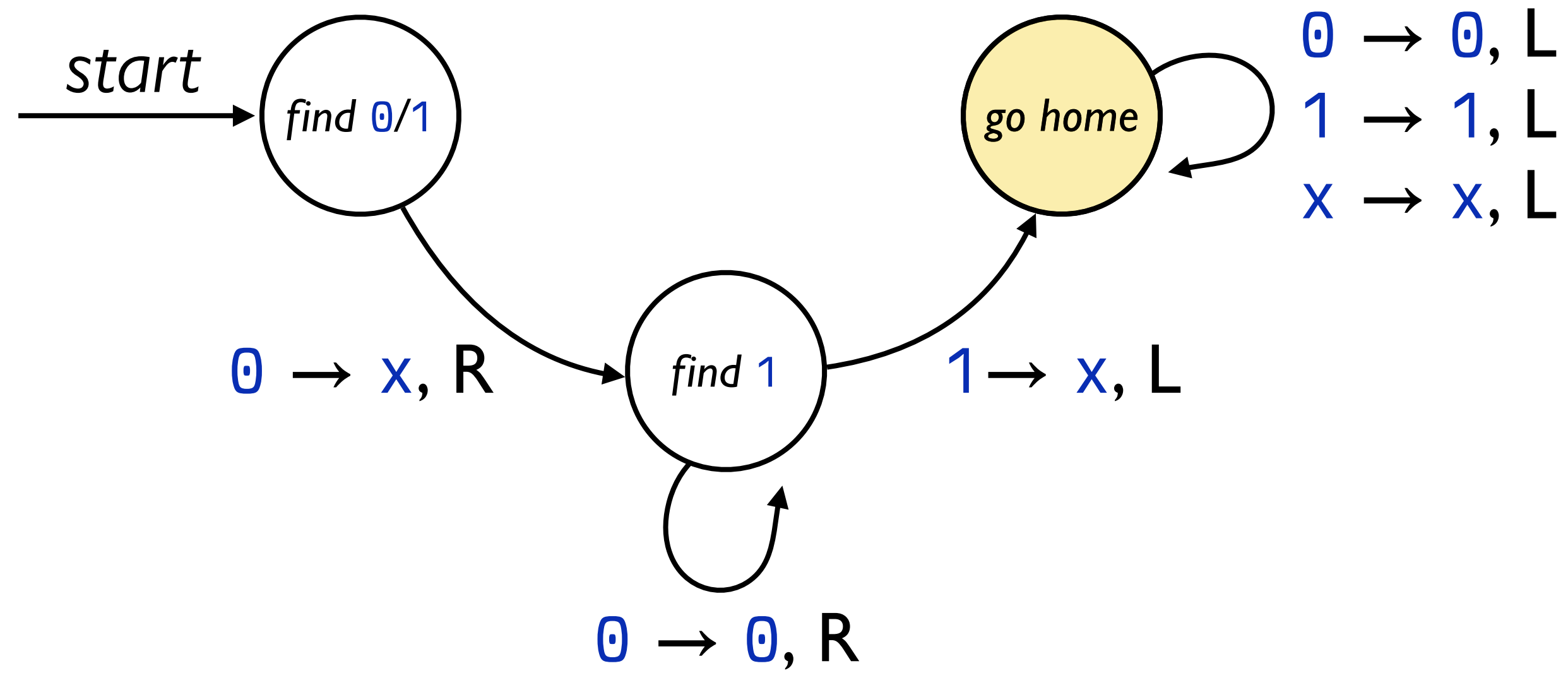


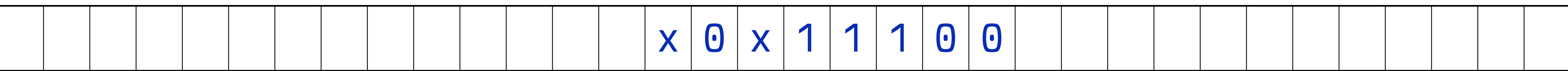
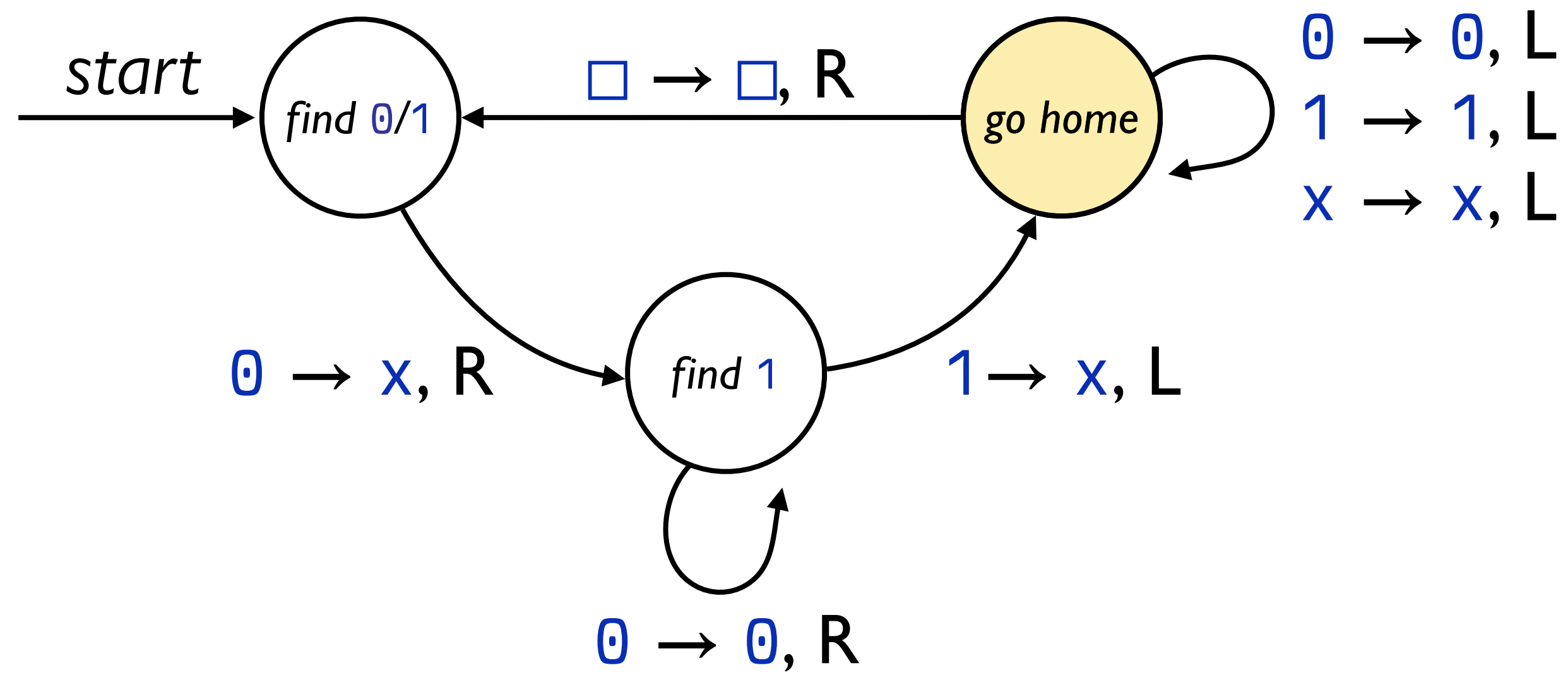










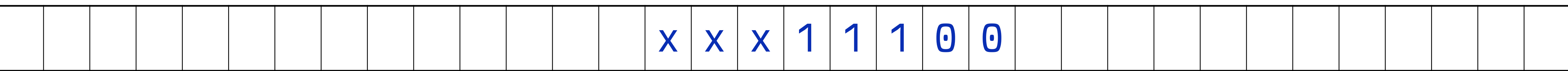
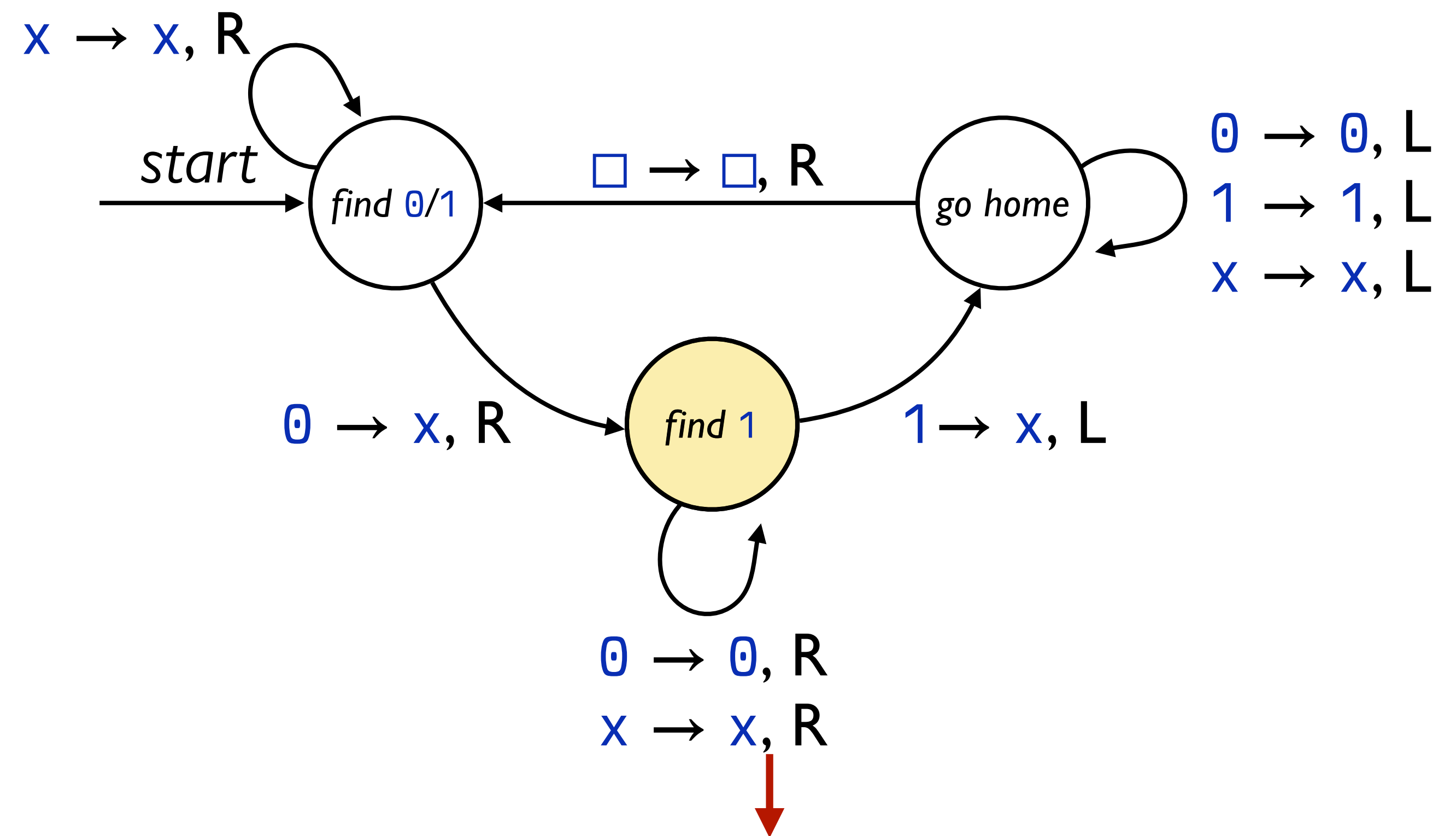


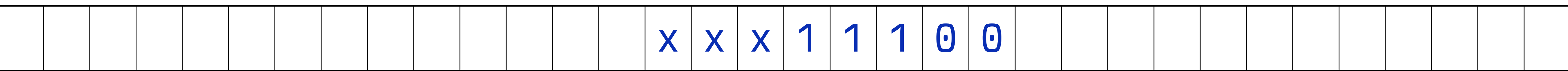
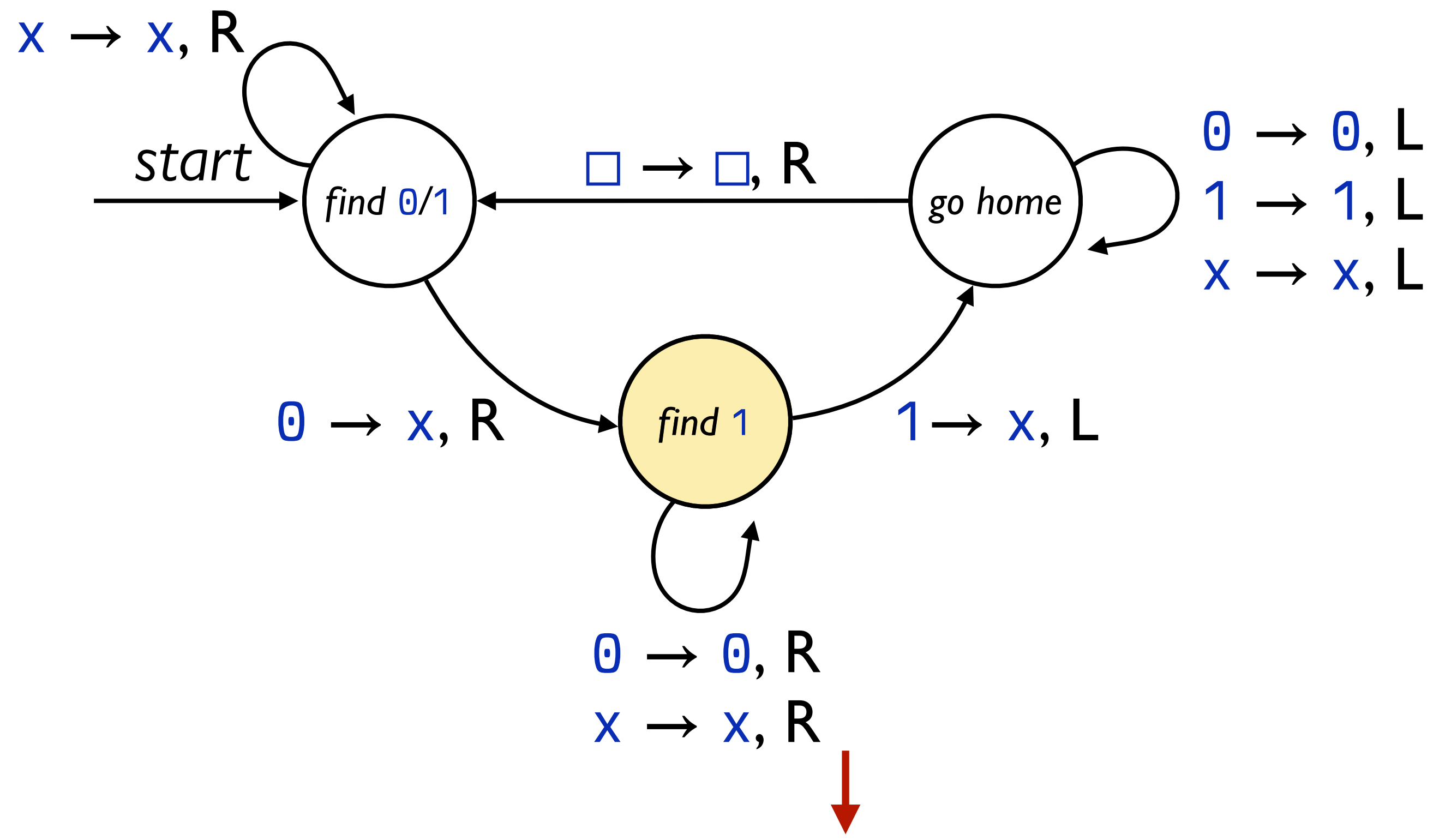


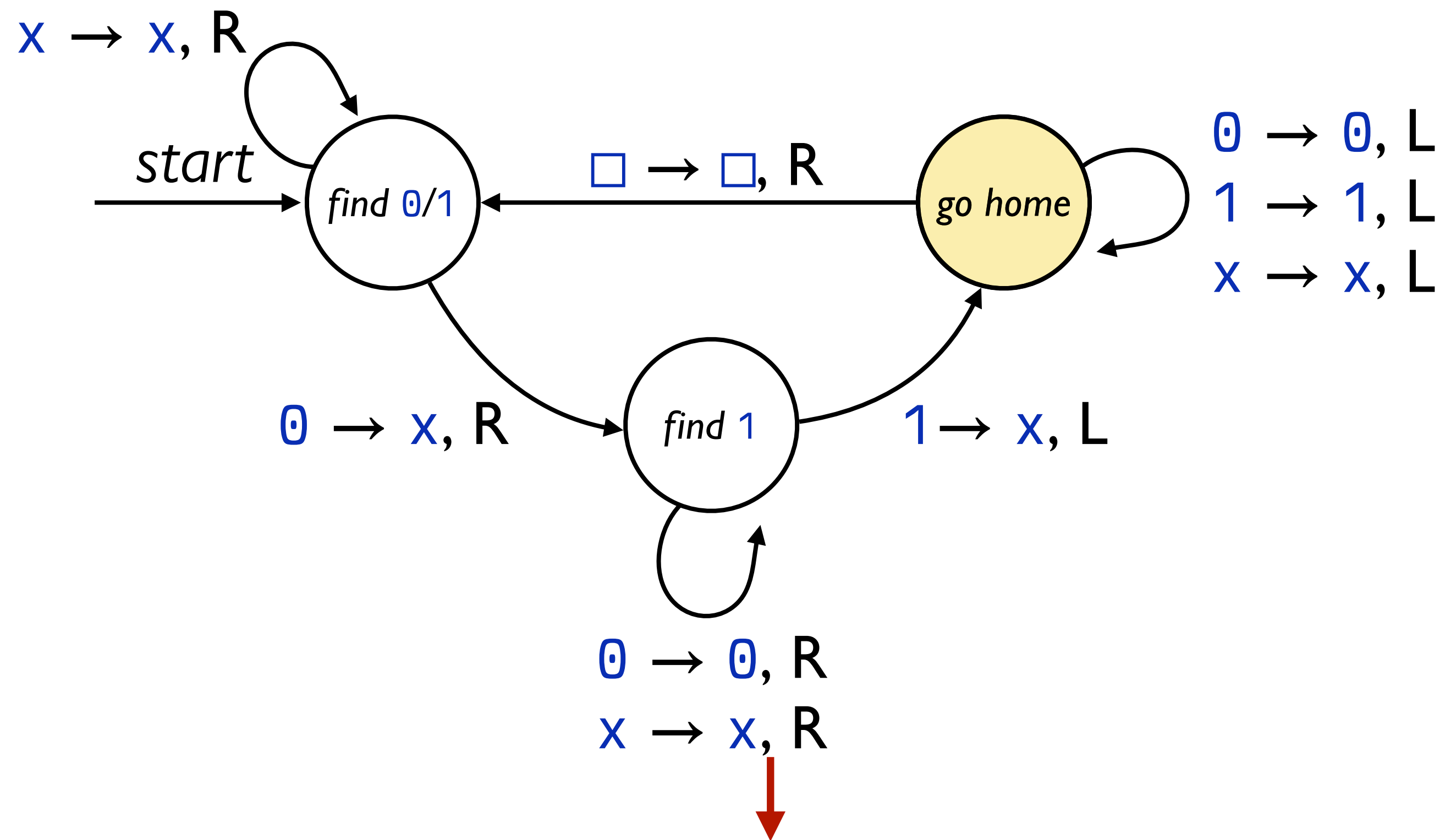








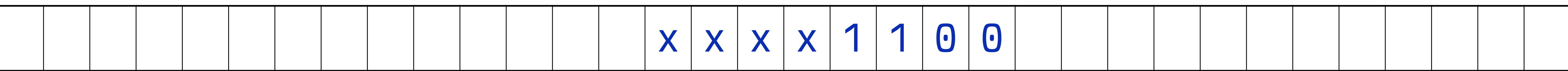
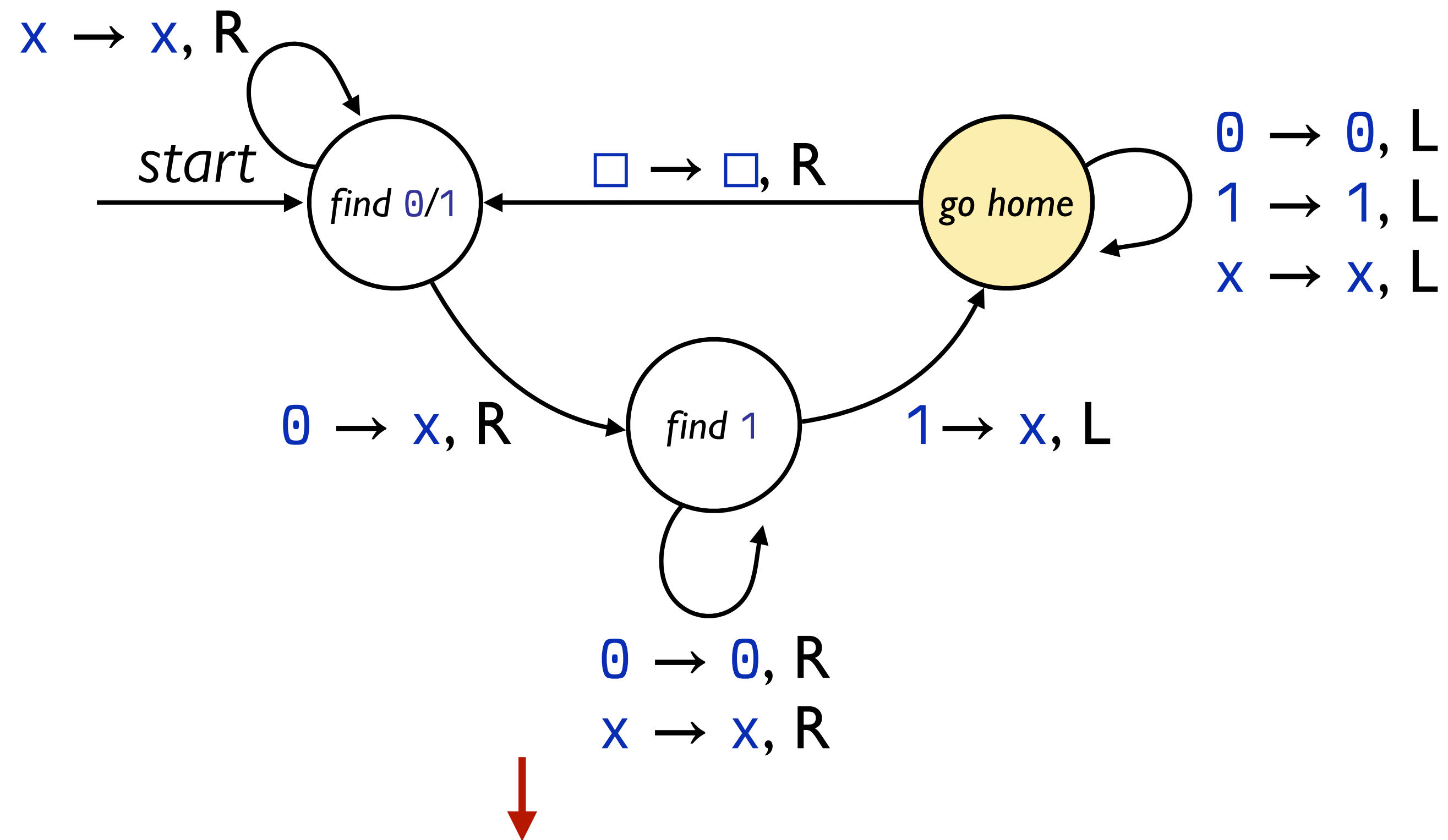




x x x x 1 1 0 0



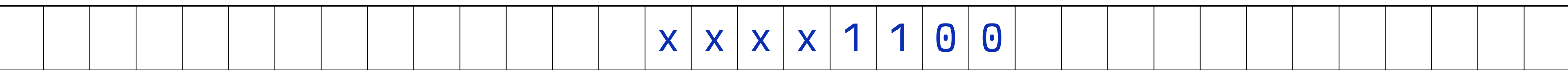
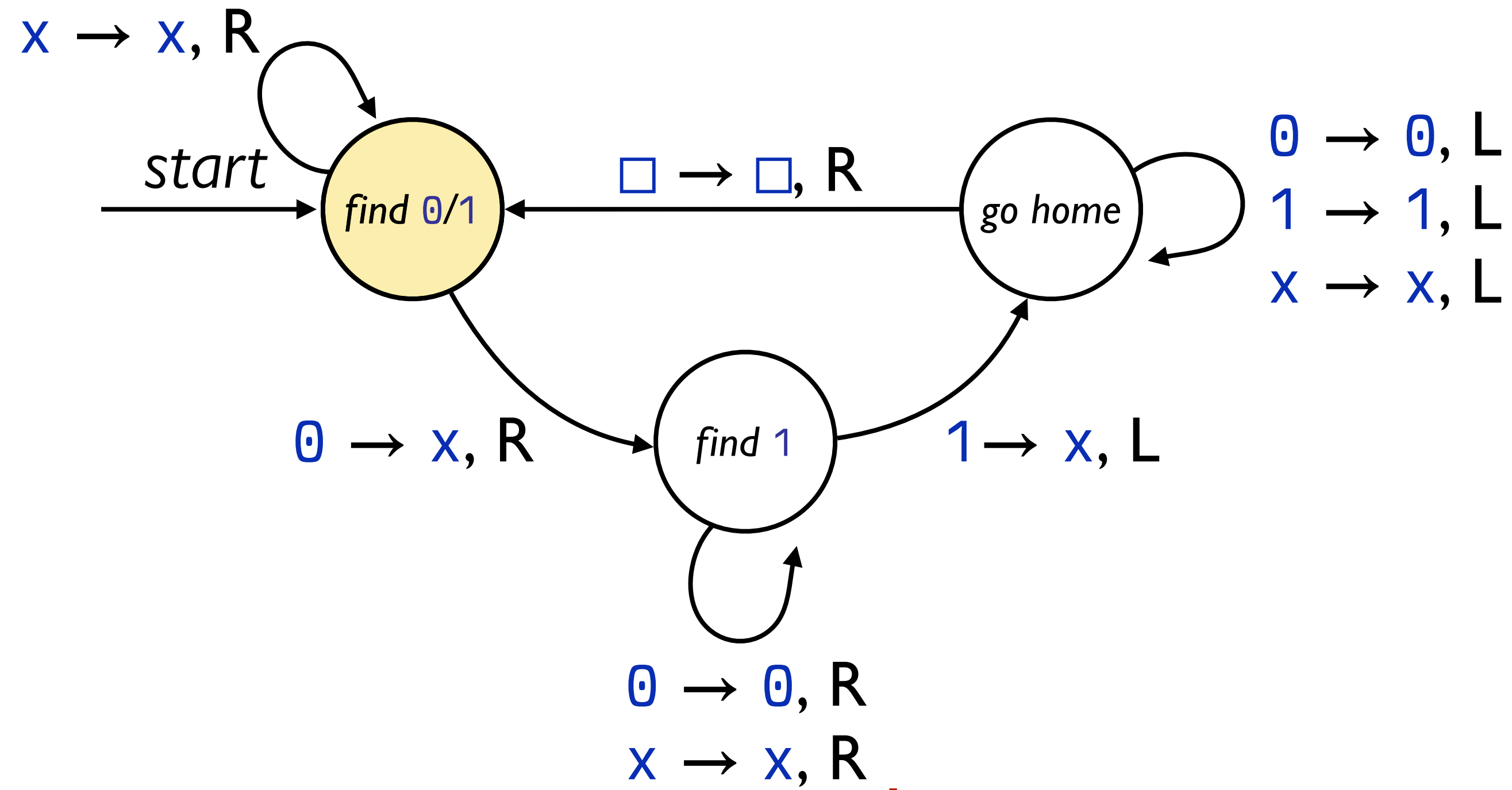


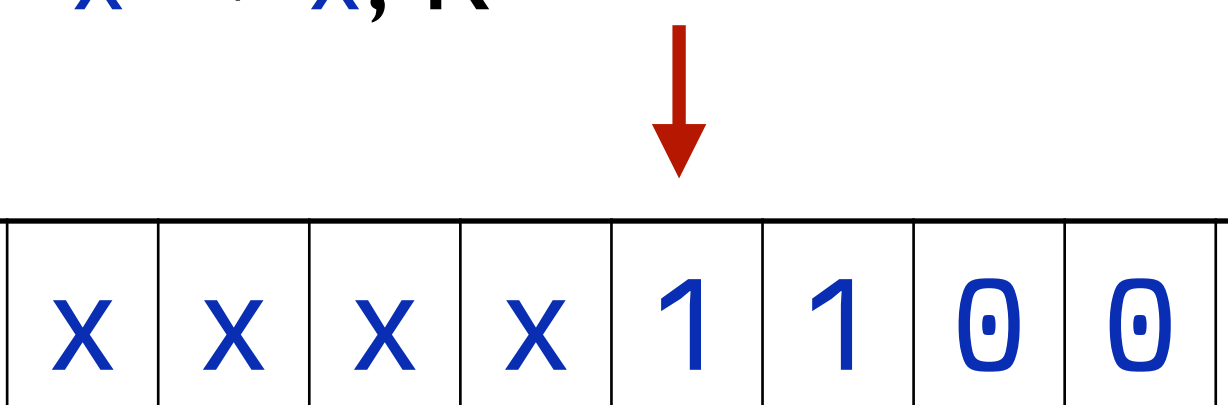
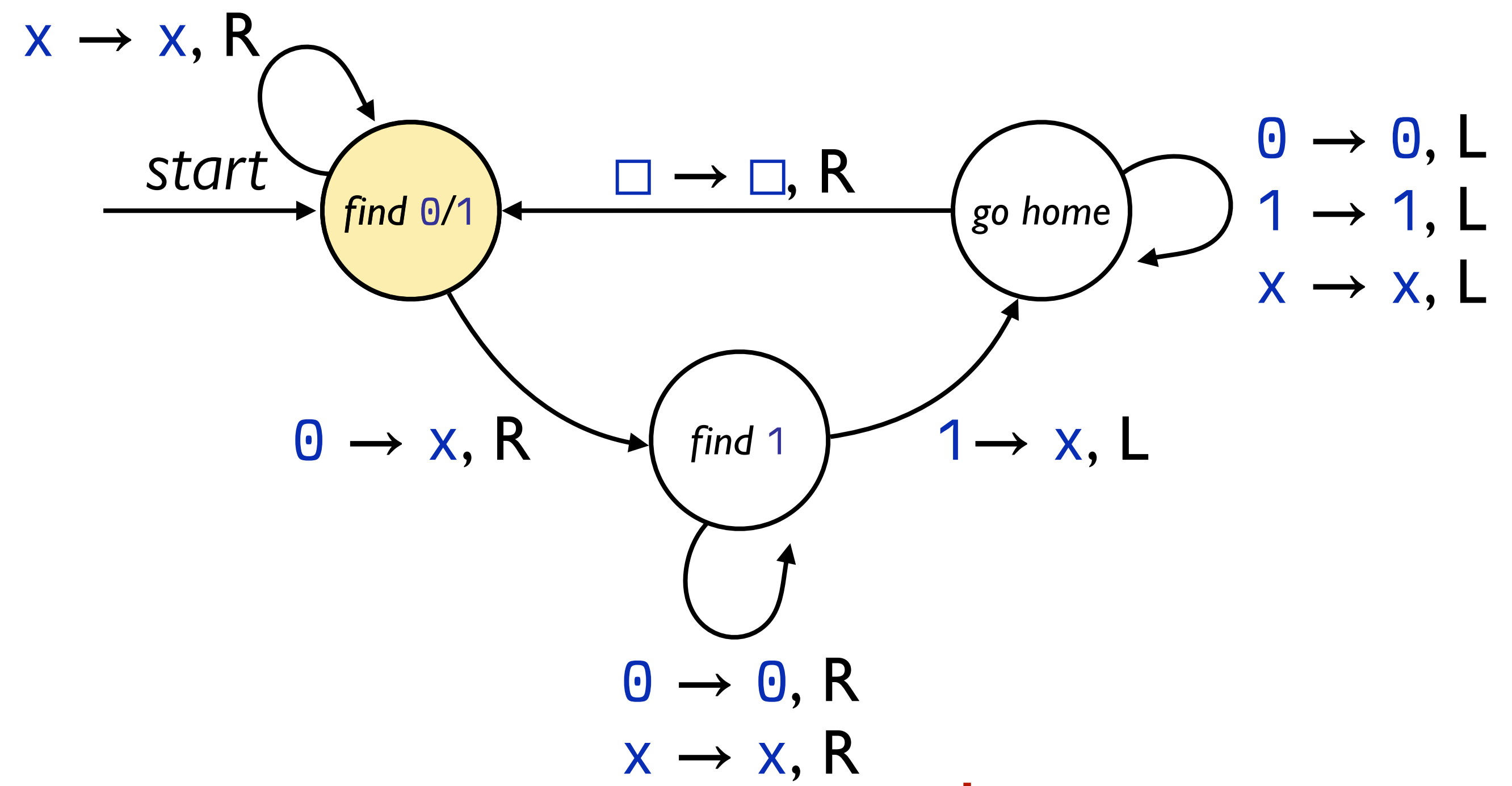


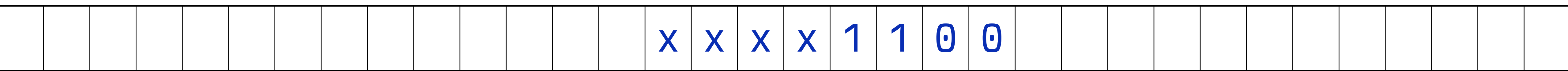
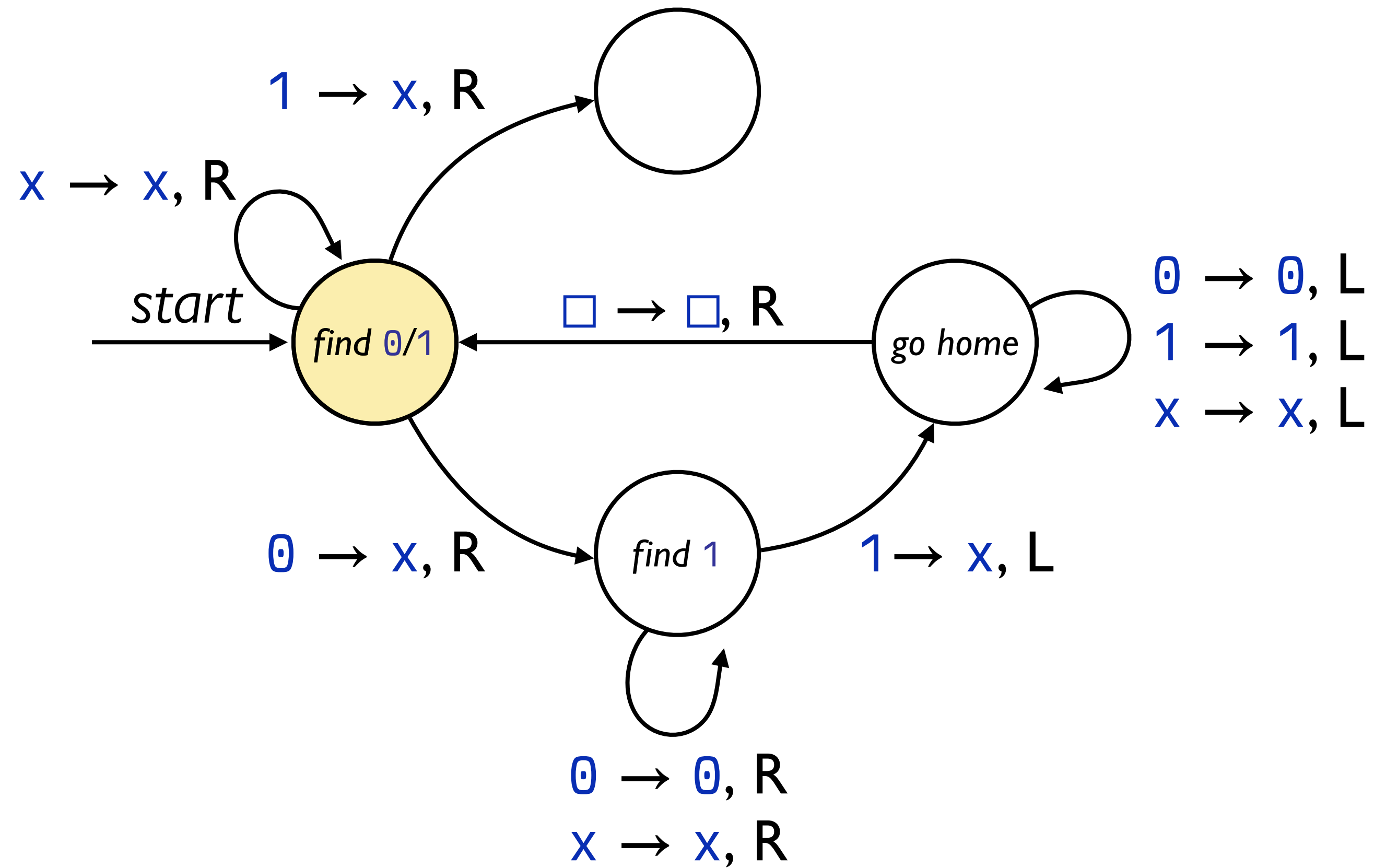


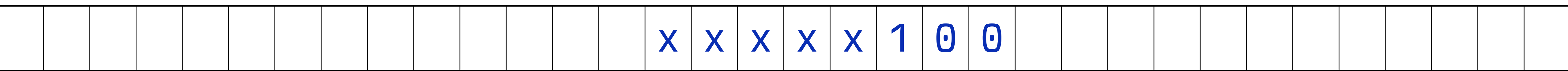
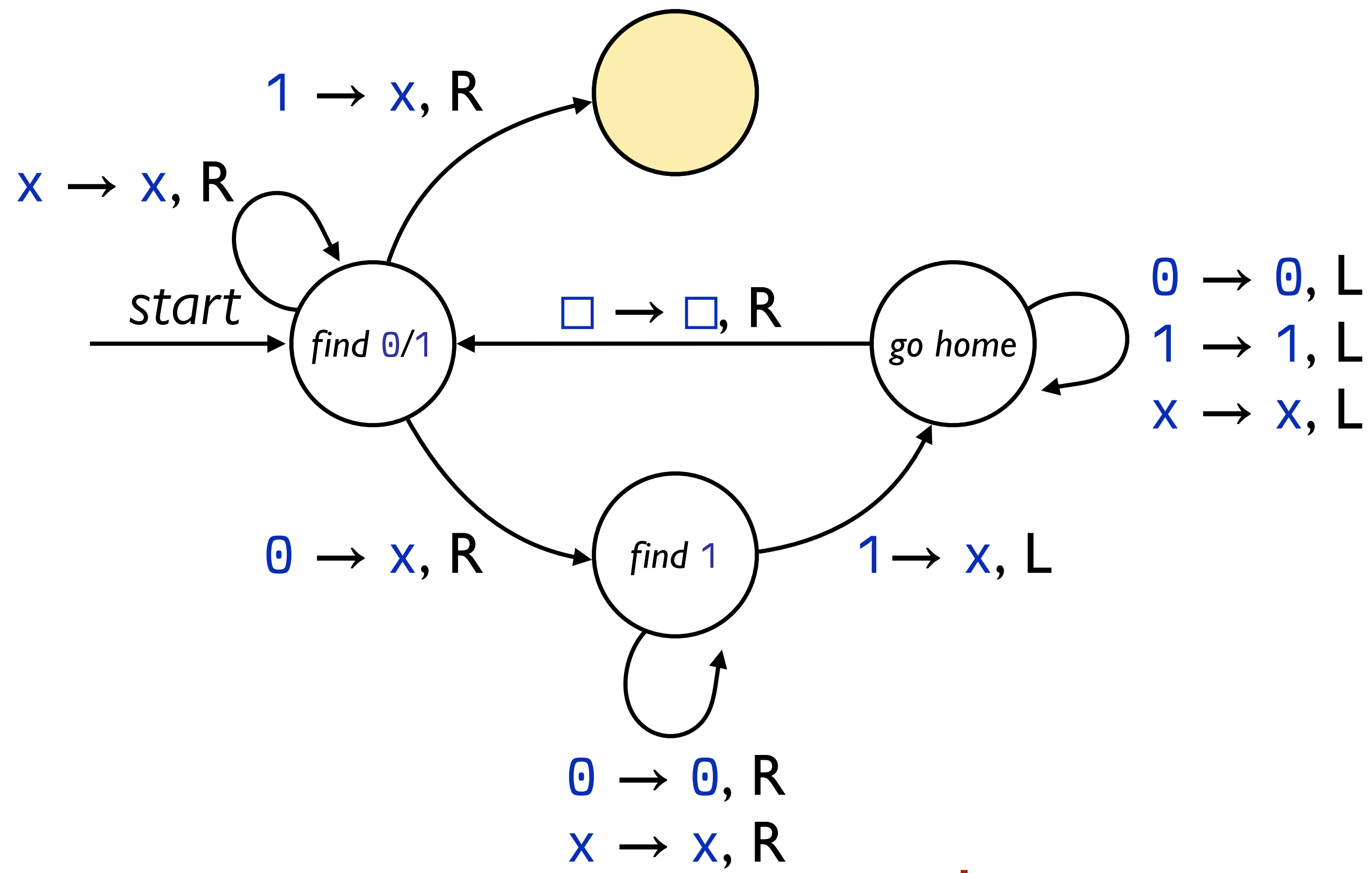






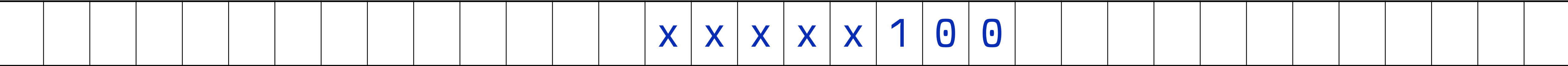
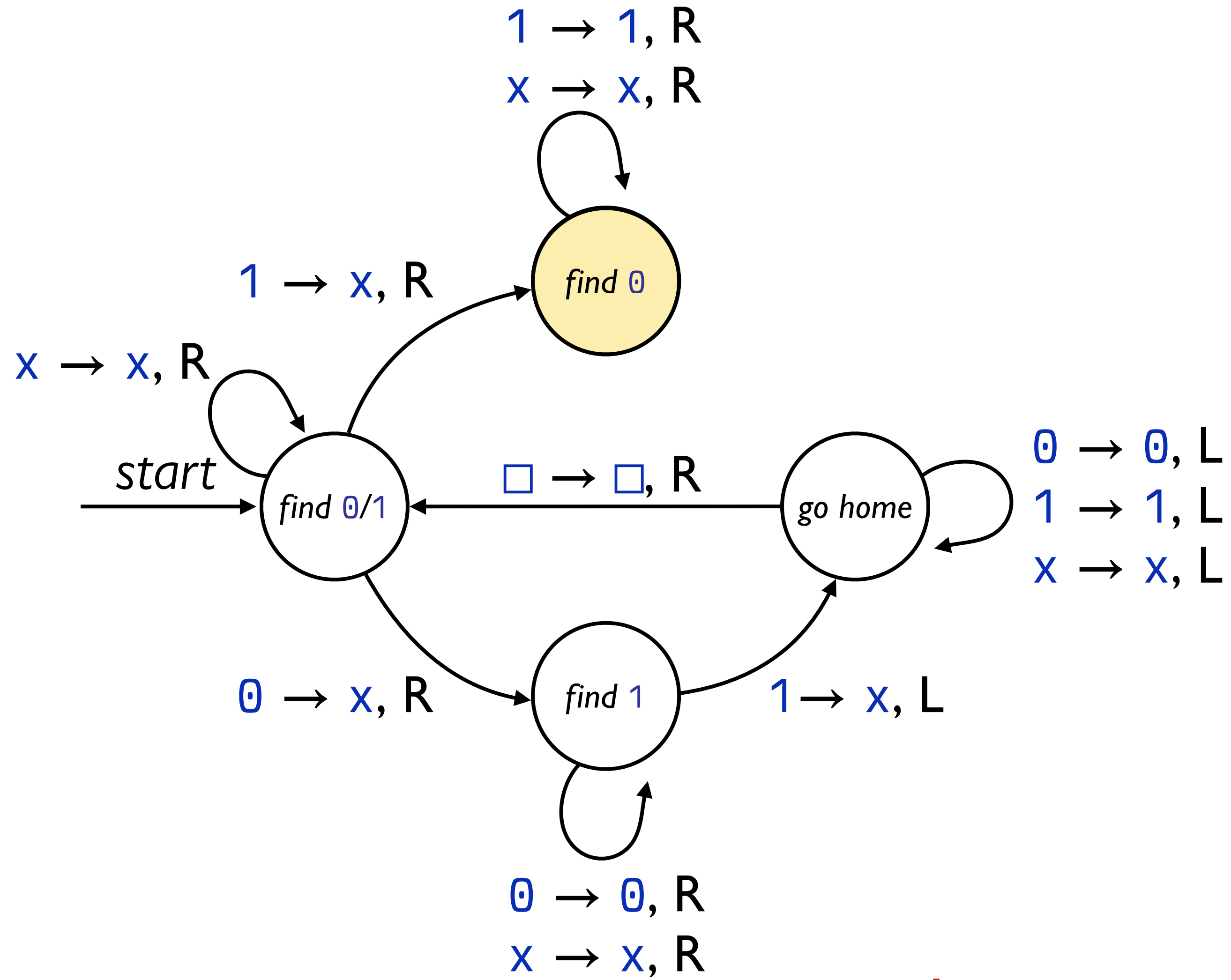






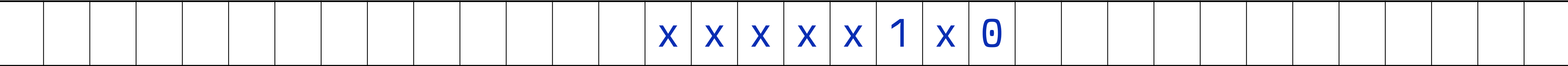
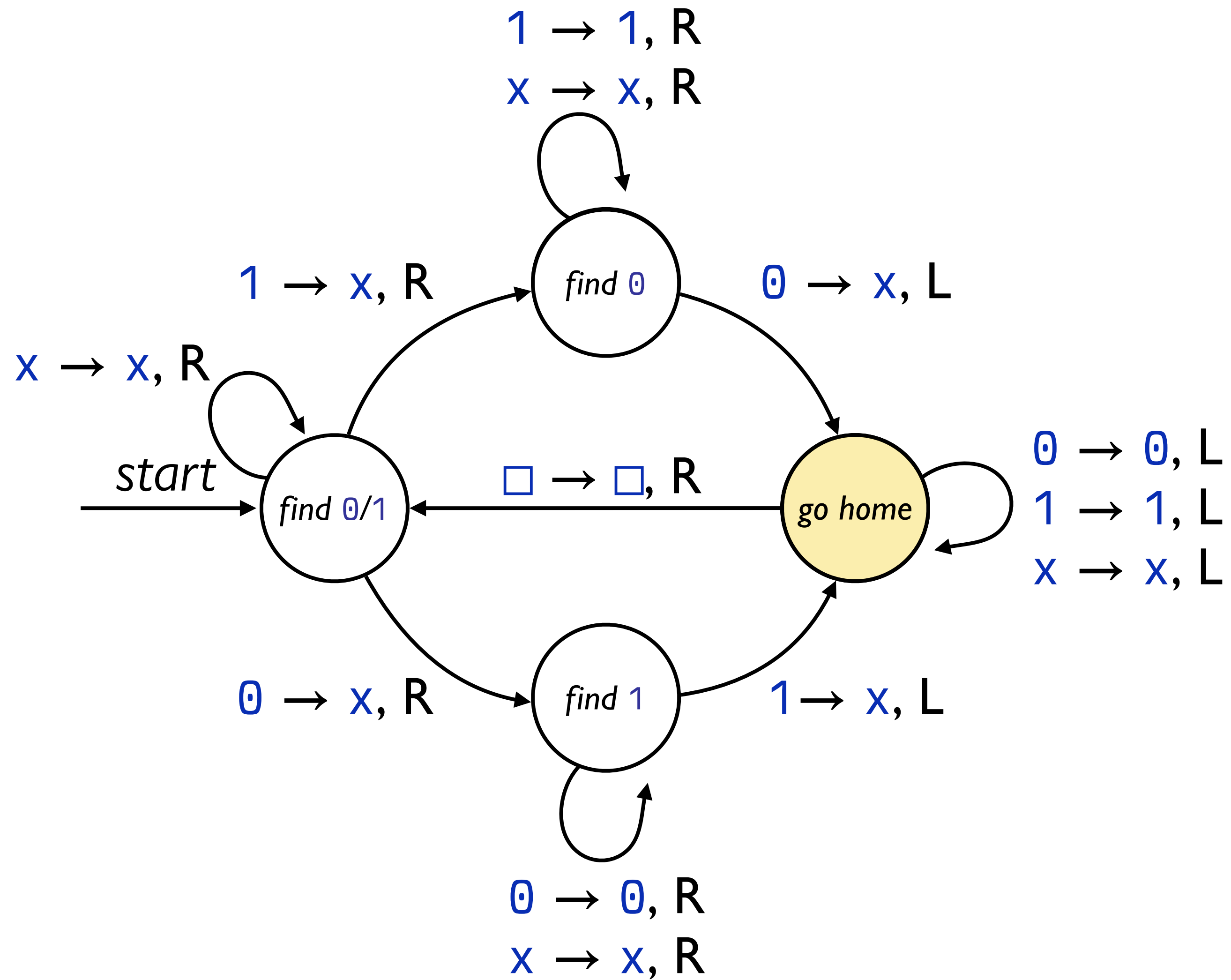


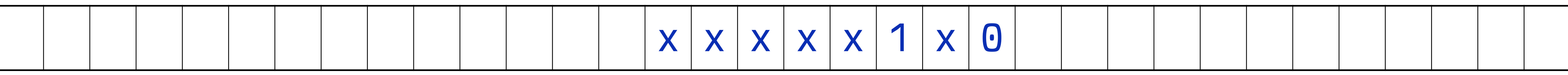
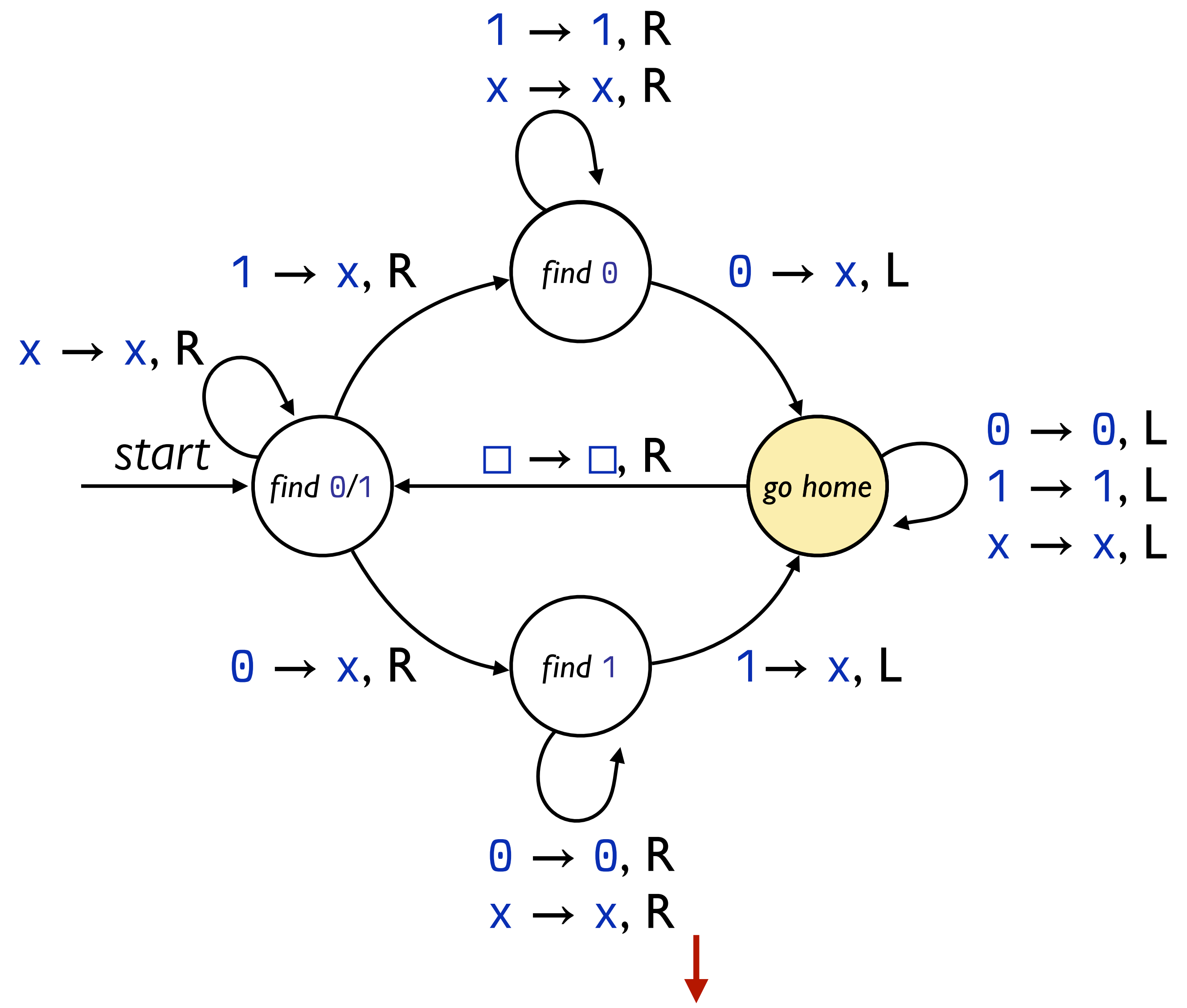




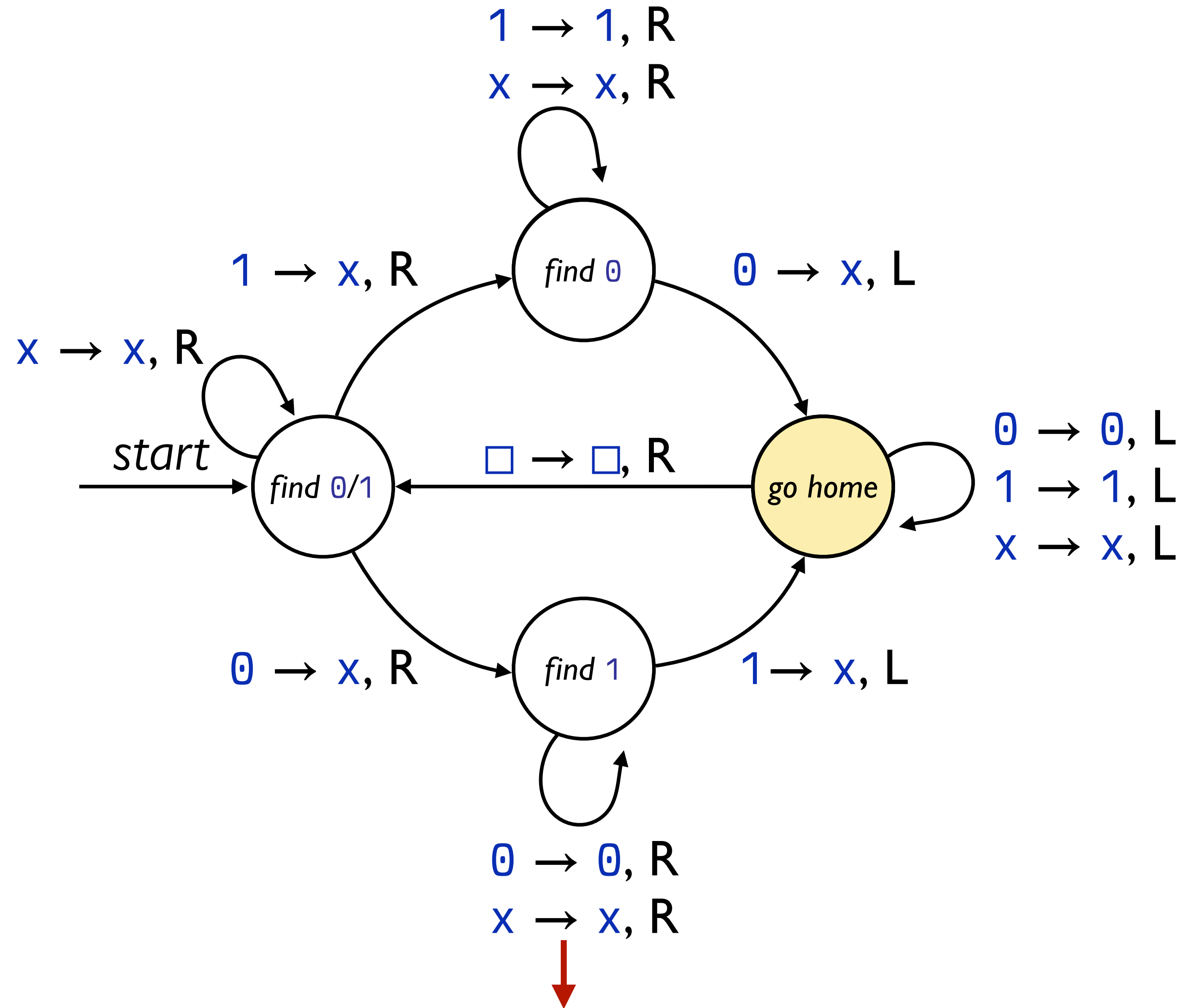








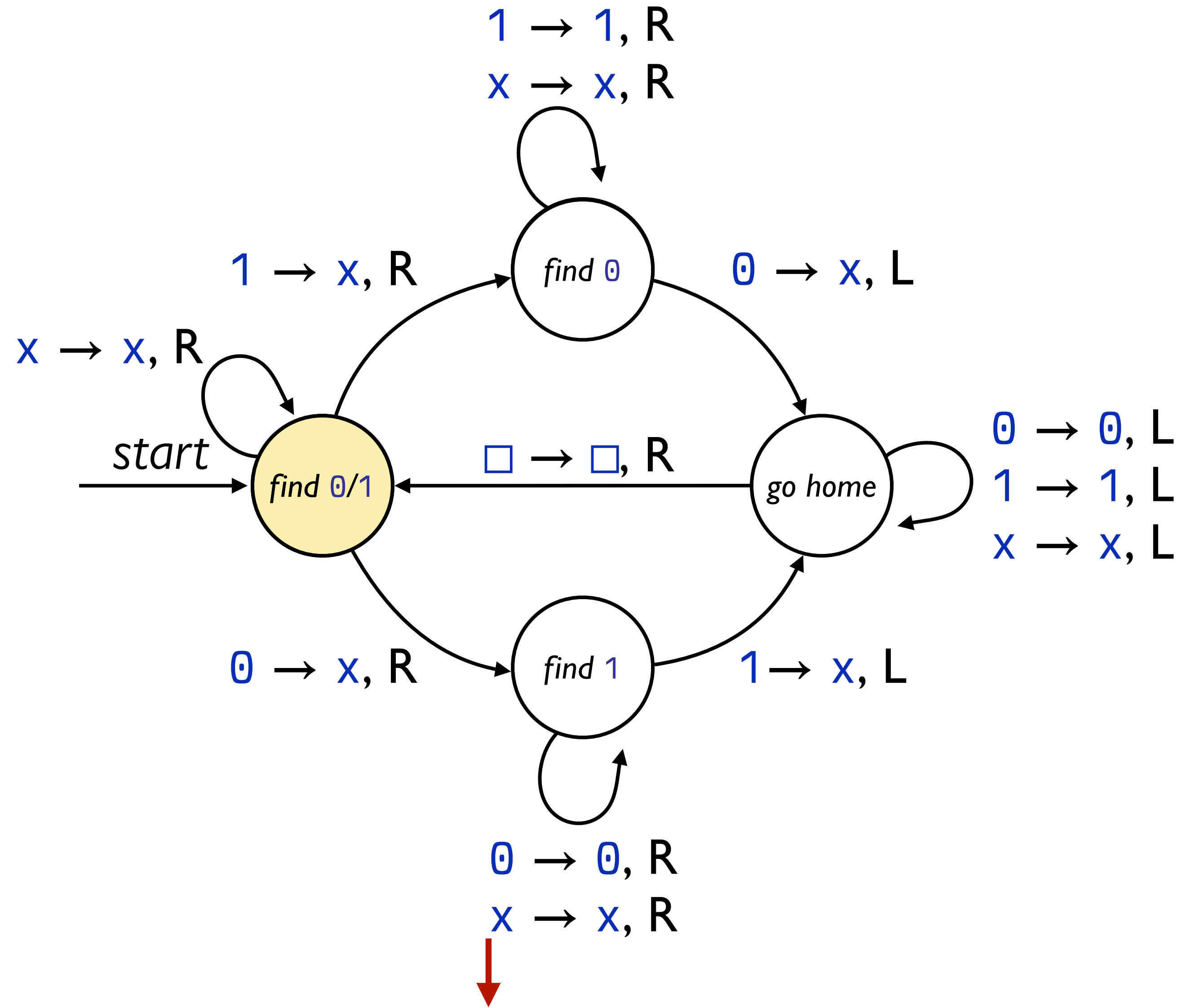




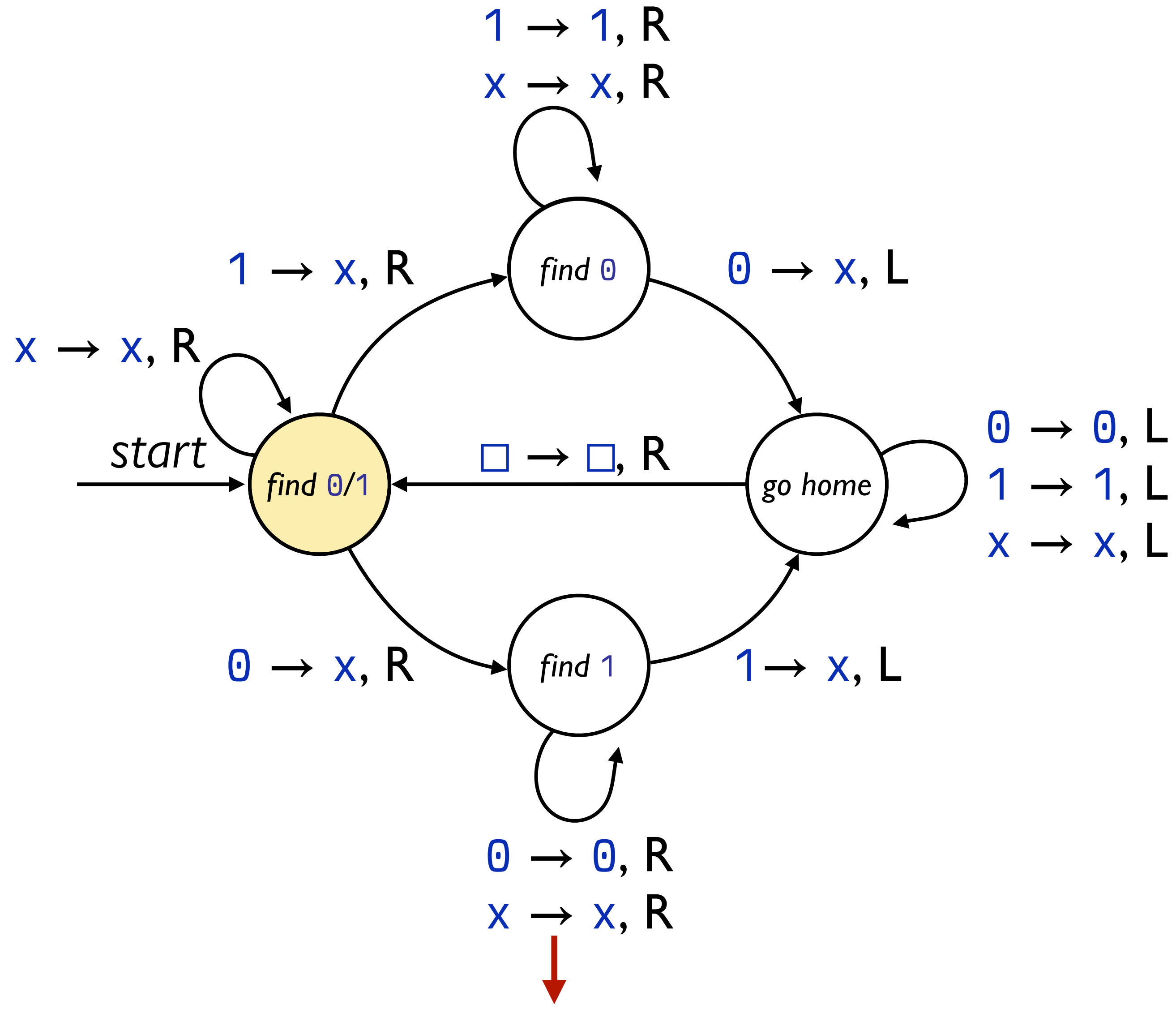
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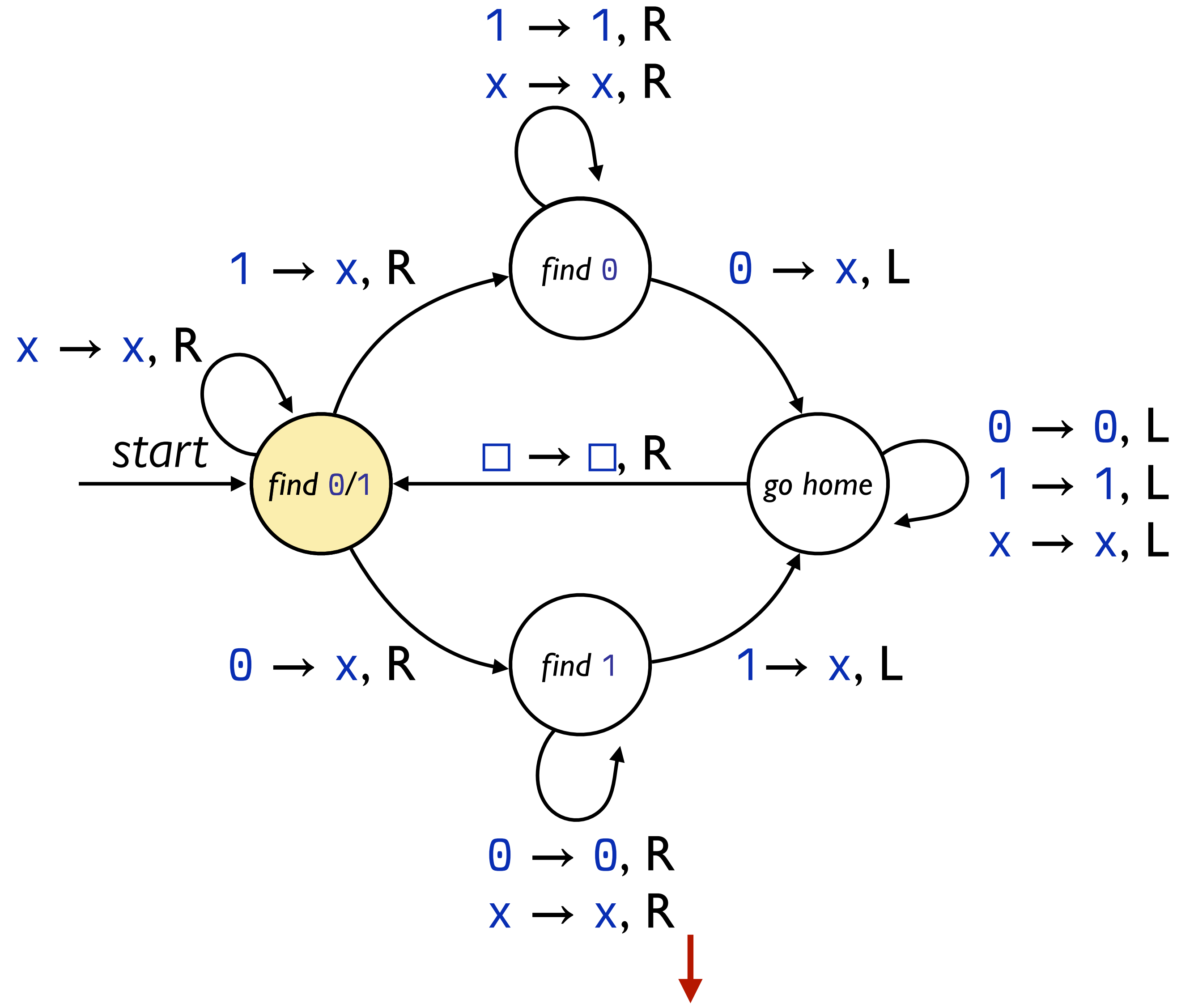


x x x x x 1 x 0



x x x x x 1 x 0





$x \ x \ x \ x \ x \ 1 \ x \ 0$



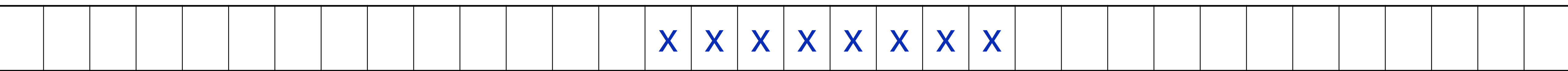
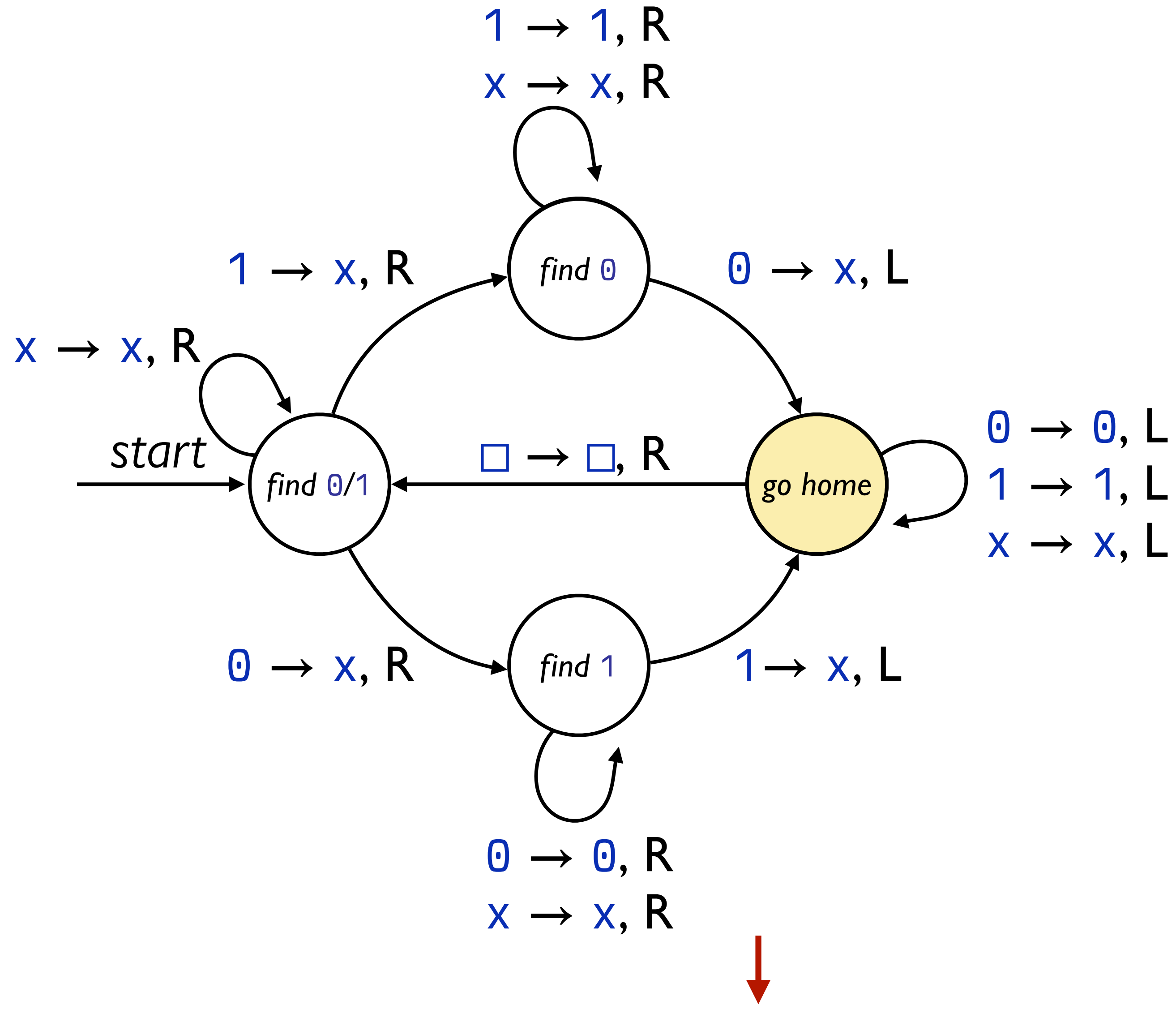


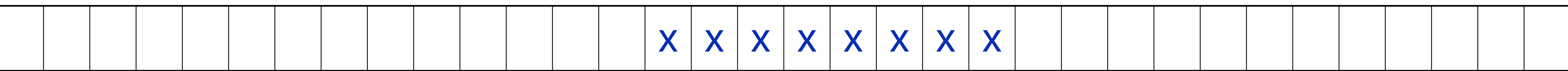
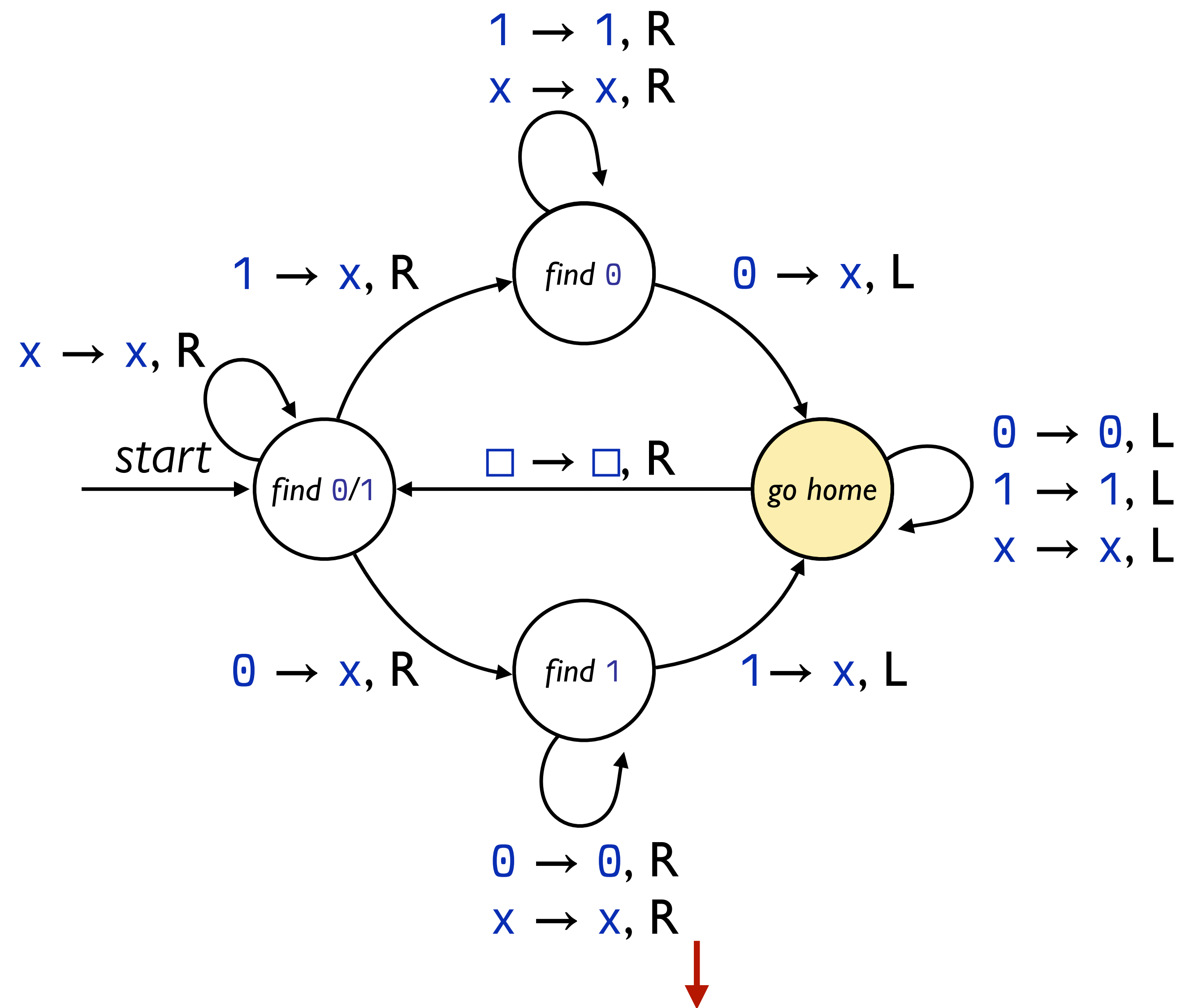


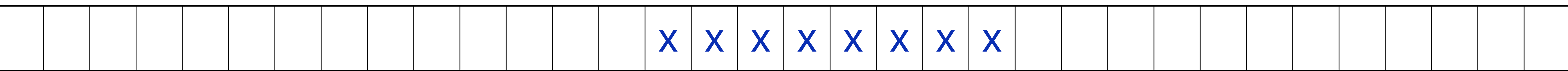
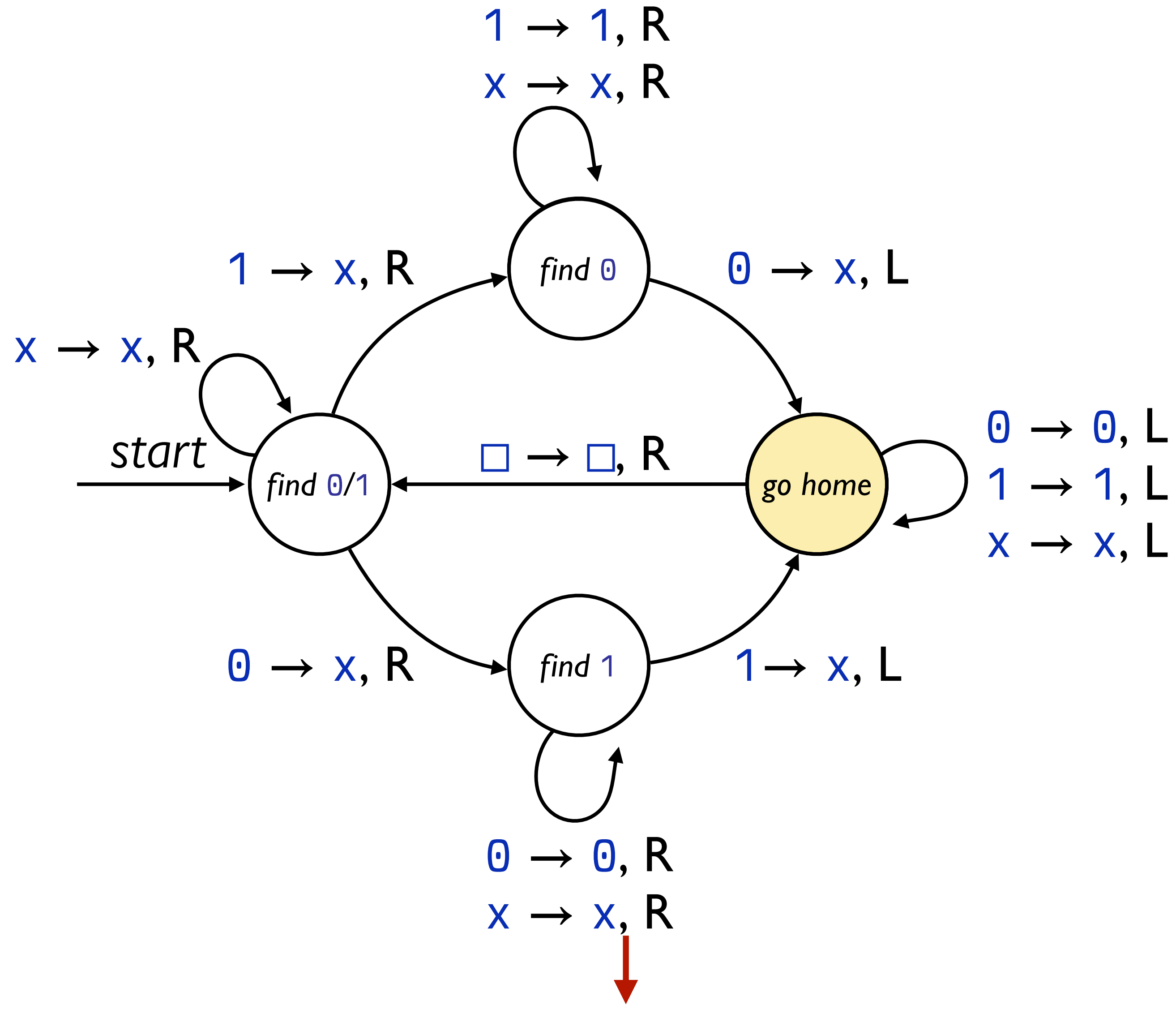


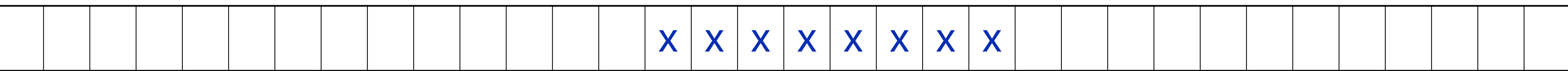
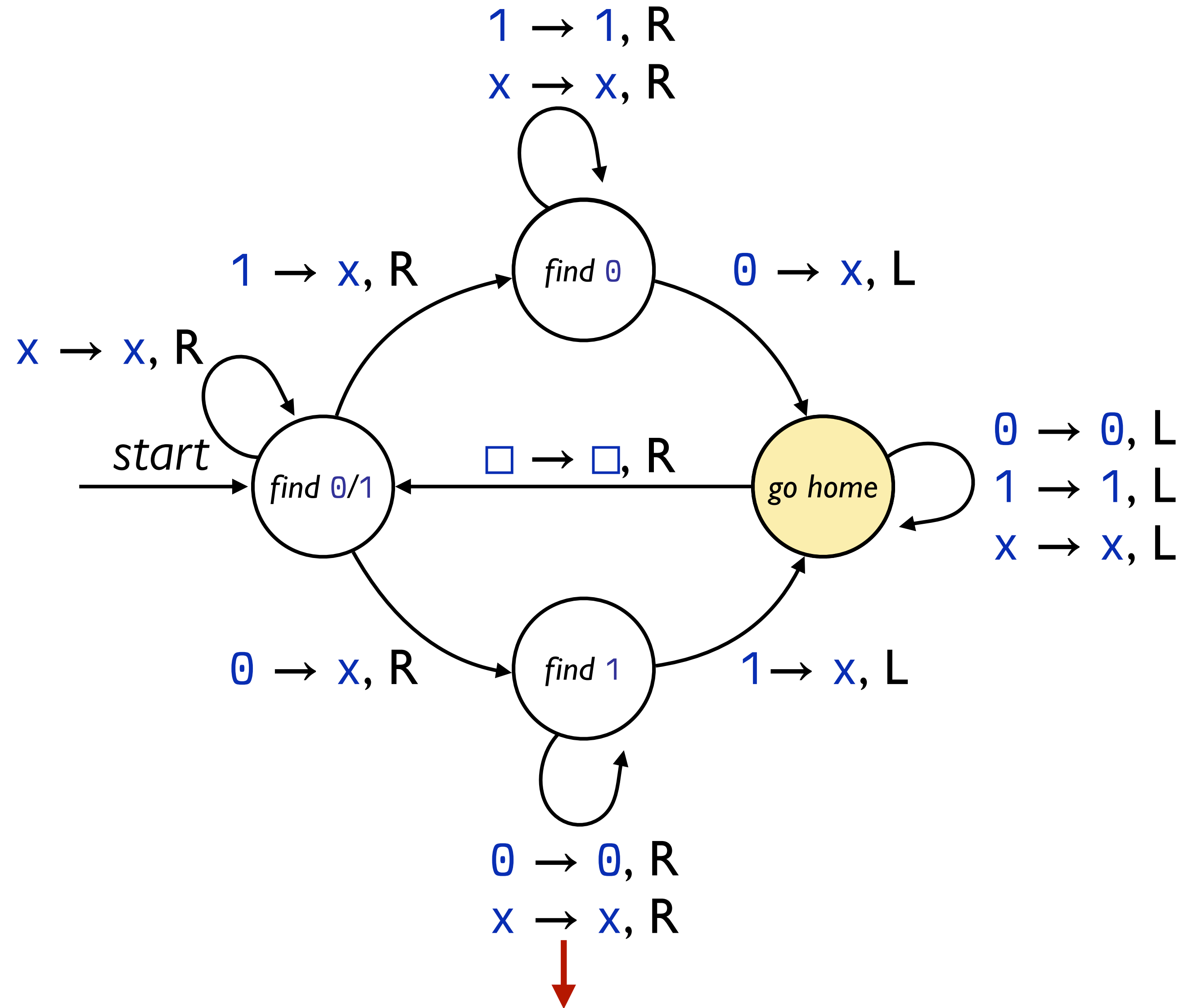


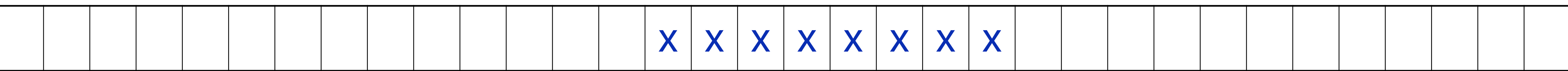
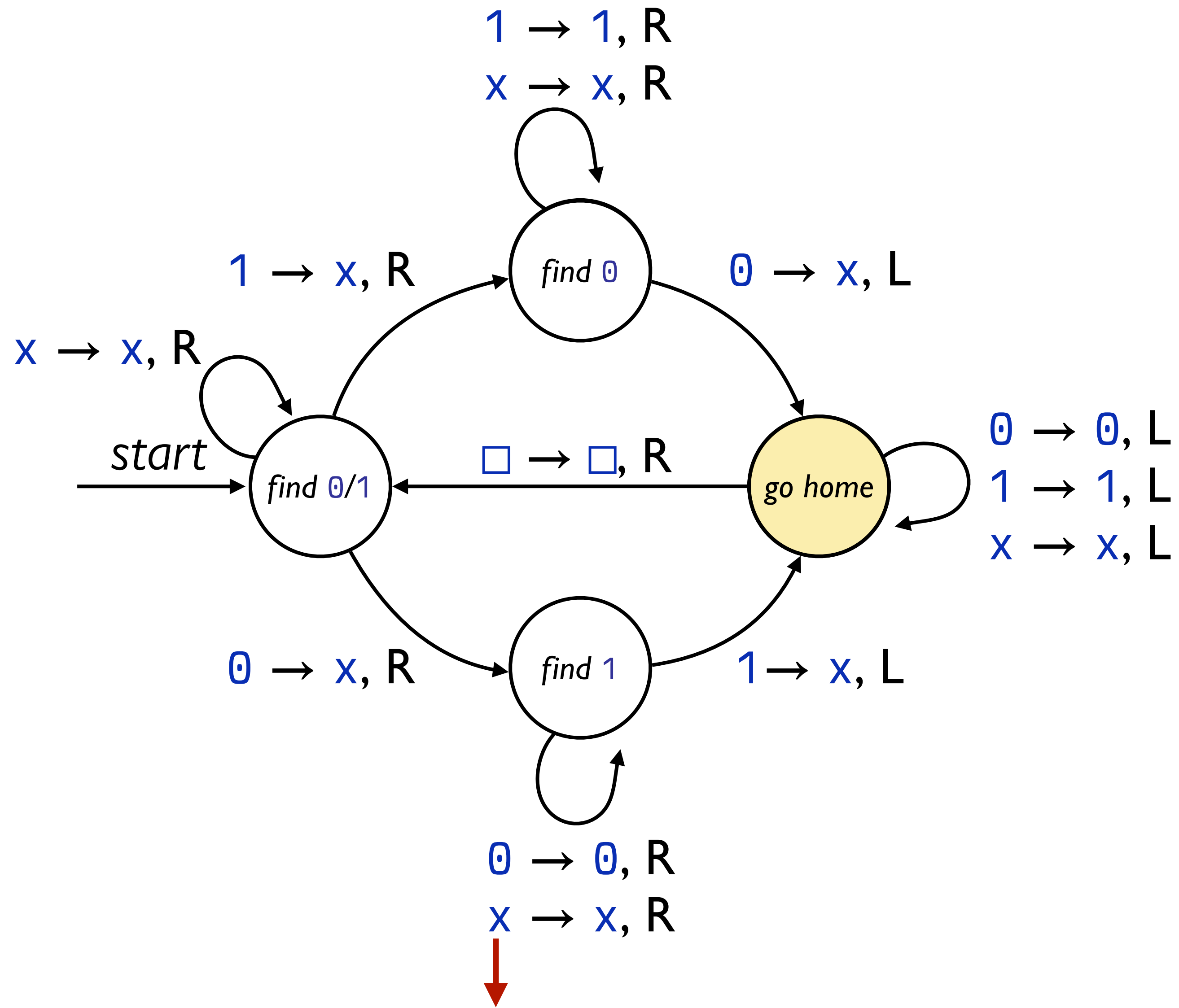








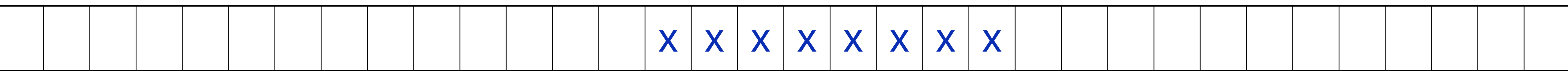
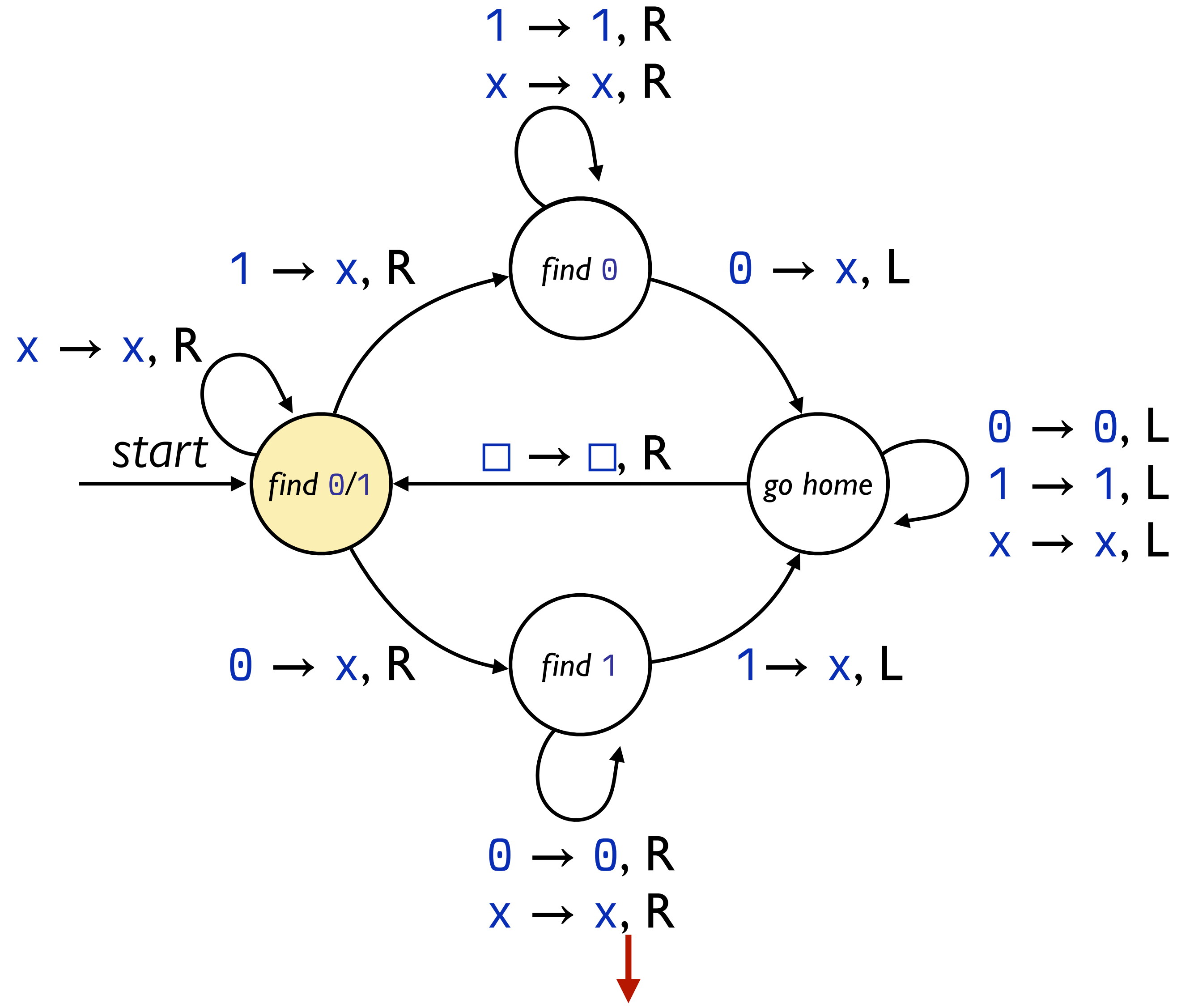


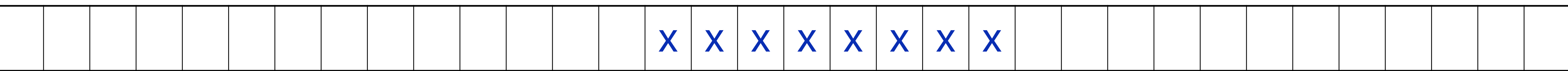
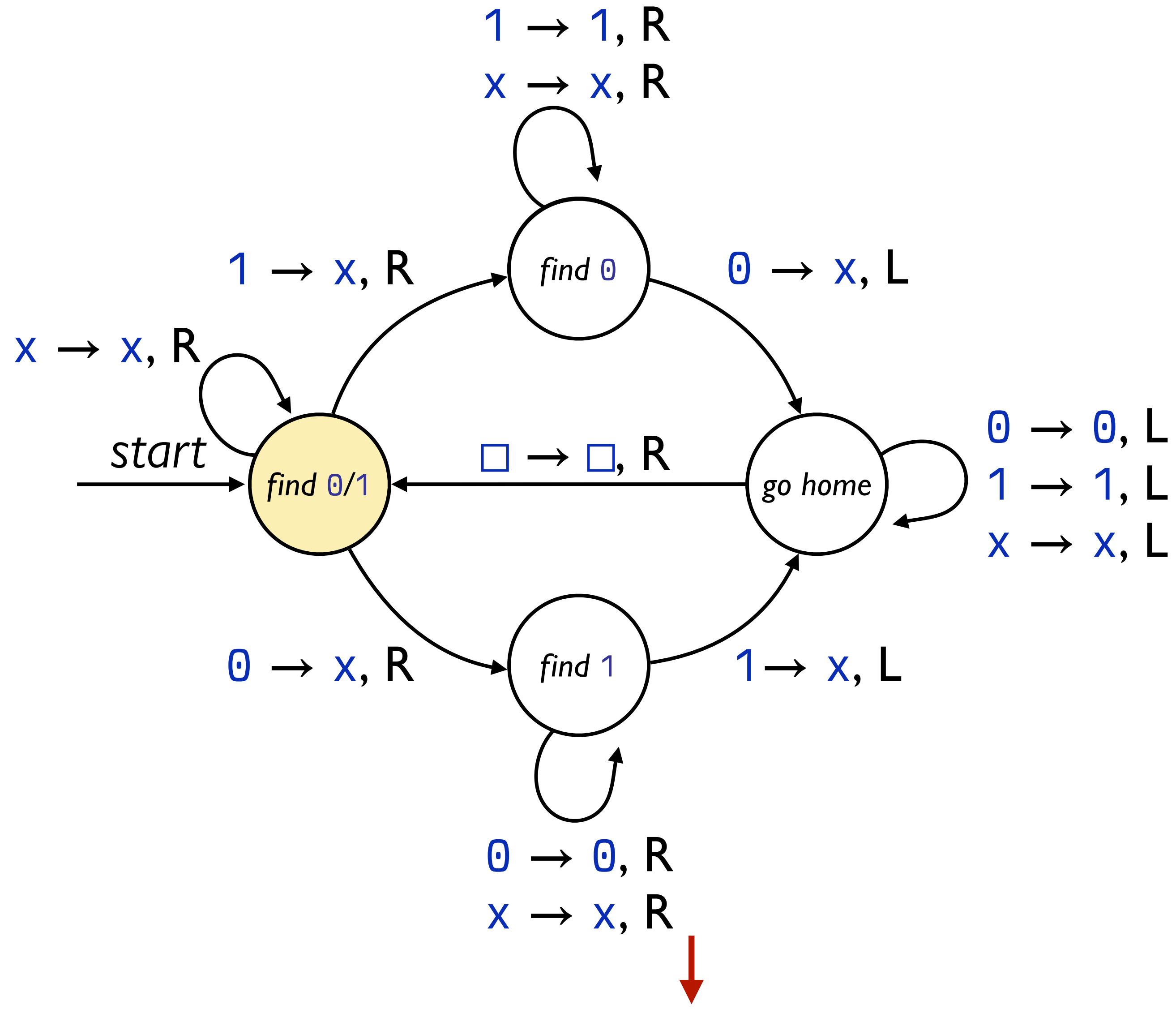


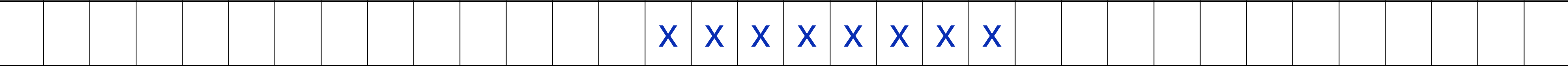
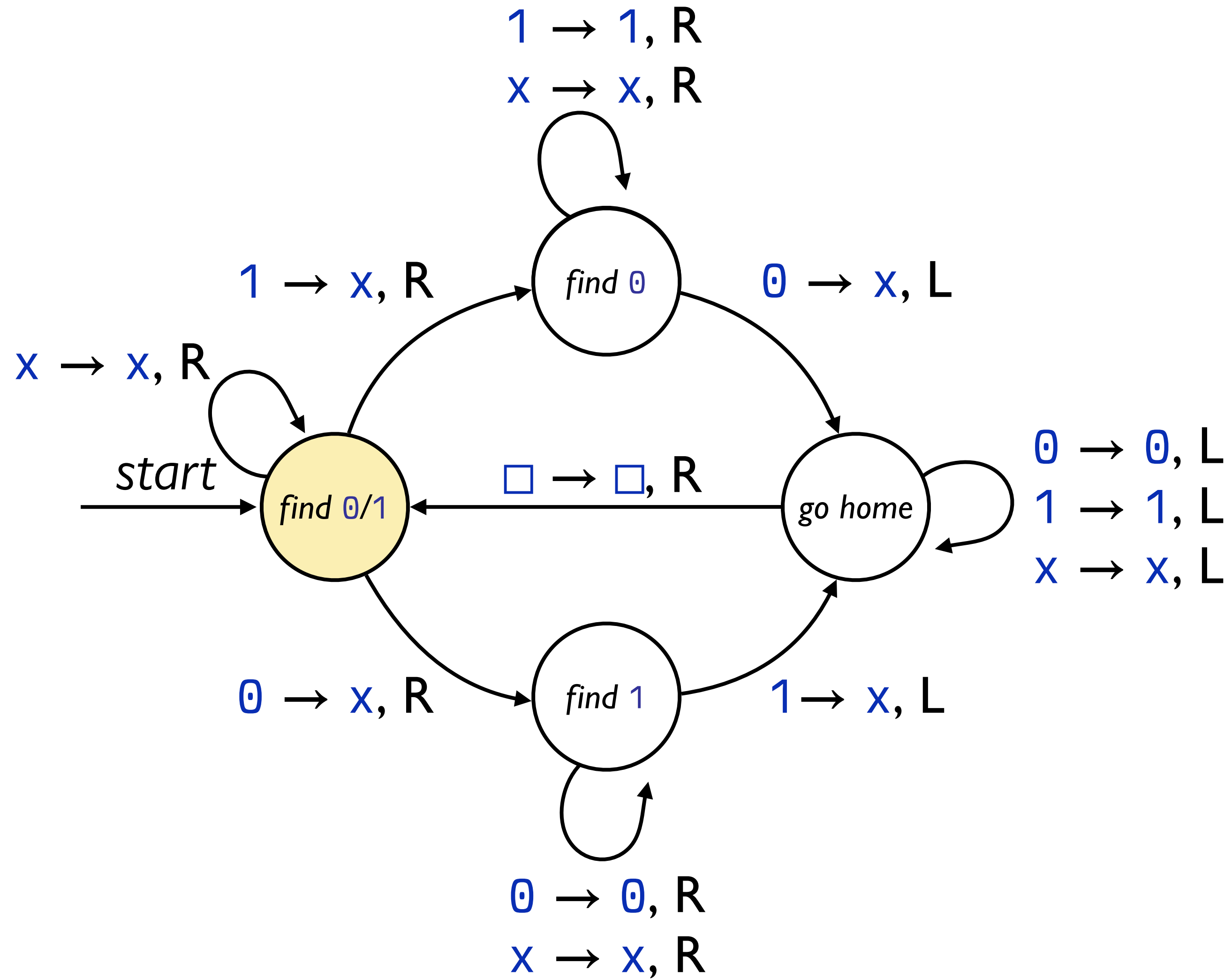


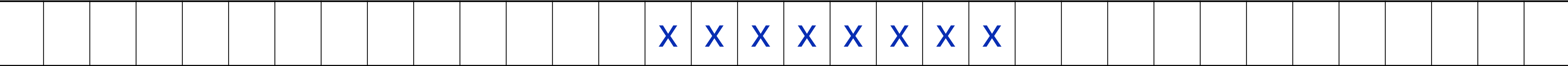
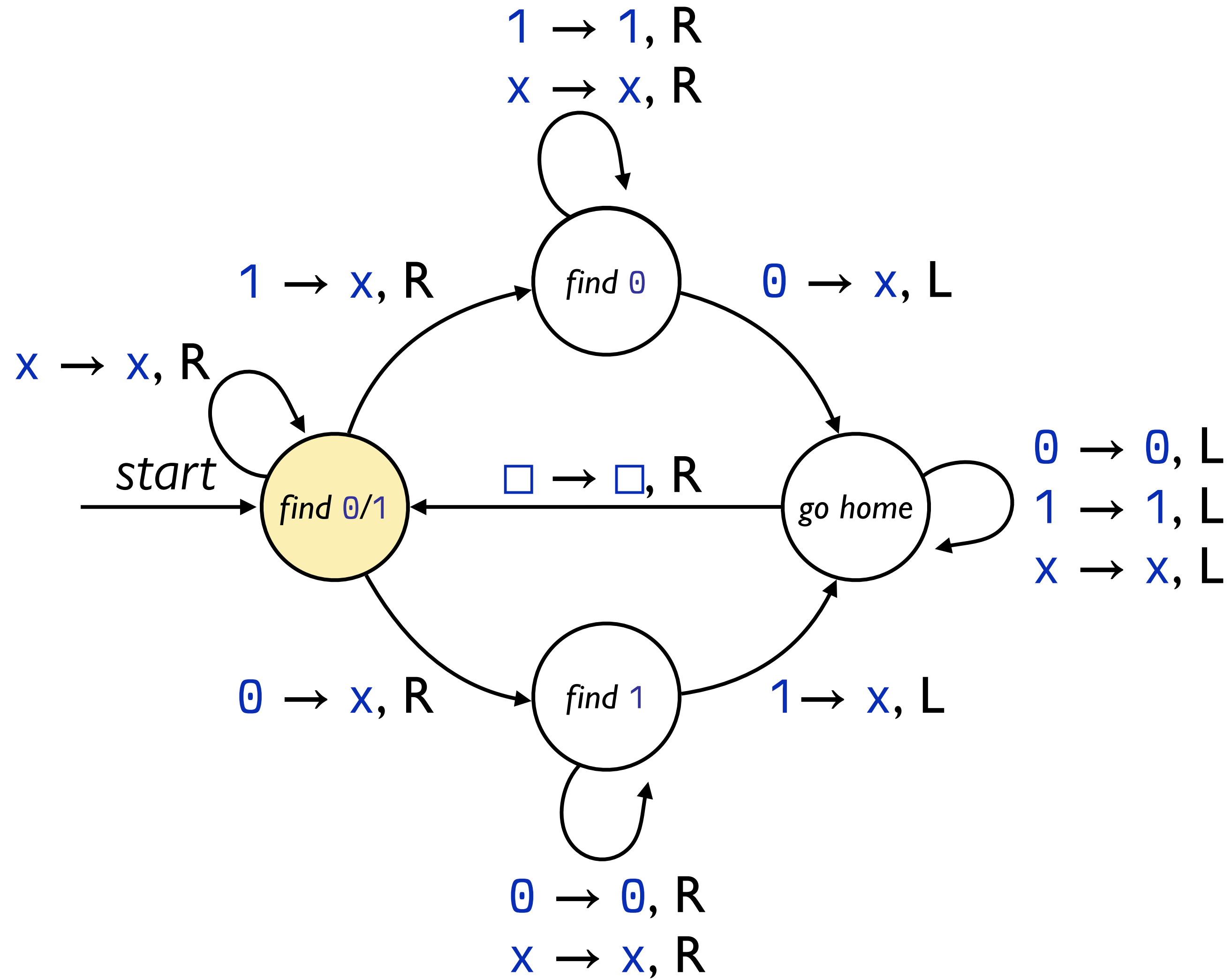










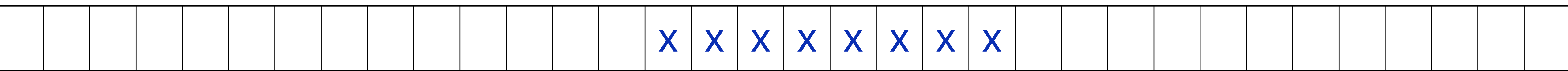
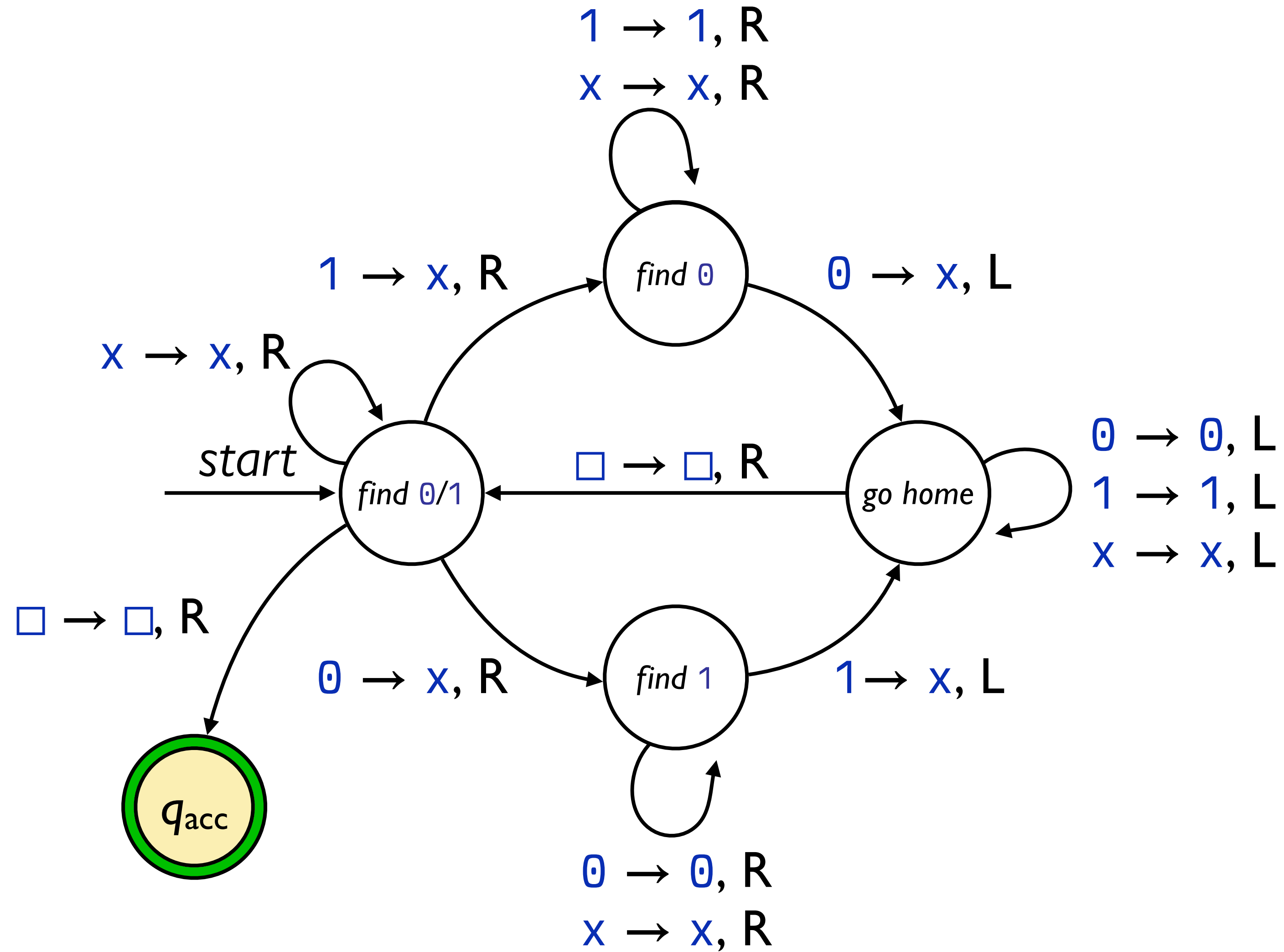




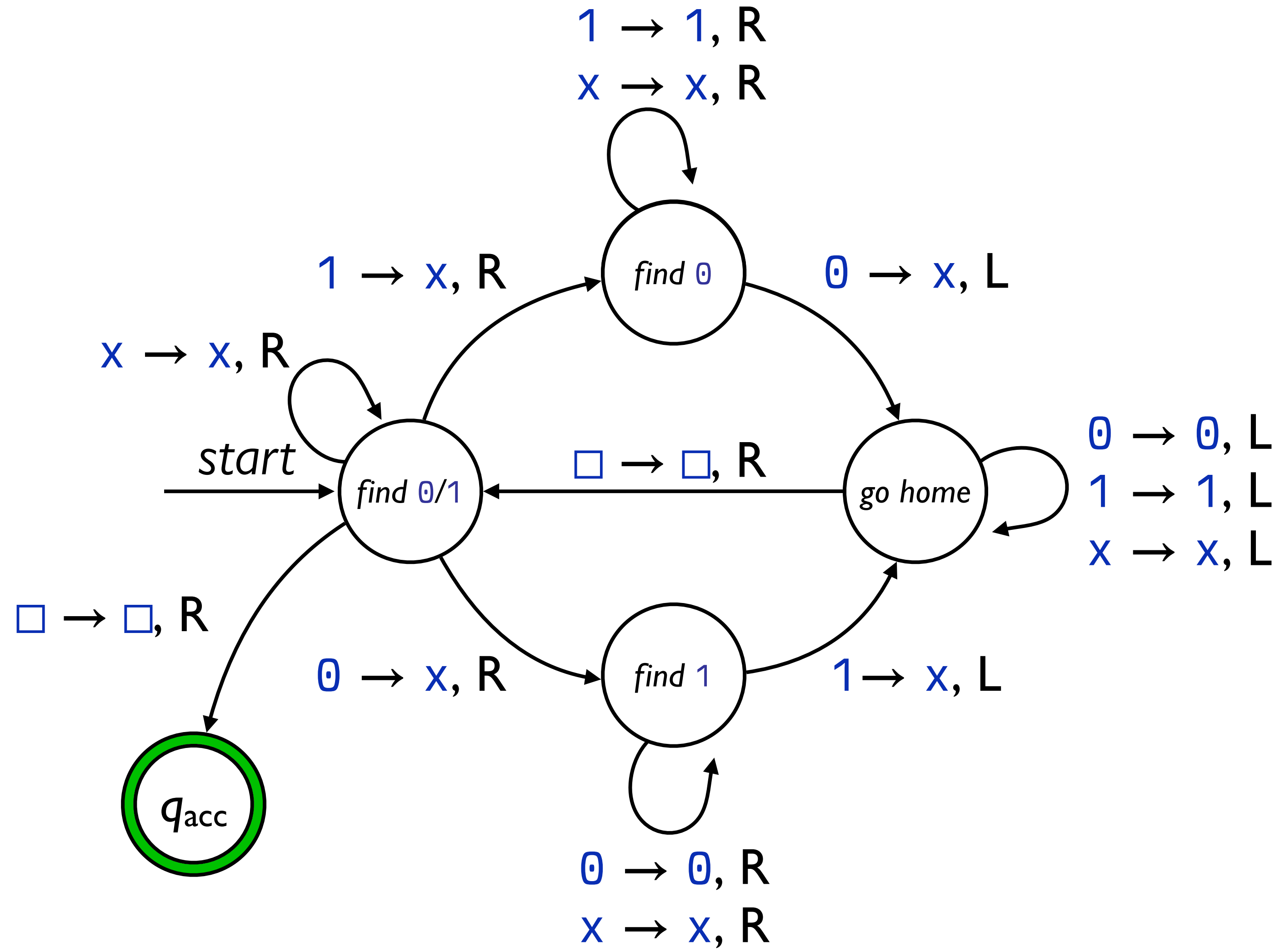


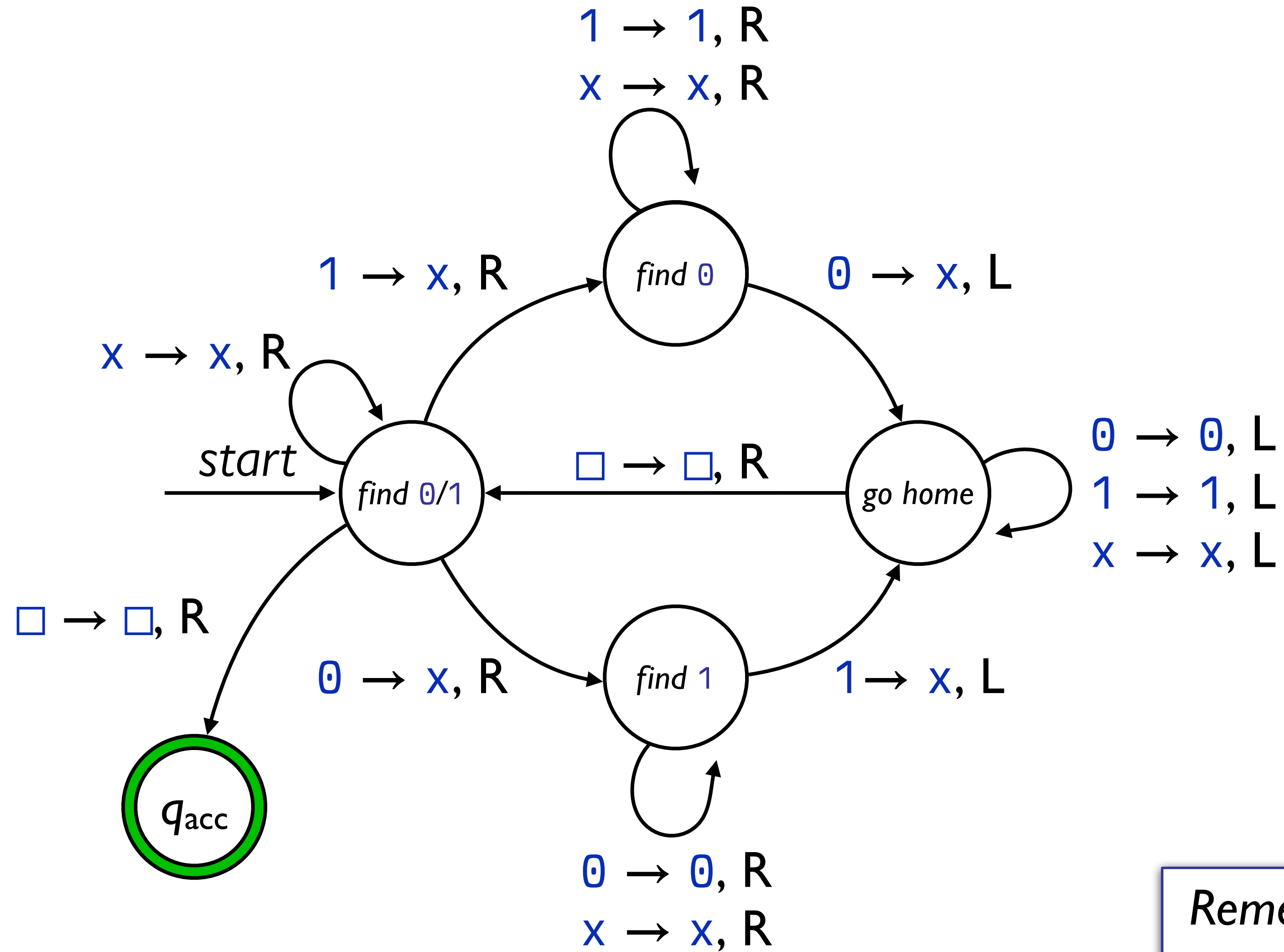




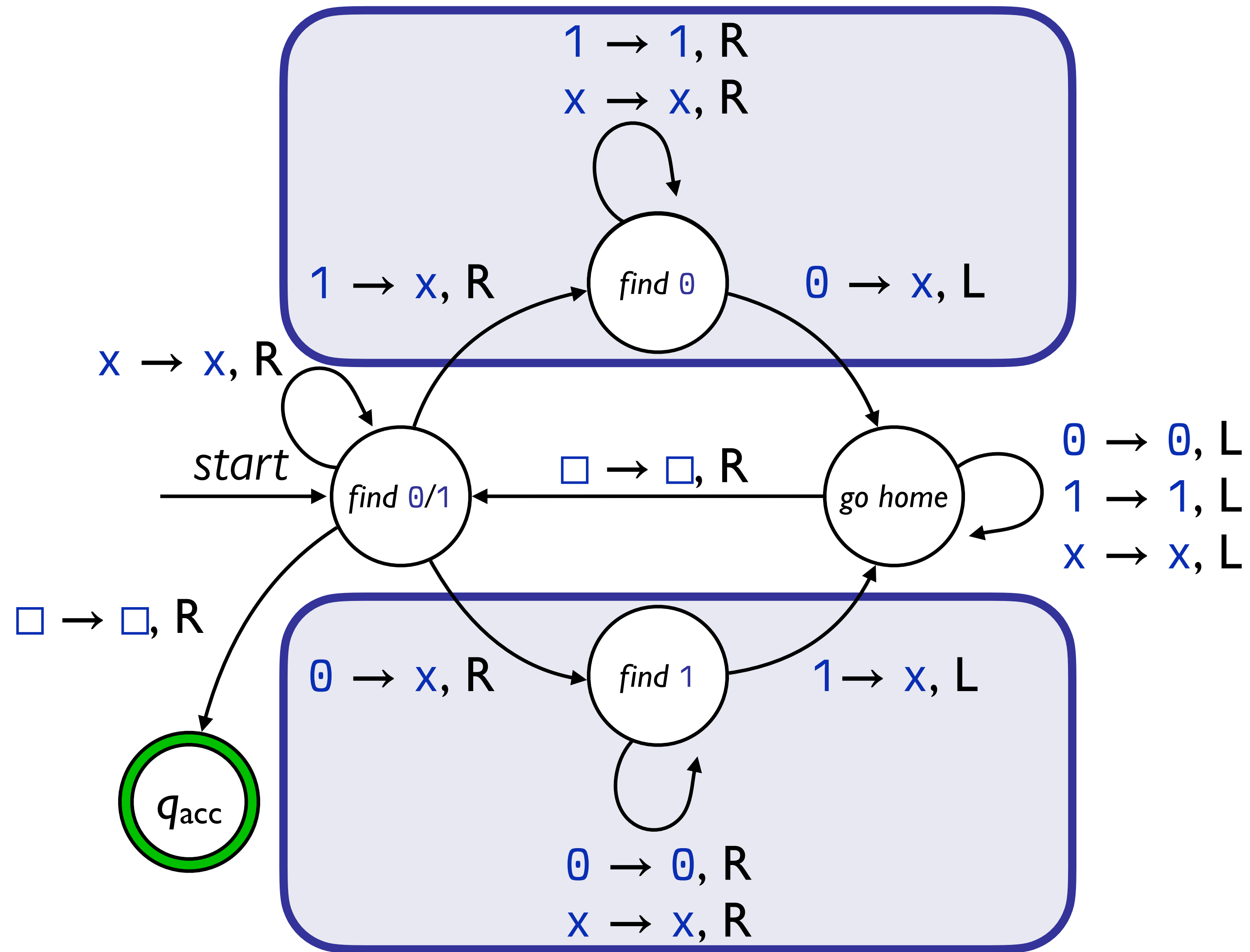








Remember that all missing transitions implicitly reject.



*These parts of the Turing machine are two “cases” it handles, though the tapes they operate on are indistinguishable.*

# Constant storage

Sometimes a Turing machine needs to remember some additional information that can't (at least conveniently) be put on the tape.

In this case, you can use the same techniques you used in designing DFAs and introduce extra states into the Turing machine's finite-state control.

The finite-state control can only remember one of finitely many things, but that might be all you need!

One more TM for practice

Design a Turing machine for the language  $\{ww^R \mid w \in \{a, b\}^*\}$



